



The Abdus Salam  
International Centre for Theoretical Physics



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## School and Workshop on Structure and Function of Complex Networks

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### Modeling Strategic Formation of Social Networks

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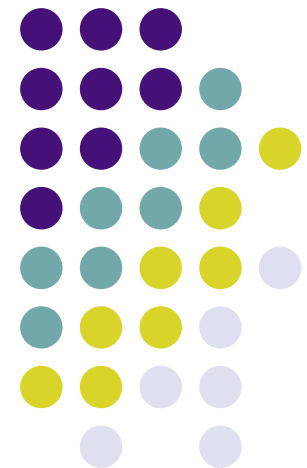
These are preliminary lecture notes, intended only for distribution to participants

# Modeling Strategic Formation of Social Networks

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Matthew O Jackson

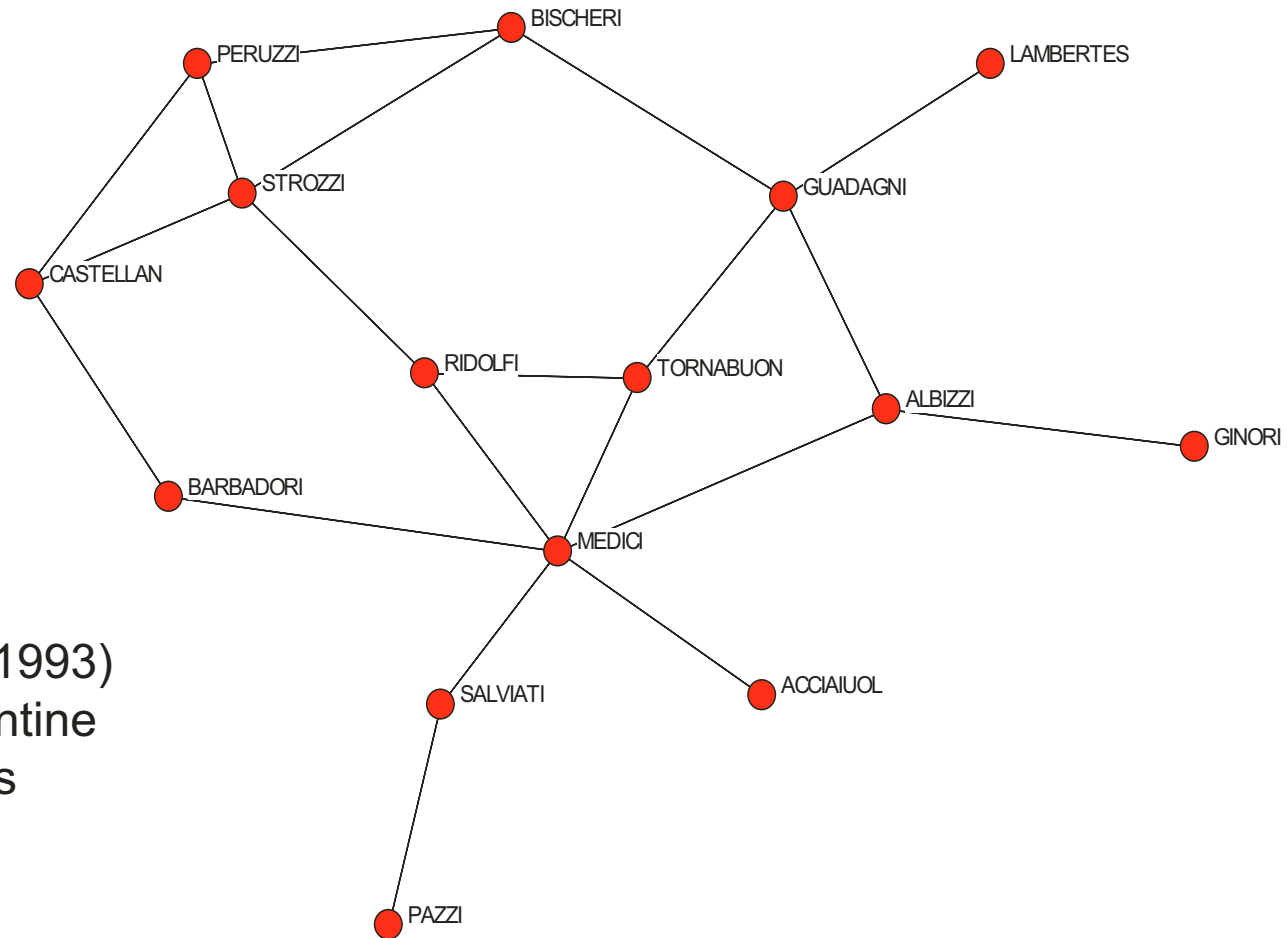
Trieste Summer School on Networks



# Example of Social Network with Strategic Formation



● PUCCI



Padgett -Ansell (1993)  
[from Kent] Florentine  
Marriages, 1430's



# What do we know?

- Networks are prevalent
  - Job contact networks, crime, trade, politics, ...
- Network position and structure matters
  - rich sociology literature
  - Padgett example – Medicis not the wealthiest nor the strongest politically, but the most central
- “Social” Networks have special characteristics
  - small worlds, degree distributions...

# Questions:



- How does network structure affect interaction and behavior? (tomorrow, Tuesday)
- Which networks form?
  - random modeling
  - Game theoretic reasoning
- When do efficient networks form?
  - Intervention - design incentives?

# Economic/Game Theoretic Modeling - distinguishing ingredients



- Welfare analysis – agents get utility from networks
- Decision making agents form links and/or choose actions

# Notation



- $\{1, \dots, i, \dots, j, \dots, n\}$  – nodes or players
- $g$  - network,  $n$  by  $n$  matrix with entries 0 or 1
- $ij$  in  $g$  iff  $g_{ij}=1$
- $g+ij$  – add link/edge  $ij$  to  $g$ ;  $g-ij$  – delete link  $ij$  from  $g$
- $u_i(g)$  utility to  $i$  from network  $g$



# Efficiency and Stability

(Jackson and Wolinsky (1996))

- Efficiency:
  - $\operatorname{argmax} \sum u_i(g)$
- Pairwise Stable networks:
  - $u_i(g) \geq u_i(g-ij)$  for each  $i$  and  $ij$  in  $g$
  - $u_i(g+ij) > u_i(g)$  implies  $u_j(g+ij) < u_j(g)$  for each  $ij$  not in  $g$





# “Connections Model”

- $0 \leq \delta_{ij} \leq 1$  a benefit parameter for  $i$  from path connection between  $i$  and  $j$
- $0 \leq c_{ij}$  cost to  $i$  of link to  $j$
- $d(i,j)$  shortest path length between  $i,j$

$$u_i(g) = \sum_j \delta_{ij}^{d(i,j)} - \sum_{j \text{ in } N_i(g)} c_{ij}$$

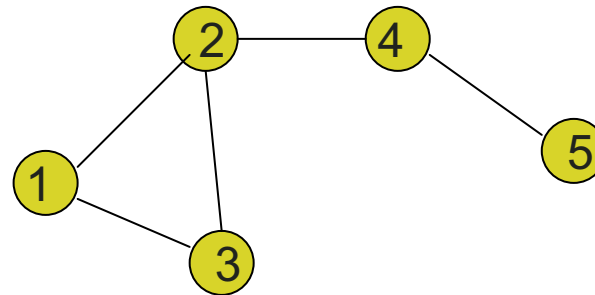
# Example: Symmetric Connections Model



- benefit from a friend is  $\delta < 1$
- benefit from a friend of a friend is  $\delta^2, \dots$
- cost of a link is  $c > 0$

$$u_2 = 3\delta + \delta^2 - 3c$$

$$u_1 = 2\delta + \delta^2 + \delta^3 - 2c$$



$$u_5 = \delta + \delta^2 + 2\delta^3 - c$$

# Efficient Networks



- low cost:  $c < \delta - \delta^2$ 
  - complete network is efficient
- medium cost:  $\delta - \delta^2 < c < \delta + (n-2)\delta^2/2$ 
  - star network is efficient
- high cost:  $\delta + (n-2)\delta^2/2 < c$ 
  - empty network is efficient



# “Proof”

- $c < \delta - \delta^2$  - obvious that complete is efficient
- $\delta - \delta^2 < c$ 
  - players who are connected should be connected in a star:
    - minimal number of links to connect
    - connection at length 2 is more valuable than at 1 ( $\delta - c < \delta^2$ )
  - value of a star is
    - $2(n-1)\delta + (n-1)(n-2)\delta^2 - 2(n-1)c = 2(n-1) [\delta + (n-2)\delta^2/2 - c]$
  - Star of size  $m+n$  more valuable than  $\text{star}(m) + \text{star}(n)$
  - Star has positive value only when  $c < \delta + (n-2)\delta^2/2$

# Pairwise Stable Networks:

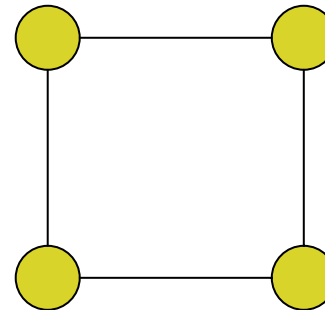


- low cost:  $c < \delta - \delta^2$ 
  - complete network is pairwise stable
- medium/low cost:  $\delta - \delta^2 < c < \delta$ 
  - star network is pairwise stable
  - others are also pairwise stable
- medium/high cost:  $\delta < c < \delta + (n-2)\delta^2/2$ 
  - star network is not pairwise stable (no loose ends)
  - nonempty pairwise stable networks are over-connected and may include too few agents
- high cost:  $\delta + (n-2)\delta^2/2 < c$ 
  - empty network is pairwise stable

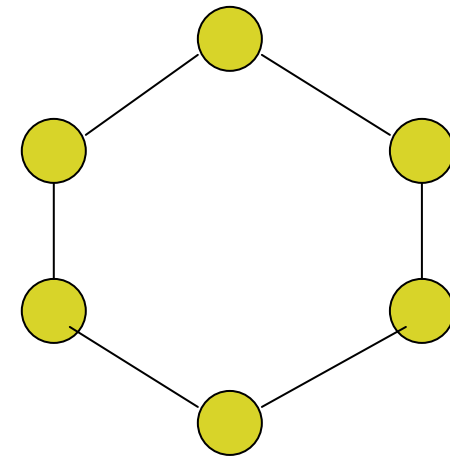
# Example – Pairwise Stable but not efficient



- $\delta - \delta^2 < c < \delta - \delta^3$ ,  $n=4$



- $\delta - \delta^3 < c < (\delta + \delta^2 + \delta^3)(1 - \delta^2)$



(unique nonempty pairwise stable structure if  $\delta < c < (\delta + \delta^2 + \delta^3)(1 - \delta^2)$ ,  $n=6$ )



# Directed Connections

[Bala and Goyal (2000)]

- same payoffs as before except
  - directed network and
    - one way flow – link is only useful to whom incurs cost
    - two way flow – one player pays, but link is useful to both
- Now links are formed unilaterally
  - use Nash equilibrium to model stability

# Two way flow



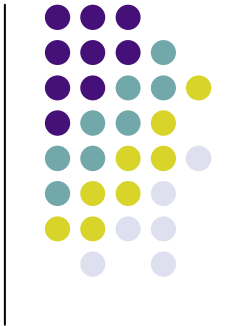
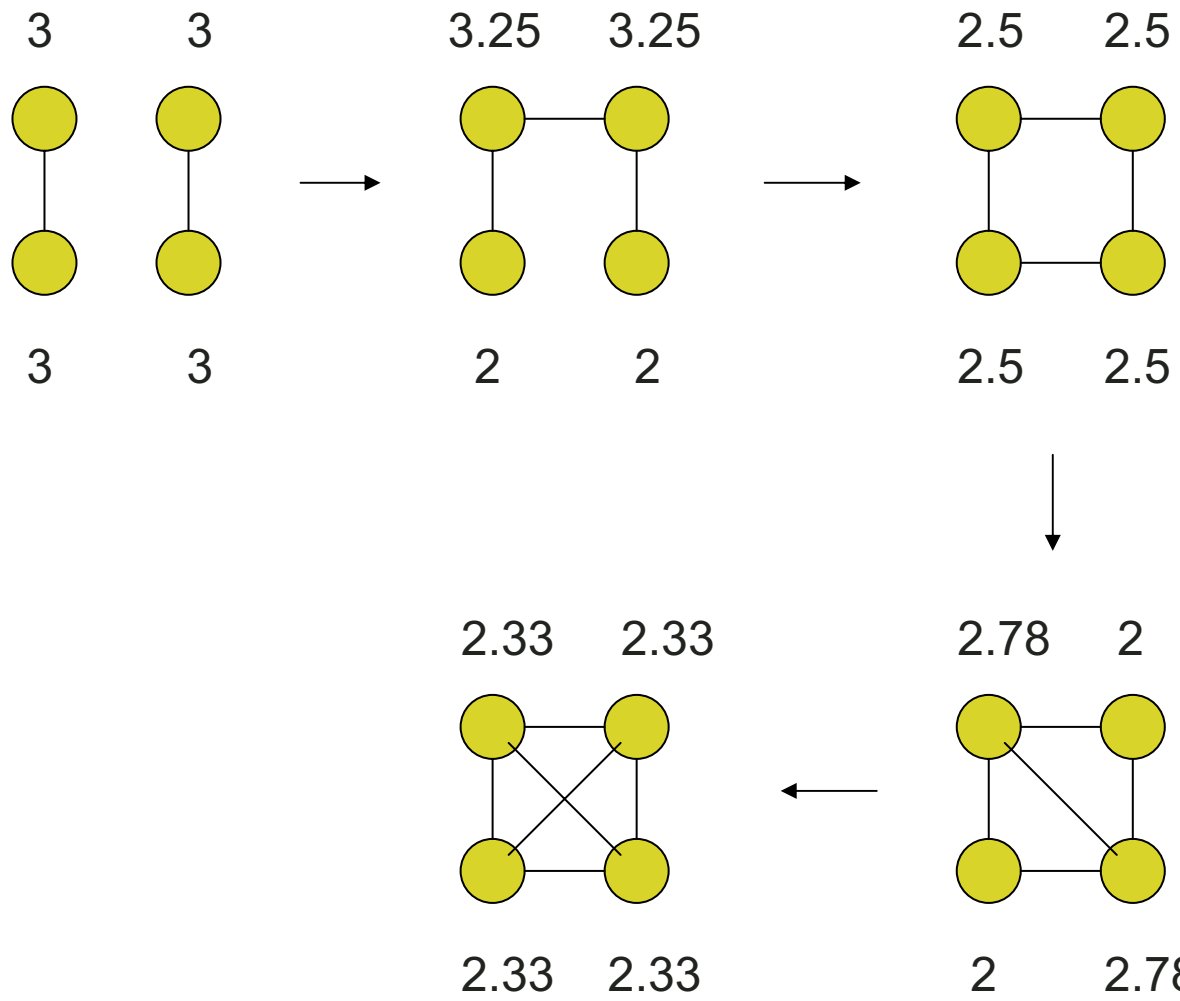
- Efficiency as before, except  $c/2$  and link in either direction (but not both)
- Nash Stable:
  - low cost:  $c < \delta - \delta^2$ 
    - two-way “complete” networks are pairwise stable
  - medium/low cost:  $\delta - \delta^2 < c < \delta$ 
    - all star networks are pairwise stable, plus others
  - medium/high cost:  $\delta < c < \delta + (n-2)\delta^2/2$ 
    - peripherally sponsored star networks are stable (no other stars, but sometimes other networks)
  - efficient and stable can be empty:
    - $\delta - \delta^2 < c < 2(\delta - \delta^2)$  complete is efficient, not equilibrium





# A toy co-author model

- $u_i(g) = \sum_{j: ij \text{ in } g} [1/n_i + 1/n_j + 1/(n_i n_j)]$
- $n$  is even:
  - efficient networks: pairs
  - pairwise stable networks: completely connected components, where for any two components, one has more than the square of the number of nodes in the other



# Some Settings stable=efficient



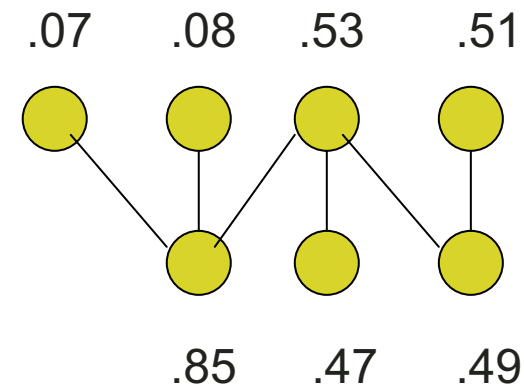
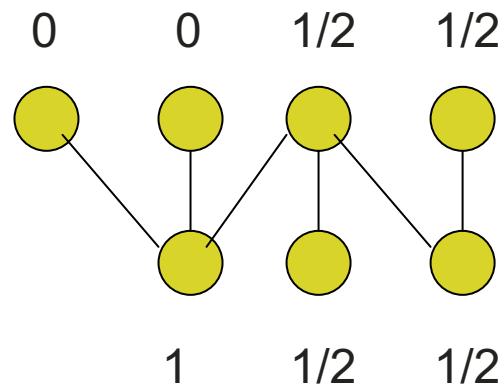
Buyer-Seller Networks: Corominas Bosch (2002):

- Sellers each with one identical object value 0
- Buyers each desire one object at value 1
- Alternating offers bargaining



# Experiments

## Charness, C-B, and Frechette



(normalized from their 2500 base and 200 reservation value)

# general algorithm for solution



- is there group of at least two sellers connected to just one buyer?
  - if so, remove, buyer gets 1, sellers 0
- is there group of at least  $k+1$  sellers connected to just  $k$  buyers?
  - if so, remove, buyers get 1, sellers 0
- iterate
- repeat with buyers/sellers reversed
- remaining players get  $1/2$



# Pairwise stable = efficient

[I] buyer gets 1 implies some linked seller gets 0 (and vice versa)

[II] component – all buyers get  $\frac{1}{2}$  if and only if for all  $k$  and subsets of  $k$  buyers they are linked to at least  $k$  distinct sellers and vice versa

- pairwise stable and cost to link implies no 0's. so by [I] must be that all players get  $\frac{1}{2}$
- if not pairs find subnetwork that links them into pairs – extra links can be deleted and still satisfy [II], same payoff at lower cost



- If bargaining is such that split is not 0,1 in ``unbalanced'' networks, then could get inefficiency



# Stable and Efficient only coincide in special cases



- Can transfers help in other cases?
- What can we say about when conflict exists
- What can we say about transfers helping?
- What about other formation processes?



# Transfers can help

- Change utilities from  $u_i(g)$  to  $u_i(g)+t_i(g)$ 
  - $\sum_{i \text{ in } C} t_i(g) = 0$  within each component  $C$
  - respect anonymity
- E.g., peripheral players pay center of star in connections model to maintain connections

Can Prove: Transfers can lead to at least one efficient network being pairwise stable provided:

- $v$  is anonymous, and
- in some efficient network all nodes have degree  $>1$ .

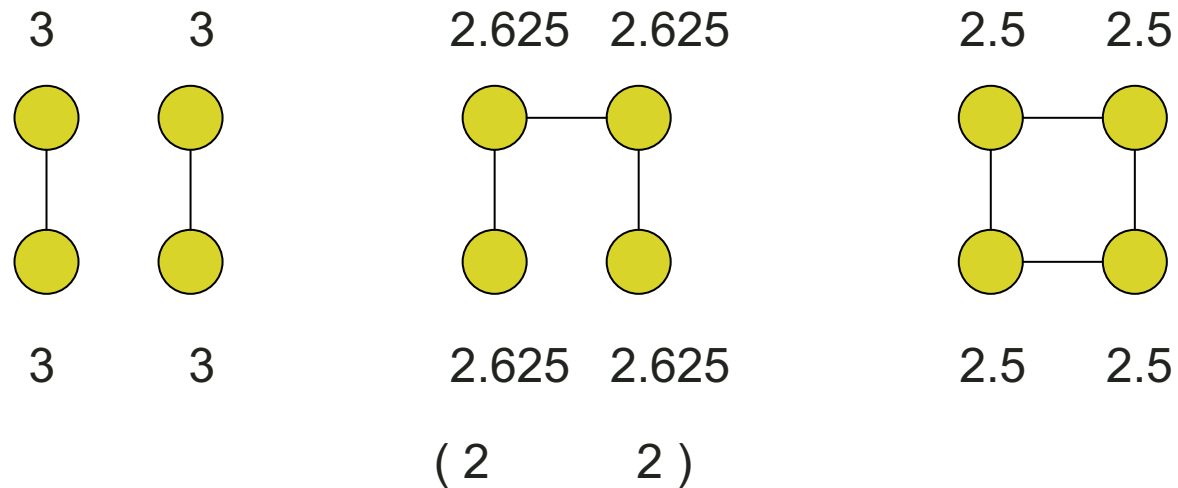
[ugly proof]

# Transfers in Co-author - equalizing works



[tax on having a second link]

( 3.25 3.25 )



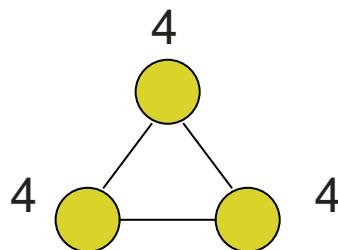
But note: need to charge players who form links - not just subsidize (here penalizing for forming links)

# Transfers cannot always help

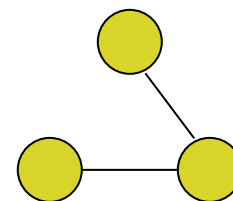
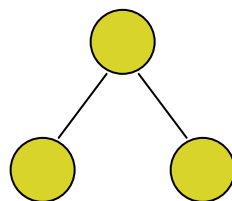
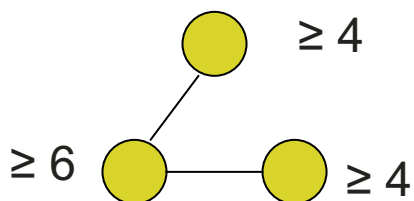


anonymity: same transfers to identical players

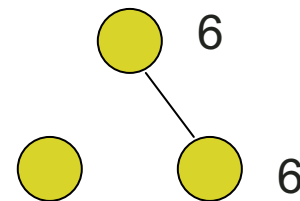
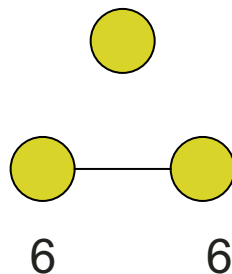
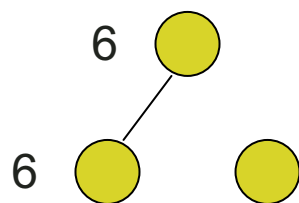
balance: no transfers outside of component



value 12



value 13  
efficient



value 12

# What is needed to avoid this?



- Break Anonymity
- Make transfers part of formation process or allow transfers to be contingent on network

# Bargaining when forming links Can Help



- Currarini and Morelli (2000)
  - Order players (breaks anonymity)
  - Player demands payoff and suggests links
  - Link is formed if both players announce it, and component is feasible in terms of demands



# Bargaining in the example

- 1 announces  $\{2,3\}$  and demands 6.5
  - 1 cannot ask for more than  $x > 6.5$  or 2 can respond with  $\{3\}$  and slightly below  $x$  and will get  $x$  and 1 will get nothing
- 2 announces  $\{1\}$  and demands 6.5
  - 2 cannot say  $\{3\}$  (or  $\{1\}$  or  $\{1,3\}$ ) and ask for more than 6.5
- 3 announces  $\{1\}$  and demands 0
  - 3 cannot do better



If  $v(g+ij) > v(g)$  whenever  $g+ij$  has more components than  $g$ , then all (subgame perfect) equilibria of the above game lead to efficient networks.

- Game is contrived – introduces asymmetries, requires endpoint
- But ability to endogenize transfers as part of the formation process is important



# Can economic models match observables?



- Small worlds derived from costs/benefits
  - low costs to local links – high clustering
  - high value to distant connections – low diameter
  - high cost of distant connections – few distant links

## Geographic Connections (Johnson-Gilles (2000), Carayol-Roux (2003), Galeotti-Goyal-Kamphorst (2004), ...)



### Islands connections model (Jackson-Rogers (2004))

- players live on islands
- cost  $c$  of link to player on the island
- cost  $C > c$  of link to player on another island

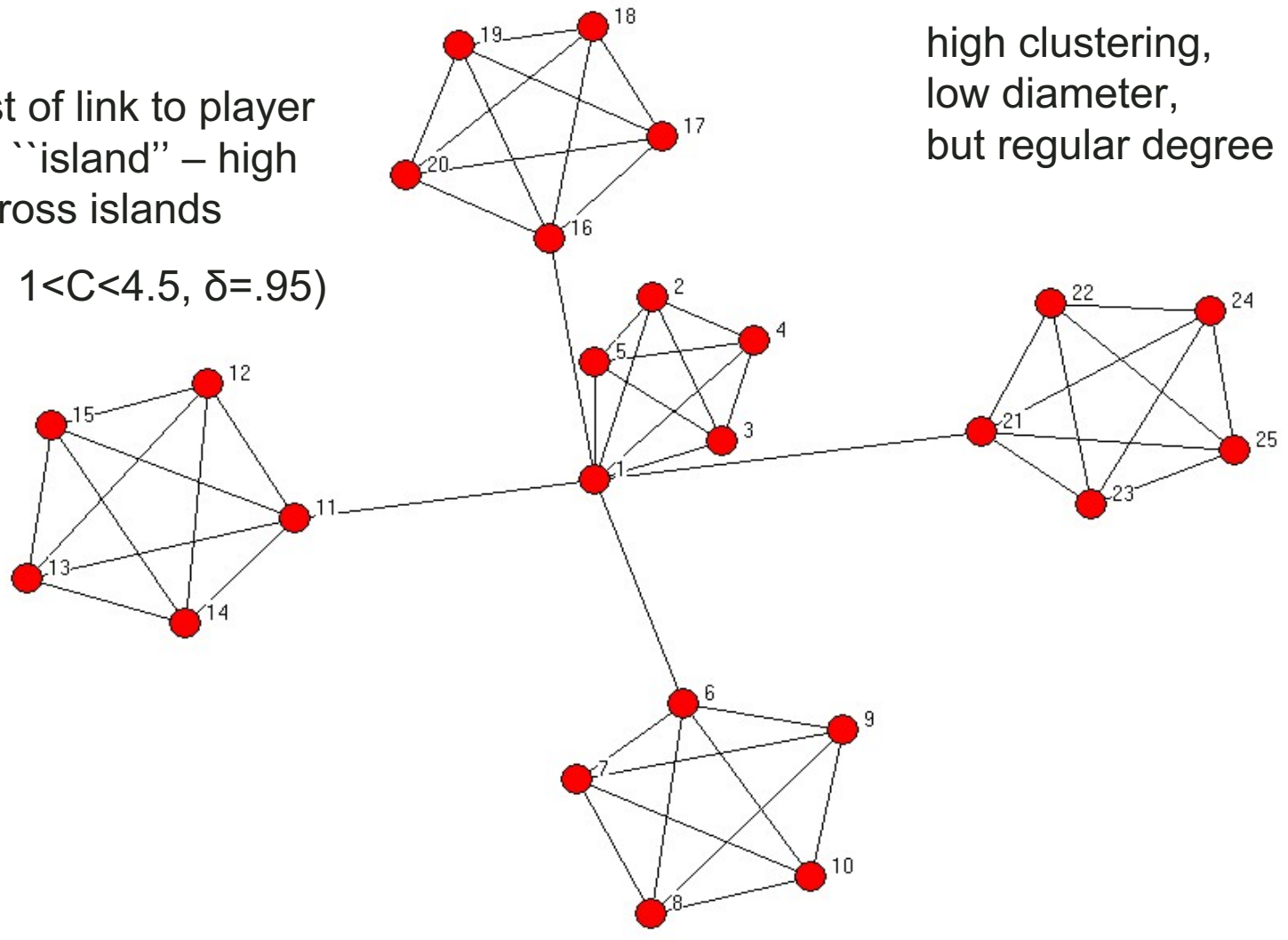
### Results:

- High clustering within islands, few links across
- small distances



low cost of link to player  
on own "island" – high  
cost across islands

( $c < .04$ ,  $1 < C < 4.5$ ,  $\delta = .95$ )



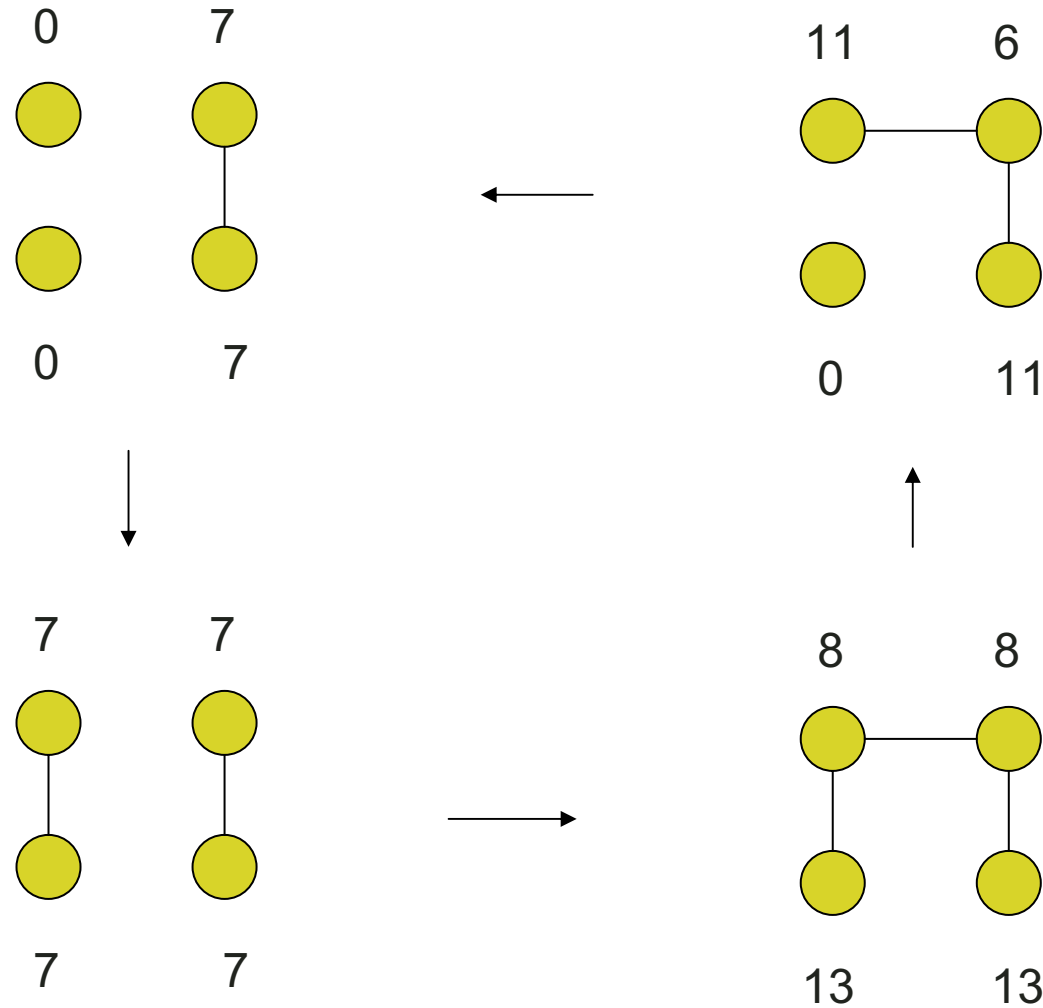
high clustering,  
low diameter,  
but regular degree

# Rich literature on strategic formation



- loosen anonymity (Dutta-Mutuswami (1997))
- directed networks (Bala-Goyal (2000), Dutta-Jackson (2000),...)
- bargaining when forming links (Currarini-Morelli(2000), Slikker-van den Nouweland (2000), Mutuswami-Winter(2002), Bloch-Jackson (2004))
- dynamic models (Aumann-Myerson (1988), Watts (2001), Jackson-Watts (2002ab), Goyal-Vega-Redondo (2004), Feri (2004), Lopez-Pintado (2004),...)
- farsighted models (Page-Wooders-Kamat (2003), Dutta-Ghosal-Ray (2003), Deroian (2003),...)
- allocating value (Myerson (1977), Meessen (1988), Borm-Owen-Tijs (1992), van den Nouweland (1993), Qin (1996), Jackson-Wolinsky (1996), Slikker (2000), Jackson (2005)...)
- modeling stability (Dutta-Mutuswami (1997), Jackson-van den Nouweland (2000), Gilles-Sarangi (2003ab), Calvo-Armengol and Ilicic (2004),...)
- experiments (Callander-Plott (2001), Corbae-Duffy (2001), Pantz-Zeigelmeyer (2003), Charness-Corominas-Bosch-Frechette (2001), Falk-Kosfeld (2003), ...)

# Nonexistence of Pairwise Stable



(Lending Model  
Jackson-Watts (2002))

# Advantages of an economic approach



- Payoffs allow for a welfare analysis
  - Identify tradeoffs – incentives versus efficiency
- Tie the nature of externalities to network formation...
- Put network structures in context - outcomes of network interaction
- Account for (and *explain*) some observables

# What's missing from Game theoretic formation models?



- Stark network structures emerge
  - need to mix with random models
- over-emphasize choice versus chance determinants for some *large* applications?

# Models of Networks in Context



- crime networks (Glaeser-Sacerdote-Scheinkman (1996), Ballester, Calvo, Zenou (2003),...)
- markets (Kirman (1997), Tesfatsion (1997), Weisbach-Kirman-Herreiner (2000), Kranton-Minehart (2002), Corominas-Bosch (2005), Wang-Watts (2002), Galeotti (2005),Kakade et al (2005)...) )
- labor networks (Boorman (1975), Montgomery (1991, 1994), Calvo (2000), Arrow-Borzekowski (2002), Calvo-Jackson (2004ab,2005), Cahuc-Fontaine (2004),...)
- insurance (Fafchamps-Lund (2000), DeWeerd (2002), Bloch-Genicot-Ray (2004),...)
- IO (Bloch (2001), Goyal-Moraga (2001), Goyal-Joshi (2001), Belleflamme-Bloch (2002),Billard-Bravard (2002), ...)
- international trade (Casella-Rauch (2001), Furusawa-Konishi (2003),...)
- public goods (Bramouille-Kranton (2004), Galeotti and Vega (2005),...)
- airlines (Starr-Stinchcombe (1992), Hendricks-Piccione-Tan (1995))
- network externalities in goods (Katz-Shapiro (1985), Economides (1989, 1991) , Sharkey (1991)...) )
- organization structure (Radner (), Radner-van Zandt (), Demange (2004)...) )
- learning (Bala-Goyal (1998), Morris (2000), DeMarzo-Vayanos-Zweibel (2003), Gale-Kariv (2003), Choi-Gale-Kariv (2004),...)



# Whither now?



- Bridging random/mechanical – economic/strategic
- Networks in Applications - predictions of behavior as dependent on network structure
  - Labor, mobility, voting, trade, collaboration, crime, public goods, ...
- Empirical/Experimental
  - many case studies lack key economic variables, tie networks to outcomes, utilities
  - enrich modeling of social interactions from a structural perspective
- Furthering game theoretic modeling, equilibrium, dynamics,...
- Foundations and Tools– centrality, power, allocation rules, community structures, ...

# Setting stable=efficient depends on who bears link cost



Buyer-Seller Networks: Kranton-Minehart (2002):

- Sellers each with one identical object
- Buyers each desire one object, private valuation
- links only costly to buyers
- sellers hold simultaneous ascending auctions

# Example: values iid $U[0,1]$ , 1 seller



	Each buyer's expected utility	Seller's expected utility	Total social value
n buyers	$1/[n(n+1)]$	$(n-1)/(n+1)$	$n/(n+1)$
n+1 buyers	$1/[(n+1)(n+2)]$	$n/(n+2)$	$(n+1)/(n+2)$
change	$-2/[n(n+1)(n+2)]$	$2/[(n+1)(n+2)]$	$1/[(n+1)(n+2)]$

