

Division - Direct Method

In IEEE floating point standard a real number is represented as :

$$(-1)^S \times M \times 2^E$$

In 32-bit representation :

$$\begin{array}{l|l} S \in \{0, 1\} & \\ M \in [1, 2[& \text{Or } M \in [0, 2[\\ E \in \{-126, \dots, 127\} & \text{if } E = -127 \end{array}$$



normal

sub-normal

Square Root

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If U or V are sub-normal numbers, the calculation of $\frac{U}{V}$ may lead to a lost of precision

Therefore if $M_U = 0$ $\frac{U}{V} = 0$

and if M_U or $M_V \in]0, 1[$

E_U or E_V are decreased and M_U or M_V are $\times 2$ until they can fit within $[1, 2[$



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Let U and V ($V \neq 0$) be two real numbers

We seek to calculate the real number $\frac{U}{V}$

$$U = (-1)^{S_U} \times M_U \times 2^{E_U}$$

$$V = (-1)^{S_V} \times M_V \times 2^{E_V}$$

then $\frac{U}{V} = (-1)^{S_U \oplus S_V} \times \frac{M_U}{M_V} \times 2^{E_U - E_V}$



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Since $M_U \in [1, 2[$ and $M_V \in [1, 2[$ $\frac{M_U}{M_V} \in]0.5, 2[$

Therefore if $M_U < M_V$

E_U is decreased and M_U is $\times 2$ so $\frac{M_U}{M_V} \in [1, 2[$



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Then, the problem can be stated as :

Given two positive real numbers $A \in [1, 4[$ and $B \in [1, 2[$ such as $A \geq B$ we seek to calculate $X_{Th} = \frac{A}{B}$ $X_{Th} \in [1, 2[$

Let X_n be an approximation of X_{Th} coded on $n+1$ bits

$$X_n = \sum_{j=0}^n x_{-j} \times 2^{-j} \quad \text{such as} \quad X_n \leq \frac{A}{B} < (X_n + 2^{-n})$$



We propose to calculate X_n digit-by-digit

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Let $X_k = \sum_{j=0}^n x_{-j} \times 2^{-j}$ such as $X_k \leq \frac{A}{B} < (X_k + 2^{-k})$

At each iteration X_k is obtained from X_{k-1}

$$X_k = X_{k-1} + x_{-k} 2^{-k} \quad x_{-k} \in \{0, 1\}$$

2^{-k} is denoted W_k

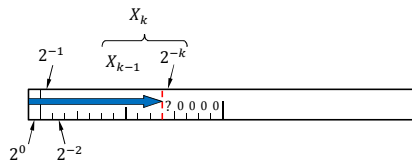


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Let $X_{k-1} = \sum_{j=0}^{k-1} x_{-j} \times 2^{-j}$ such as $X_{k-1} \leq \frac{A}{B} < (X_{k-1} + 2^{-(k-1)})$



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$$X_k \leq \frac{A}{B} < (X_k + 2^{-k})$$

$$BX_k \leq A < B(X_k + 2^{-k})$$

$$0 \leq A - BX_k < 2^{-k}B \quad \text{Let } \Delta_k = A - BX_k$$

$$0 \leq \Delta_k < 2^{-k}B \quad \text{yet } B < 2$$

then $0 \leq \Delta_k < 2^{-k} \times 2$

- $0 \leq \Delta_0 < 2$
- At each iteration the upper bound of Δ_k is divided by 2



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$$\Delta_k = A - BX_k$$


$$\Delta_k = A - B(X_{k-1} + x_{-k}2^{-k})$$

$$\Delta_k = A - BX_{k-1} - x_{-k}2^{-k}B$$

$$\Delta_k = \Delta_{k-1} - x_{-k}2^{-k}B$$

$$2^k \Delta_k = 2^k \Delta_{k-1} - x_{-k}B \quad \text{Let } D_k = 2^k \Delta_k$$

$$D_k = 2D_{k-1} - x_{-k}B \quad x_{-k} = \begin{cases} 0 \\ 1 \end{cases} \text{ such as } 0 \leq D_k$$



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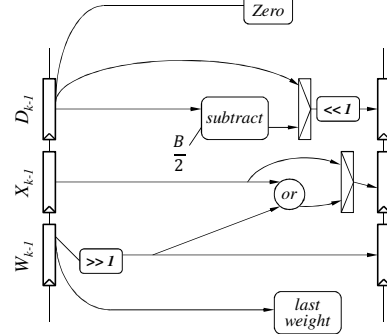
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Implementation

$$D_k = 2 \left(D_{k-1} - x_{-k} \frac{B}{2} \right)$$

$$X_k = X_{k-1} + x_{-k} 2^{-1} W_{k-1}$$

$$W_k = 2^{-k} = 2^{-1} W_{k-1}$$






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Iteration scheme :

$$\begin{cases} D_k = 2 \left(D_{k-1} - x_{-k} \frac{B}{2} \right) \\ X_k = X_{k-1} + x_{-k} 2^{-k} \end{cases} \quad x_{-k} = \begin{cases} 0 \\ 1 \end{cases} \text{ such as } 0 \leq D_k$$

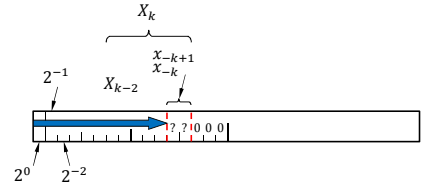



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Division - Direct Method - Improvement

radix 4

$$X_k = X_{k-2} + (2x_{-k+1} + x_{-k})2^{-k}$$




Square Root

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radix 4


$$\Delta_k = A - BX_k$$

$$\Delta_k = A - B(X_{k-2} + (x_{-k+1}2^{-k+1} + x_{-k}2^{-k}))$$

$$\Delta_k = \Delta_{k-2} - B(x_{-k+1}2^{-k+1} + x_{-k}2^{-k})$$

$$2^k \Delta_k = 2^k \Delta_{k-2} - B(2x_{-k+1} + x_{-k}) \quad \text{Let } D_k = 2^k \Delta_k$$


$$D_k = 4 \left(D_{k-2} - \frac{B}{4} (2x_{-k+1} + x_{-k}) \right)$$



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Square Root

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radix 4

Iteration scheme :

$$\begin{cases} D_k = 4 \left(D_{k-2} - \frac{B}{4} (2x_{-k+1} + x_{-k}) \right) \\ X_k = X_{k-2} + (2x_{-k+1} + x_{-k})2^{-k} \end{cases}$$


$$x_{-k+1}, x_{-k} = \begin{cases} 00 \\ 01 \\ 10 \\ 11 \end{cases} \text{ such as } 0 \leq D_k$$



Square Root

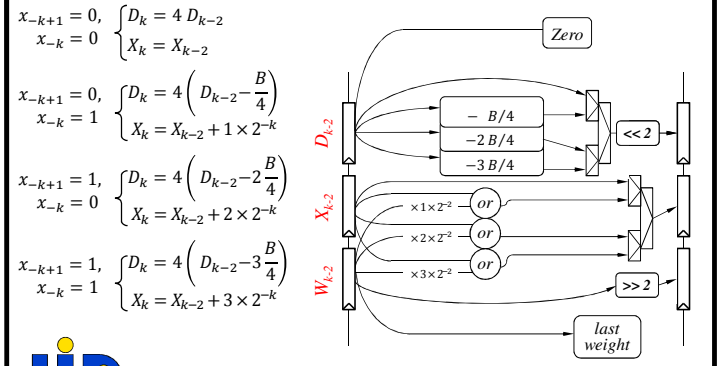
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radix 8

$$X_k = X_{k-3} + (4x_{-k+2} + 2x_{-k+1} + x_{-k})2^{-k}$$

LIP 6

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Division - Direct Method - Improvement

radix 8

Iteration scheme :

$$\begin{cases} D_k = 8 \left(D_{k-3} - \frac{B}{9} (4x_{-k+2} + 2x_{-k+1} + x_{-k}) \right) \\ X_k = X_{k-3} + (4x_{-k+2} + 2x_{-k+1} + x_{-k})2^{-k} \end{cases}$$

$$x_{-k+2}, x_{-k+1}, x_{-k} = \begin{cases} 000 \\ 001 \\ 010 \\ 011 \\ 000 \\ 001 \\ 010 \\ 011 \end{cases} \text{ such as } 0 \leq D_k$$

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radix 8

$$\Delta_k = A - BX_k$$

$$\Delta_k = A - B(X_{k-3} + (x_{-k+2}2^{-k+2} + x_{-k+1}2^{-k+1} + x_{-k}2^{-k}))$$

$$\Delta_k = \Delta_{k-3} - B(x_{-k+2}2^{-k+2} + x_{-k+1}2^{-k+1} + x_{-k}2^{-k})$$

$$2^k \Delta_k = 2^k \Delta_{k-3} - B(4x_{-k+2} + 2x_{-k+1} + x_{-k}) \quad \text{Let } D_k = 2^k \Delta_k$$

$$D_k = 8 \left(D_{k-3} - \frac{B}{8} (4x_{-k+2} + 2x_{-k+1} + x_{-k}) \right)$$

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$$D_k = 8 \left(D_{k-3} - \frac{B}{8} (4x_{-k+2} + 2x_{-k+1} + x_{-k}) \right)$$

$$X_k = 8 (X_{k-3} + (4x_{-k+2} + 2x_{-k+1} + x_{-k}))$$

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