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Properties of Quantum Entropy and Related Convex Trace Functions

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Properties of Quantum Entropy and Related Convex Trace Functions

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Lecture Plan

1. Introduction and Background on Quantum Information
2. Properties of Entropy and Relative Entropy
3. Simple Proof of Joint Convexity of Relative Entropy
4. a) Noise in Quantum Information
   b) More Inequalities and conjectures

I. Introduction and Background on Quantum Information

1. Preliminaries
2. Dirac bra and ket notation
3. Overview of quantum information
4. Some quantum mechanics basics
5. Quantum entropy
6. Tensor products and entanglement
7. Aside on SVD and "Schmidt" decompostion
8. Quantum relative entropy

Hilbert space set-up

Full quantum theory associates a particle, e.g., electron with

$\mathcal{H} = L_2(\mathbf{R}_3) \otimes \mathbf{C}_d \qquad d = 2s+1,\ s$ denotes "spin"

Quantum info suppress "spatial" part in $L_2(\mathbf{R}_3)$ and focus on spin,

typically spin $\frac{1}{2}$ or $\mathbf{C}_2$. Many particles or qubits use $\mathbf{C}_2^{\otimes n}$

can also consider $d > 2$ and $\mathbf{C}_d \otimes \mathbf{C}_{d'} \otimes \ldots$ etc.

Will work with $\mathcal{H} = \mathbf{C}_d$ or tensor products of these

and $\mathcal{B}(\mathcal{H}) = M_d =$ space of $d \times d$ matrices

$M_d =$ also a Hilbert space with inner product $\langle A, B \rangle = \operatorname{Tr} A^* B$

will also consider linear maps $\Phi : M_d \mapsto M_{d'}$

Most results extend to ∞ dim. $\mathcal{H}$ in suitable way

and can even replace M_d by a von Neumann algebra

Basic notation

A, B, C self-adjoint matrices

$A > B$ means $A - B$ positive definite $\langle v, (A-B)v\rangle > 0 \quad \forall\, v \neq 0$

$A \geq B$ or $A \gneq B$ means $A - B$ positive semi-definite, but $A \neq B$.

$A \geqq B$ means $A - B$ positive semi-definite (with $A = B$ allowed)

difference rarely significant $\begin{cases} f \text{ obviously diverges when } B = 0 \\ \lim f(B + \epsilon I) \text{ well defined as } \epsilon \to 0 \end{cases}$

operator (or matrix) inequality has one of above forms

trace inequality has form $\operatorname{Tr} AC > \operatorname{Tr} BC$ or $\operatorname{Tr} AC \geq \operatorname{Tr} BC$

with usual meanings for $>, \geq$ when $\operatorname{Tr} X \in \mathbf{R}$ (usually $[0, \infty)$

study trace inequalities related to maps $M_d \mapsto R$ that are convex

Functions of operators

For $A = UDU^*$, define $f(A) = U f(D) U^*$

$$A = U \begin{pmatrix} \lambda_1 & 0 & \dots & 0 \\ 0 & \lambda_2 & \dots & 0 \\ \vdots & & \ddots & \vdots \\ 0 & \dots & 0 & \lambda_m \end{pmatrix} U^\dagger \qquad f(A) = U \begin{pmatrix} f(\lambda_1) & 0 & \dots & 0 \\ 0 & f(\lambda_2) & \dots & 0 \\ \vdots & & \ddots & \vdots \\ 0 & \dots & 0 & f(\lambda_m) \end{pmatrix} U^\dagger$$

equiv. to any reasonable def using power series, integral rep., etc.

also applies to operators, e.g., L_Q acting on M_d space of matrices

For example, for $L_A(X) = AX$ find $L_{\log A}(X) = \log L_A(X)$

Examples of trace functions

$R, Q > 0$ but often extend to $R, Q \geq 0$

applications often use $\operatorname{Tr} R = \operatorname{Tr} Q = 1$ but not essential

$S(R) = -\operatorname{Tr} R \log R$ quantum entropy (concave)

$H(R, Q) = \operatorname{Tr} R(\log R - \log Q)$ relative entropy (jointly convex)

K fixed $\operatorname{Tr} K^* R^p K Q^{1-p}$ jointly $\begin{cases} \text{concave} & 0 < p < 1 \quad \text{(Lieb)} \\ \text{linear} & p = 0, 1 \\ \text{convex} & 1 < p < 2 \quad \text{(Ando)} \end{cases}$

(i) math – lead to many interesting matrix inequalities

(ii) original interest in quantum statistical mechanics

now important in quantum information theory

Dirac notation: $\mathbf{u}.\mathbf{v}$ in finite-dim vector space $\mathbf{C}_d$

Inner product $\langle v, u\rangle = \mathbf{v}^*\mathbf{u} = \begin{pmatrix} & \bar{v} & \end{pmatrix} \begin{pmatrix} \\ u \\ \\ \end{pmatrix}$

Reverse order $|u\rangle\langle v| = \mathbf{u}\,\mathbf{v}^* = \begin{pmatrix} \\ u \\ \\ \end{pmatrix} \begin{pmatrix} & \bar{v} & \end{pmatrix}$

get $n \times n$ matrix of map $w \mapsto \langle v, w\rangle u$

$P_u = \dfrac{1}{\|u\|^2} uu^* = \dfrac{|u\rangle\langle u|}{\|u\|^2}$ projection onto 1-dim subspace span$\{u\}$.

"ket" $|u\rangle \leftrightarrow \mathbf{u} \leftrightarrow$ col vec dual "bra" $\langle u| \leftrightarrow \mathbf{u}^* \leftrightarrow$ row vec

- Can be justified by Riesz rep. theorem.
- Put complex conj. on "left" i.e., $\langle v, u\rangle$ linear in u; anti-lin in v.
- Use any convenient label, e.g. $|\lambda_k\rangle$ or $|k\rangle$ for eigenfctn v_k of λ_k

conventions

in view of above $\langle u, v\rangle = \mathbf{u}^*\mathbf{v}$ anti-linear in u and linear in v

$$a\langle u, v\rangle = \langle u, av\rangle = \langle \bar{a}u, v\rangle$$

math	physics	MBR	this lecture	
A^*	$A^\dagger$	$A^\dagger$	A^*	adjoint
$\bar{a}$	a^*	$\bar{a}$	$\bar{a}$	complex conjugate
		a^*		never – too confusing
		$\widehat{\Phi}$	$\widehat{\Phi}$	adjoint wrt H-S inner prod

$\Phi : M_d \mapsto M_{d'}$ linear map

$$\operatorname{Tr} A^*\Phi(B) = \langle A, \Phi(B)\rangle = \langle \widehat{\Phi}(A), B\rangle = \operatorname{Tr}[\widehat{\Phi}(A)]^* B$$

Quantum vs classical information

Classical – "bit" takes values in $\mathbf{Z}_2 = \{0, 1\}$, e..g, "on" or "off"

encode info in strings of 0 & 1, elements of $\mathbf{Z}_2^{\otimes n} = \mathbf{Z}_2 \otimes \mathbf{Z}_2 \ldots \otimes \mathbf{Z}_2$

Quantum – "qubit" takes values in $\mathbf{C}_2$ (up to norm. and phase)

$0 \sim \begin{pmatrix} 1 \\ 0 \end{pmatrix} \equiv |0\rangle_z$ $\uparrow$ spin "up" or $\uparrow$ vertical polar

$1 \sim \begin{pmatrix} 0 \\ 1 \end{pmatrix} \equiv |1\rangle_z$ $\downarrow$ spin "down" or $\rightarrow$ horiz polar

Isomorphism between $\mathbf{Z}_2^{\otimes n}$ and O.N. prod basis for $\mathbf{C}_2^{\otimes n}$

"Computational basis", e.g., $|0110\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix} \otimes \begin{pmatrix} 1 \\ 0 \end{pmatrix}$

Embed class in quant — but can do much more

$0 \sim \begin{pmatrix} 1 \\ 0 \end{pmatrix} \equiv |0\rangle_z$ $\uparrow$ spin "up" or $\uparrow$ vertical polar

$1 \sim \begin{pmatrix} 0 \\ 1 \end{pmatrix} \equiv |1\rangle_z$ $\downarrow$ spin "down" or $\rightarrow$ horiz polar

Now consider $\begin{pmatrix} 1 \\ \pm 1 \end{pmatrix}$ "spin" right $\rightarrow$ or left $\leftarrow$ (in x-direction)

$$\frac{1}{\sqrt{2}}\begin{pmatrix} 1 \\ \pm 1 \end{pmatrix} = \frac{1}{\sqrt{2}}\begin{pmatrix} 1 \\ 0 \end{pmatrix} + \frac{1}{\sqrt{2}}\begin{pmatrix} 0 \\ 1 \end{pmatrix} = \frac{1}{\sqrt{2}}(|0\rangle_z + |1\rangle_z)$$

Then measure spin in z-direction. Get either $\uparrow$ or $\downarrow$ (i.e., 0 or 1)

each with probability $\frac{1}{2}$. But not classical prob, i.e.,

Not a classical mixture but a superposition of vectors.

In some sense $2^{-n/2}\begin{pmatrix} 1 \\ 1 \end{pmatrix}^{\otimes n} = \frac{1}{\sqrt{2}}\begin{pmatrix} 1 \\ 1 \end{pmatrix} \otimes \frac{1}{\sqrt{2}}\begin{pmatrix} 1 \\ 1 \end{pmatrix} \otimes \ldots \otimes \frac{1}{\sqrt{2}}\begin{pmatrix} 1 \\ 1 \end{pmatrix}$

contains all 2^n strings of 00110101...., each with prob 2^{-n}

Can one encode more than $\{0, 1\}$ in qubit ??

4 states $\begin{pmatrix} 1 \\ 0 \end{pmatrix} \begin{smallmatrix} \uparrow \\ \uparrow \end{smallmatrix}$ $\frac{1}{\sqrt{2}}\begin{pmatrix} 1 \\ 1 \end{pmatrix} \begin{smallmatrix} \rightarrow \\ \nearrow \end{smallmatrix}$ $\begin{pmatrix} 0 \\ 1 \end{pmatrix} \begin{smallmatrix} \downarrow \\ \rightarrow \end{smallmatrix}$ $\frac{1}{\sqrt{2}}\begin{pmatrix} 1 \\ -1 \end{pmatrix} \begin{smallmatrix} \leftarrow \\ \searrow \end{smallmatrix}$??

Most general qubit $|v\rangle = a|0\rangle + b|1\rangle = \begin{pmatrix} a \\ b \end{pmatrix} = \begin{pmatrix} \sin\theta \\ e^{i\varphi}\cos\theta \end{pmatrix}$

use to encode $[0, \pi]$ or $[-1, 1]$ or more ??

NO — will see can only reliably distinguish orthog states

Thm: (Holevo bound) Accessible Information or max info/qubit one can extract ≤ 1. Will give formal math thm and proof

non-orthog encoding can't reliably distinguish BUT advantages!

- Noisy communication (use quant particles to sed class info)

  non-orthog input may yield more distinguishable outputs

- Quant cryptography — sacrifice info to detect eavesdropper

Quantum basics and von Neumann measurement

Fund Postulate of Q.M.: Observable represented by self-adj op A

spectral decomp $A = \sum_k a_k E_k = \sum_k a_k |\alpha_k\rangle\langle\alpha_k|$

Measurement of A with system in some state ψ.

(i) get some e-value (only possibility)

(ii) leave system in e-state α_k

(iii) probability is $|\langle\alpha_k, \psi\rangle|^2 = \operatorname{Tr} E_k |\psi\rangle\langle\psi|$

Write $|\psi\rangle = \sum_k c_k |\alpha_k\rangle$ as a superposition of e-states, $c_k = \langle\alpha_k, \psi\rangle$

Coefficients c_k in superpos. give probs $|c_k|^2$ not classical

Average result of meas in state $|\psi\rangle$ is $\langle\psi, A\psi\rangle = \operatorname{Tr} A|\psi\rangle\langle\psi|$

set $\{E_k\}$ orthog projections $E_j E_k = E_k \delta_{jk}$ with $\sum_k E_k = I$ called

von Neumann measurement or projection valued measure (PVM)

corresponds to "yes-no" experiment (e.g., polarization filter)

Density matrices: mixed vs. pure states

pure state often rep by vector $|\psi\rangle \in \mathcal{H}$ up to phase, with $\|\psi\| = 1$.

better to use rank-1 projection $|\psi\rangle\langle\psi|$ (with $\|\psi\| = 1$).

mixed state $\rho = \sum_k p_k |\phi_k\rangle\langle\phi_k|$ is convex comb of pure states

$p_k > 0, \sum_k p_k = 1$ and $\|\phi_k\| = 1$ but ϕ_k not nec orthog

ρ called density matrix (D.M.) or density operator

$\rho \in M_d$ is D.M. if and only if $\rho \geq 0$ and $\operatorname{Tr} \rho = 1$.

Interp: a) ensemble in quantum statistical mechanics

b) know only part (subsystem) $\rho = \rho_A = \operatorname{Tr}_B \rho_{AB}$

c) $A \mapsto \operatorname{Tr} A \rho$ positive linear functional on M_d.

two kinds of probability in Q.M. (i) p_k traditonal prob interp,

but (ii) $|\phi_k\rangle$ can be superposiition with different interp.

von Neumann's quantum entropy

von Neumann (1927) defined mixed quantum state and its entropy

$$S(\rho) \equiv -\operatorname{Tr} \rho \log \rho = -\sum_k \lambda_k \log \lambda_k$$

where ρ spectral decomp $\rho = \sum_k \lambda_k |\chi_k\rangle\langle\chi_k|$ so λ_k e-vals

Density matrix $\rho > 0$ and $\operatorname{Tr} \rho = 1 \quad \Rightarrow \quad S(\rho) \geq 0$

also find $\rho = |\psi\rangle\langle\psi|$ pure $\Leftrightarrow \rho^2 = \rho \Leftrightarrow S(\rho) = 0$

But $S(P)$ well-defined for and concave any pos semi-def ops

will give three (3) proofs

$S(\rho) \geq 0$ is result of normalization and/or phys interp

Shannon (1948): classical info with entropy equiv. to diag matrix

Next Time: more, including subadditivity properties

Back to Measurement: Role of non-commutativity

Now consider two non-commuting observables

$$A = \sum_j a_j |\alpha_j\rangle\langle\alpha_j| = \sum_j a_j E_j, \qquad B = \sum_k b_k |\beta_k\rangle\langle\beta_k| = \sum_k b_k F_k$$

start in $|\psi\rangle$ measure A, then B ends in e-state $|\beta_k\rangle$ of B

start in $|\psi\rangle$ measure B, then A ends in e-state $|\alpha_j\rangle$ of A

in mixed $\rho = \sum_k p_k |\phi_k\rangle\langle\phi_k|$ average result of measuring A

is $\sum_k p_k \langle\phi_k, A\phi_k\rangle = \operatorname{Tr} \rho A$

Define map $\Omega_{\mathcal{M}}$ describes result of PVM or vN measurement

$$\Omega_{\mathcal{M}} : \rho \mapsto \sum_j E_j \rho E_j = \sum_j |\alpha_j\rangle\langle\alpha_j, \rho\, \alpha_j\rangle\langle\alpha_j| = \sum_j |\alpha_j\rangle\langle\alpha_j| \operatorname{Tr} \rho E_j$$

Measure B, then A ends with $F_k \mapsto \Omega_{\mathcal{M}}(F_k) = \sum_j E_j F_k E_j$

Quantum measurement: POVM

$$\sum_{jk} E_j F_k E_j = \sum_j E_j I E_j = I$$

$\{E_j F_k E_j\}$ example of POVM positive operator valued measurement

Def: (Davies and Lewis) POVM $\mathcal{M} = \{G_b\}$ $\quad G_b > 0, \ \sum_b G_b = I$

Result of POVM depends on order in which G_b performed

QC map for von Neumann measurement

$$\Omega_{\mathcal{M}} : \gamma \mapsto \sum_j (\mathrm{Tr}\,\gamma E_j)|\alpha_j\rangle\langle\alpha_j| \qquad E_j = |\alpha_j\rangle\langle\alpha_j| \ \text{ O.N.}$$

QC map for POVM using instrument with "pointer" $|f_b\rangle$

$$\Omega_{\mathcal{M}} : \gamma \mapsto \sum_j (\mathrm{Tr}\,\gamma G_b)|\phi_b\rangle\langle\phi_b| \otimes |f_b\rangle\langle f_b| \text{ where } |f_b\rangle \ \text{ O.N.}$$

Tensor products and entanglement

Quant Info typical $\mathcal{H} = \mathcal{H}_1 \otimes \mathcal{H}_2 \otimes \ldots \otimes \mathcal{H}_n$ or $\mathcal{H} = \mathcal{H}_A \otimes \mathcal{H}_B \otimes \ldots$

- Quantum Computer $\mathcal{H} = \mathbf{C}_2^{\otimes n}$ for n qubits
- Quantum Communication $\mathcal{H}_A \otimes \mathcal{H}_B \otimes \mathcal{H}_E$

A = sender "Alice" B = receiver "Bob"

E = "Eve" (Eavesdropper – sexist) , but also

E = Environment (can be Evil or Friendly)

Partial trace $\mathrm{Tr}_B : \mathcal{B}(\mathcal{H}_A \otimes \mathcal{H}_B) \mapsto \mathcal{B}(\mathcal{H}_A)$ or $M_{d_A} \otimes M_{d_B} \mapsto M_{d_A}$

$\mathrm{Tr}_B\, A \otimes B = A(\mathrm{Tr}\, B)$ extend by linearity

Formal partial inner product $|\phi_k\rangle$ O.N. basis for $\mathcal{H}_B$

$\mathrm{Tr}_B X_{AB} = \sum_k \langle\phi_k, X_{AB}\,\phi_k\rangle_B$ means $\quad \forall\ \chi, \psi \in \mathcal{H}_A$

$X_A = \mathrm{Tr}_B X_{AB}$ iff $\langle\chi, X_A\,\psi\rangle = \sum_k \langle\chi\otimes\phi_k, X_{AB}\,\psi\otimes\phi_k\rangle_B$

Aside: entanglement measure

pure state $|\psi\rangle \in \mathcal{H} = \mathcal{H}_A \otimes \mathcal{H}_B$ called "entangled" if it is

not a tensor prod, i.e., can not be written $|\psi\rangle \neq \phi_A \otimes \phi_B$

Example max entang Bell states on $\mathbf{C}_2 \otimes \mathbf{C}_2$

$\frac{1}{\sqrt{2}}(|00\rangle \pm |11\rangle) \qquad \frac{1}{\sqrt{2}}(|01\rangle \pm |10\rangle)$

have non-classical correlations long regarded as mysterious

modern quant info attitude: accept as way world is and ask

What nifty new things can we do with entanglement?

Measure entanglement of pure state

$E(\psi) = S\big(\mathrm{Tr}_B\, |\psi\rangle\langle\psi|\big) = S\big(\mathrm{Tr}_A\, |\psi\rangle\langle\psi|\big)$ essent unique

mixed ρ_{AB} separable if convex comb of pure product states

many (industry) of inequiv. entang measures for mixed states

Aside: SVD and "Schmidt" decomposition

Singular Value Decomposition: Recall $B^*B = \sum_k \mu_k^2 |b_k\rangle\langle b_k| \equiv |B|^2$

Then $B = U|B| = \sum_k \mu_k |a_k\rangle\langle b_k| \qquad |a_k\rangle = U|b_k >$

U partial isometry – restriction to $(\ker B)^{\perp}$ unique unitary

Isomorphism $\mathcal{B}(\mathcal{H}) \simeq \mathcal{H} \otimes \mathcal{H} \qquad |v\rangle\langle w| \leftrightarrow |v \otimes w\rangle$

apply SVD + iso to $|\psi\rangle \in \mathcal{H} \otimes \mathcal{H}$ $|\psi\rangle = \sum_k \mu_k |\alpha_k \otimes \beta_k\rangle$

pure $\rho_{AB} = |\psi\rangle\langle\psi| \Rightarrow$ reduced density matrices $\rho_A \equiv \mathrm{Tr}_B\, \rho_{AB}$ etc.

$\rho_A = \sum_k |\mu_k|^2 |\alpha_k\rangle\langle\alpha_k| \quad \rho_B = \sum_k |\mu_k|^2 |\beta_k\rangle\langle\beta_k|$

Cor 1: $\rho_{AB} = |\psi\rangle\langle\psi|$ pure $\Rightarrow \rho_A, \rho_B$ have same non-zero e-vals

Cor 2: $\rho_{AB} = |\psi\rangle\langle\psi|$ pure $S(\rho_A) = S(\rho_B)$

Can reverse to get "purification" start with $\rho = \sum_k \lambda_k |\phi_k\rangle\langle\phi_k|$

Define $|\psi\rangle = \sum_k \sqrt{\lambda_k}\, |\phi_k \otimes \phi_k\rangle \in \mathcal{H} \otimes \mathcal{H}$ $\quad \mathrm{Tr}_B |\psi\rangle\langle\psi| = \rho$

Quant Info view: mystical result of Schmidt about tensor products

SVD for matrices started 1870's (Horn and Johnson, Chap. 3)

Schmidt(1907) equiv. result interp $K(x,y)$ as kernal of operator

$g(y) \mapsto f(x) = \int K(x,y)g(y)dy$

Rediscovered by quantum chemists called Carleson-Keller (1961)

John Coleman (1963) pointed out due to Schmidt

OK interp for $\psi(x_1 \ldots x_m, y_1 \ldots y_n) \in L_2(\mathbf{R}^{m+n})$ wave function

wrong-headed to look for extension to higher order tensor products

More info: See Appendix A of King & Ruskai

IEEE Trans. Info. Theory, 47, 192209 (2001) quant-ph/9911079

Aside: G Homogenous of degree one

Consequences of $G(\lambda A) = \lambda G(A)$

- $G(A)$ convex $\Leftrightarrow$ $G(A+B) \le G(A) + G(B)$

convex $\Rightarrow$ "subadditive" [in sense $G(A+B) \le G(A) + G(B)$]

$$\tfrac{1}{2}G(A+B) = G(\tfrac{1}{2}A + \tfrac{1}{2}B) \le \tfrac{1}{2}G(A) + \tfrac{1}{2}G(B)$$

"subadditive" $\Rightarrow$ convex

$$\begin{aligned} G\big[xA + (1-x)B\big] &\le G(xA) + G[(1-x)B] \\ &= xG(A) + (1-x)G(B) \end{aligned}$$

- $G(x)$ convex $\Rightarrow$ $\lim_{x \to 0} \frac{1}{x}\big[G(A+xB) - G(A)\big] \le G(B)$

$$\lim_{x\to 0} \frac{G(A+xB) - G(A)}{x} \le \lim_{x \to 0} \frac{G(A) + xG(B) - G(A)}{x}$$

Relative Entropy

Def: $H(\rho,\gamma) = \mathrm{Tr}\, \rho\big(\log\rho - \log\gamma\big)$

ρ, γ pair of density matrices, but can be any pos semi-def ops.

Use Greek $\rho, \gamma, \ldots$ for density matrices,

Use Roman R, Q for arb pos semi-def matrices

Joint Convexity $\quad H\Big(\sum_j R_j, \sum_j Q_j\Big) \le \sum_j H(R_j, Q_j)$

Cor: Monotone under Partial Trace (MPT)

$$H(R_1, Q_2) \le H(R_{12}, Q_{12})$$

$$R_{12}, Q_{12} \in \mathcal{B}(\mathcal{H}_1 \otimes \mathcal{H}_2) \qquad R_1 = \mathrm{Tr}_2\, R_{12} \text{ etc.}$$

II. Properties of Entropy and Relative Entropy

1. Fundamental Properties of Entropy
   - 1.1 Concavity and subadditivity
   - 1.2 Minor properties
   - 1.3 Strong subadditiivty (SSA)
   - 1.4 First proof of concavity of $S(\rho)$
   - 1.5 Triangle inequality
2. Klein's inequality and second proof of concavity of $S(\rho)$
3. Fundamental Properties of Relative Entropy
4. Aside on Information Theory Expressions
5. Aside on Monotone and Convex Operator Functions
6. Third proof of concavity of $S(\rho)$

Fundamental Properties of Quantum Entropy

Def: $S(R) = -\text{Tr}\, R \log R$ for $R \geq 0$ pos semi-def ($0 \log 0 \equiv 0$)

- Concave: $\quad x\, S(R_1) + (1-x) S(R_2) \leqq S(xR_1 + (1-x)R_2)$

but $\quad \leq x\, S(R_1) + (1-x)S(R_2) + x \log x + (1-x)\log(1-x)$

- Subadditive: $\quad S(R_{AB}) \leqq S(R_A) + S(R_B)$

  with $= \Leftrightarrow \; R_{AB} = R_A \otimes R_B$

- Strongly Subadditive $\quad S(R_B) + S(R_{ABC}) \leqq S(R_{AB}) + S(R_{BC})$

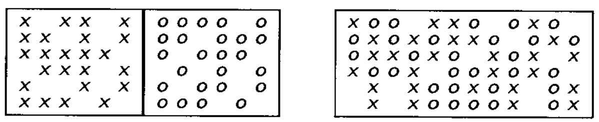
All indep of norm. if consistent, e.g., $\text{Tr}\, R_B = \text{Tr}\, R_{AB} = \text{Tr}\, R_{ABC}$

Also have (a) $\quad S(\mu R) = \mu S(R) - \mu \log \mu \text{Tr}\, R$

(b) $R \in M_d$ and $\text{Tr}\, R = 1 \;\Rightarrow\; 0 \leq S(R) \leq \log d$

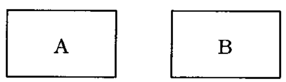
Concave: $\quad x\, S(\rho_1) + (1-x)S(\rho_2) \leq S(x\rho_1 + (1-x)\rho_2)$

refers to mixture



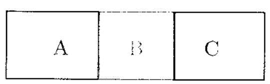
Subadditive: $\quad S(\rho_{AB}) \leq S(\rho_A) + S(\rho_B)$

refers to regions or subsystems



SSA: $\quad S(\rho_B) + S(\rho_{ABC}) \leq S(\rho_{AB}) + S(\rho_{BC})$

overlapping regions



Subadditive $\Rightarrow$ Concave

Proof # 1 of $S(\rho)$ concave **once subadd shown** (easy)

$$\rho_{AB} = \begin{pmatrix} x\rho_1 & 0 \\ 0 & (1-x)\rho_2 \end{pmatrix}$$

$$\rho_A = x\rho_1 + (1-x)\rho_2, \qquad \rho_B = \left(\begin{smallmatrix} x & 0 \\ 0 & 1-x \end{smallmatrix}\right)$$

$$\begin{aligned} S(\rho_{AB}) &= S(x\rho_1) + S((1-x)\rho_2) \\ &= xS(\rho_1) - x\log x + (1-x)S(\rho_2) - (1-x)\log(1-x) \end{aligned}$$

$$\begin{aligned} S(\rho_{AB}) &\leq S(\rho_A) + S(\rho_B) \\ &= S(x\rho_1 + (1-x)\rho_2) - x\log x - (1-x)\log(1-x) \end{aligned}$$

$$\Rightarrow \quad xS(\rho_1) + (1-x)S(\rho_2) \leq S(x\rho_1 + (1-x)\rho_2)$$

Triangle inequality for $S(\rho)$

$$S(\rho_{AB}) \leqq S(\rho_A) + S(\rho_B)$$

Given ρ_{AB} can find $\mathcal{H}_C$ and $|\psi_{ABC}\rangle \in \mathcal{H}_A \otimes \mathcal{H}_B \otimes \mathcal{H}_C$

s.t. $\text{Tr}_C\, \rho_{ABC} = \text{Tr}_C\, |\psi_{ABC}\rangle\langle\psi_{ABC}| = \rho_{AB}$ purification

$\rho_{AB} = \sum_k \lambda_k |\phi_k\rangle\langle\phi_k| \qquad |\psi_{ABC}\rangle = \sum_k \sqrt{\lambda_k}\, |\phi_k \otimes e_k\rangle$

$(\mathcal{H}_A \otimes \mathcal{H}_B) \otimes \mathcal{H}_C \; S(\rho_{AB}) = S(\rho_C) \quad \mathcal{H}_A(\otimes \mathcal{H}_B \otimes \mathcal{H}_C)\; S(\rho_A) = S(\rho_{BC})$

subs above $\quad S(\rho_C) \leq S(\rho_{BC}) + S(\rho_B)$

$\Rightarrow \quad S(\rho_C) - S(\rho_B) \leq S(\rho_{BC})$

Reverse B $\leftrightarrow$ C and combine with subadd to get

$$|S(\rho_A) - S(\rho_B)| \leq S(\rho_{AB}) \leqq S(\rho_A) + S(\rho_B)$$

Klein's inequality:

$$\mathrm{Tr}\, A \log A - \mathrm{Tr}\, A \log B \geq \mathrm{Tr}\,(A - B) \quad \text{with} \; = \; \text{iff } A = B$$

g convex means diff quotients increase

$\Rightarrow \quad \frac{g(b)-g(a)}{b-a} \leq g'(b)$ for $a < b$

$\Rightarrow \quad g(b) - g(a) \leq (b-a)g'(b)$ for all a, b

$\mathrm{Tr}\,[g'(B)(B-A) - g(B) + g(A)]$ $\qquad |\beta_k\rangle$norm e-vec of B

$$= \sum_k \Big[g'(b_k)(b_k - \langle \beta_k, A\beta_k\rangle) - g(b_k) + \langle \beta_k, g(A)\beta_k\rangle \Big]$$

Jensen $g(\langle \beta_k, A\beta_k\rangle) \leq \langle \beta_k, g(A)\beta_k\rangle$

$$\geq \sum_k \Big[g'(b_k)(b_k - \langle \beta_k, A\beta_k\rangle) - g(b_k) + g(\langle \beta_k, A\beta_k\rangle) \Big] \geq 0$$

For $g(x) = x \log x, \quad g'(x) = 1 + \log x$

$$\mathrm{Tr}\,[(B-A)(I + \log B) - B \log B + A \log A] \geq 0$$

Entropy concave : Proof # 2

$A = R_1$ and $B = R = xR_1 + (1-x)R_2$ in Klein

$$\mathrm{Tr}\, A \log A - \mathrm{Tr}\, A \log B \geq \mathrm{Tr}\,(A - B)$$

Get $\quad \mathrm{Tr}\, R_1 \log R_1 - R_1 \log R \geq \mathrm{Tr}\,(R_1 - R)$

Repeat for R_2 and Mult by x and $1 - x$

$$\begin{aligned} x\,\mathrm{Tr}\, R_1 \log R_1 - \; x R_1 \log R &\geq \mathrm{Tr}\,[\, x R_1 - R] \\ (1-x)\mathrm{Tr}\, R_2 \log R_2 - (1-x) R_2 \log R &\geq \mathrm{Tr}\,[(1-x)R_2 - R] \end{aligned}$$

add to get

$$\begin{aligned} -xS(R_1) - (1-x)S(R_2) - [xR_1 + (1-x)R_2] \log R &\geq \mathrm{Tr}\,(R - R) \\ -xS(R_1) - (1-x)S(R_2) + S(R) &\geq 0 \end{aligned}$$

Relative entropy

Def: $H(R, Q) \; = \; \mathrm{Tr}\, R \log R - \mathrm{Tr}\, R \log Q$

Klein's ineq: $H(R, Q) \geq \mathrm{Tr}\,(R - Q) \qquad \geq 0$ if $\mathrm{Tr}\, R = \mathrm{Tr}\, Q$

assume $R, Q > 0$ strictly pos — well-def if $\ker(Q) \subset \ker(R)$

Can obtain entropy from rel ent. $\quad H(R, \frac{1}{d} I) = -S(R) + \log d$

Or simply $H(R, I) = -S(R)$

homogenous of degree one $\qquad H(\lambda R, \lambda Q) = \lambda H(R, Q)$

Recall $\Rightarrow$ convexity equiv to simple form $F(A + B) \leq F(A) + F(B)$

Fundamental Properties of Relative Entropy

Joint Convexity $\quad H(R_1 + R_2, Q_1 + Q_2) \leq H(R_1, Q_1) + H(R_2, Q_2)$

Cor: (MPT) $\qquad H(R_A, Q_A) \leq H(R_{AB}, Q_{AB})$

Monotone under Partial Trace

Ibinson-Winter: that's all folks!

Special Case of MPT $\quad R_{AB} \to \rho_{ABC}, \; Q_{AB} \to I_A \otimes \rho_{BC}$

gives strong subadditivity $\quad H(\rho_{AB}, \rho_B) \leq H(\rho_{ABC}, \rho_{BC})$

SSA $\qquad -S(\rho_{AB}) + S(\rho_B) \leq -S(\rho_{ABC}) + S(\rho_{BC})$

Will prove JC and then show $\Rightarrow$ MPT

Aside: Information Theory Expressions

Mutual Information: (always positive)

$$S(A:B) = S(\rho_A) + S(\rho_B) - S(\rho_{AB}) = H(\rho_{AB}, \rho_A \otimes \rho_B) \geq 0$$

Conditional Information: (always positive for classical systems)

$$\begin{aligned} S(A|B) &= S(\rho_{AB}) - S(\rho_B) = -H(\rho_{AB}, \rho_B) \\ &= -H(\rho_{AB}, \tfrac{1}{d} I_A \otimes \rho_B) + \log d \end{aligned}$$

Max entangled state $\quad S(\rho_{AB}) - S(\rho_B) = -\log d < 0$

$\rho_{AB} = |\psi_{AB}\rangle\langle\psi_{AB}| \quad \psi_{AB} = \sum_k \frac{1}{d} |\phi_k \otimes \phi_k\rangle$

Conditional Info $S(\rho_{AB}) - S(\rho_B)$ concave

surprising since diff of concave functions – equiv. to SSA

Weak monotonicity of conditional Information

Classical $\quad S(\rho_{AB}) - S(\rho_B) \geq 0$

Quantum given any ρ_{ABC}

Purify $\rho_{ABC} = \sum_k \lambda_k |\phi_k\rangle\langle\phi_k|$ to ρ_{ABCD} s.t. $\text{Tr}_B\, \rho_{ABCD} = \rho_{ABC}$

$\rho_{ABCD} = |\psi\rangle\langle\psi| \qquad |\psi_{ABCD}\rangle = \sum_k \sqrt{\lambda_k}\, |\phi_l \otimes f_k\rangle$

where $|f_k\rangle$ O.N. in space iso to $(\ker \rho_{ABC})^\perp \subseteq \mathcal{H}_{ABC}$.

SSA:

$$\begin{aligned} S(\rho_{ABC}) + S(\rho_B) &\leq S(\rho_{AB}) + S(\rho_{BC}) \\ S(\rho_D) + S(\rho_B) &\leq S(\rho_{CD}) + S(\rho_{BC}) \\ S(\rho_{CD}) - S(\rho_D) + S(\rho_{BC}) - S(\rho_B) &\geq 0 \end{aligned}$$

two cond entropies with common subsystem can't both be negative

Negative Conditional Information

QIT uses EPR states $\frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$ to transmit

a) classical information by process called "dense coding"

b) unknown quantum states by process called "teleportation"

HOW (M. Horodecki, J. Oppenheim, A. Winter) interpretation:

Nature **436**, 673–676 (2005); CMP **269** , (2007). quant-ph/0512247.

Cond info measures # of bits Alice needs to transmit message to Bob when he has partial info – same interp class. and quant info

When negative, gives # of EPR pairs A and B have left for future

Aside: Operator-Monotone and -Convex Functions

$A > B > 0 \;\Rightarrow\; \sqrt{A} > \sqrt{B} > 0$ but $\not\Rightarrow\; A^2 > B^2$

Def: $f : (a,b) \mapsto \mathbf{R}$ is operator-monotone if (a.k.a. Pick or Herglotz)

$A > B > 0 \;\Rightarrow\; f(A) > f(B).$

Thm: Let $f : (0,\infty) \mapsto \mathbf{R}$. TFAE the following are equivalent

a) f is operator monotone

b) f can be anal cont into UHP and maps UHP into UHP

$$\begin{aligned} \text{c) } f(x) &= ax + \int_0^\infty \tfrac{xu-1}{x+u}\, \nu(u)du \qquad \text{with} \quad \nu(u) \geq 0 \\ &= a'x - \int_0^\infty \tfrac{1}{x+u}(1+u^2)\nu(u)du \quad \text{if} \quad \int u\,\nu(u)du < \infty \end{aligned}$$

where UHP $= z : \text{Im} z > 0 \qquad$ (b) $\Leftrightarrow$ (c) Nevanlinna's Thm.

Operator-Monotone and -Convex Functions continued

$A > B > 0 \Rightarrow X = B^{-1/2}AB^{-1/2} > I \Rightarrow$ e-vals of $X > 1$

e-vals of $X^{-1} < 1 \Rightarrow B^{1/2}A^{-1}B^{1/2} < I \Rightarrow 0 < A^{-1} < B^{-1}$

So $f(x) = -(x+u)^{-1}$ is op-mon which gives (c) $\Rightarrow$ (a).

Def: $g : (0,\infty) \mapsto \mathbf{R}$ operator-convex if

$$g(xA_1 + (1-x)A_2) \leq xg(A_1) + (1-x)g(A_2)$$

Roughly g is op-convex iff suitable diff quot is op-mon.

$g : [0,\infty) \mapsto \mathbf{R}$ and $g(0) = 0$ is op-convex iff

$$\frac{g(x) - g(0)}{x - 0} = \frac{g(x)}{x} \text{ op-mon}$$

Theory due to Löwner   Ando 197? notes

Entropy concave : Proof # 3

$g : [0,\infty) \mapsto \mathbf{R}$ and $g(0) = 0$ is op-convex iff $\frac{g(x)}{x}$ op-mon

Apply to $g(x) = x \log x$

$\frac{g(x)}{x} = \log x$: UHP $\mapsto \{z : \mathrm{Im} z \in (0, 2\pi)\} \subset$ UHP

Find: $g(x) = x \log x$ op-convex $g(x) = -x \log x$ op concave

$\Rightarrow$ and stronger than $S(R) = -R \log R$ concave   C. Davis

	$-1 < p < 0$	op-convex	and	op-mon-dec
$f(x) = x^p$	$0 < p < 1$	op-concave	and	op-mon
	$1 < p \leq 2$	op-convex	but	not op-mon

$f(x) = x^p$ neither op-convex nor op-mon for $p > 2$.

III: Simple Proof of Joint Convexity of Relative Entropy

1. Background and WYD
2. Unified proof of joint conv. of $\mathrm{Tr}\, K^* R^p K Q^{1-p}$ and $H(R, Q)$
   - 2.1 Pedestian modular operator via left and right mult.
   - 2.2 Integral representations
   - 2.3 Convexity of $J_p(K, A, B)$
3. Comments on Schwarz inequalities
4. $q \neq 1 - p$
5. Monotonicity under partial traces
6. More history ?
7. Lieb's golden corollary
8. Equality conditions

Background to proofs

Original proof of SSA based on Lieb's result below

WYD skew entropy $\frac{1}{2}\mathrm{Tr}\, [K, \gamma^p][K, \gamma^{1-p}]$

for $K = K^*$ and γ a density matrix

Wigner-Yanase introduced for $p = \frac{1}{2}$ and proved concave in γ.

Dyson suggested $p \in (0, 1)$ – led to conjecture

Conj: $\gamma \mapsto \mathrm{Tr}\, K\gamma^p K\gamma^{1-p} - \mathrm{Tr}\, K\gamma K$ concave

Lieb dropped linear term and proved generalization

$(A, B) \mapsto \mathrm{Tr}\, K^* A^p K B^{1-p}$ concave for $p \in (0, 1)$

Claim: But advantage to retaining linear term !!

Retaining the linear term

$$J_p(K,A,B) = \frac{1}{p(1-p)}\left[\mathrm{Tr}\, K^*AK - \mathrm{Tr}\, K^*A^pKB^{1-p}\right]$$

Note: well-def for $p > 0$ and factor $(1-p)$ changes sign at $p = 1$

Thm: $(A,B) \mapsto J_p(K,A,B)$ is convex for $p \in (0,2)$

$\Rightarrow$ $\mathrm{Tr}\, K^*A^pKB^{1-p}$ concave for $p \in (0,1)$

$\Rightarrow$ $\mathrm{Tr}\, A(\log A - \log B)$ convex $p = 1$ – extend by cont $K = I$

$\Rightarrow$ $\mathrm{Tr}\, K^*A^pKB^{1-p}$ convex for $p \in (1,2]$ and $p \in [-1,0)$

will give proof, which is elementary, short, and sweet

$\mathrm{Tr}\, A = \mathrm{Tr}\, B \quad \Rightarrow \quad J_p(I,A,B) \geq 0$ with equality $\Leftrightarrow A = B$

pseudo-metric in same sense as relative entropy $H(A,B)$ gen Klein

Pedestrian modular operator

$d \times d$ matrices form Hilbert space with $\langle A, B\rangle = \mathrm{Tr}\, A^*B$

Def. Left and Right mult as linear operators on this vector space

$$L_A(X) = AX \qquad \text{and} \qquad R_B(X) = XB$$

a) L_A and R_B commute $L_A[R_B(X)] = AXB = R_B[L_A(X)]$

b) $A = A^* \Rightarrow L_A, R_A$ self-adjoint wrt H-S inner prod

For $A, B > 0$ positive definite

c) L_A, R_A pos def $\langle X, R_A(X)\rangle = \mathrm{Tr}\, X^*XA = \mathrm{Tr}\, XAX^* \geq 0$

d) $(L_A)^{-1} = L_{A^{-1}}, \quad (R_B)^{-1} = R_{B^{-1}}$

e) $f(L_A) = L_{f(A)} \quad f(R_B) = R_{f(B)}$, e.g., $L_A^p = L_{A^p}, R_A^p = R_{A^p}$

simple form of deep idea: Araki $\Delta_{AB} = L_A R_B^{-1}$ relative modular op

Back to $J_p(K.A.B)$

$$g_p(x) = \begin{cases} \frac{1}{p(1-p)}(x - x^p) & p \neq 1 \\ x\log x & p = 1 \end{cases}.$$

well-defined for $x > 0$ and $p \neq 0$, but $p \in [\frac{1}{2}, 2]$ would suffice

$$J_p(K,A,B) \equiv \mathrm{Tr}\sqrt{B}K^*\, g_p(L_A R_B^{-1})(K\sqrt{B})$$

$$= \begin{cases} \frac{1}{p(1-p)}(\mathrm{Tr}\, K^*AK - \mathrm{Tr}\, K^*A^pKB^{1-p}) & p \in (0,1)\cup(1,2) \\ \mathrm{Tr}\, KK^*A\log A - \mathrm{Tr}\, K^*AK\log B) & p = 1 \\ -\frac{1}{2}(\mathrm{Tr}\, K^*AK - \mathrm{Tr}\, AKB^{-1}K^*A) & p = 2 \end{cases}$$

$$J_1(I,A,B) = \mathrm{Tr}\, A(\log A - \log B) = H(A,B)$$

Aside: extend to $[-1.0)$

not quite symmetric around $p = \frac{1}{2} \qquad p \leftrightarrow 1-p$

$$\widetilde{g}_p(x) = w\, g_{1-p}(w^{-1}) = \begin{cases} \frac{1}{p(1-p)}(1 - w^p) & p \neq 0 \\ -\log w & p = 0 \end{cases} \qquad p \in [-1,1)$$

$J_p(K,B,A) = \widetilde{J}_{1-p}(K^*,A,B)$

$\widetilde{J}_p(K,A,B)$ jointly convex for $p \in [-1,1)$

$\widetilde{J}_0(I,A,B) = \mathrm{Tr}\, B(\log B - \log A) = H(B,A)$

Integral representations

$$\frac{g_p(x)}{x} = \begin{cases} \frac{1}{p(1-p)}(1-x^{p-1}) & p \neq 1 \\ \log x & p = 1 \end{cases}$$

well-def for $x \in (0,\infty)$ and operator monotone for $p \in (0,2]$, or,

anal cont to upper half of complex plane and UHP $\mapsto$ UHP

$\Rightarrow$ $g_p(x)$ has integral rep of form

$$\begin{aligned} g_p(x) &= ax^2 + \int_0^\infty \frac{x^2 t - x}{x+t}\,\nu(t)\,dt \\ &= ax^2 + \int_0^\infty \left[\frac{x^2}{x+t} - \frac{1}{t} + \frac{1}{x+t}\right] t\nu(t)\,dt \end{aligned}$$

with $\nu(t) \geq 0$

Specific integrals – elementary

$$\int_0^\infty \frac{x^{p-1}}{x+1} = \frac{\pi}{\sin p\pi} \qquad 0 < p < 1 \qquad c_p = \frac{\sin p\pi}{\pi}$$

allows us to give the following explicit representations

$$g(x) = \begin{cases} \frac{1}{p(1-p)}\left[x + c_p \int_0^\infty \left(\frac{t}{x+t} - 1\right) t^{p-1} dt\right] & p \in (0,1) \\ \int_0^\infty \left(\frac{x^2}{x+t} - 1 + \frac{t}{x+t}\right) \frac{1}{1+t} dt & p = 1 \\ \frac{1}{p(1-p)}\left[x - c_{p-1} \int_0^\infty \frac{x^2}{x+t} t^{p-2} dt\right] & p \in (1,2) \\ \frac{1}{2}(-x + x^2) & p = 2 \end{cases}$$

Important: For $p \in (0,2)$ integrand supported on $(0,\infty)$.

Integal representation using L_A and R_B

Recall $J_p(K,A,B) = \operatorname{Tr} \sqrt{B} K^* g_p(L_A R_B^{-1})(K\sqrt{B})$

$$g_p(x) = ax + \int_0^\infty \left[\tfrac{x^2}{x+t} - \tfrac{1}{t} + \tfrac{1}{x+t}\right] t\nu(t)\,dt$$

$$\begin{aligned} \operatorname{Tr} \sqrt{B} K^* \frac{1}{L_A R_B^{-1} + tI}(K\sqrt{B}) &= \operatorname{Tr} \sqrt{B} K^* \frac{R_B}{L_A + tR_B}(K\sqrt{B}) \\ &= \operatorname{Tr} BK^* \frac{1}{L_A + tR_B}(KB) \end{aligned}$$

$$\begin{aligned} J_p(K,A,B) &= \operatorname{Tr} K^*AK - \operatorname{Tr} KBK^* \int_0^\infty \nu(t)\,dt \\ &+ \int_0^\infty \left[\operatorname{Tr} K^*A \frac{1}{L_A + tR_B}(AK) + \operatorname{Tr} BK^* \frac{1}{L_A + tR_B}(KB)\right] t\,\nu(t)\,dt \end{aligned}$$

Suffices to show $(A,B,X) \mapsto \operatorname{Tr} X^* \frac{1}{L_B + tR_A}(X)$ jointly convex

Proof:

Note: $\operatorname{Tr} (\lambda X)^* \frac{1}{L_{\lambda B} + tR_{\lambda A}}(\lambda X) = \lambda \operatorname{Tr} X^* \frac{1}{L_B + tR_A}(X)$

Homo of degree 1 $\Rightarrow$ suffices to prove "subadditivity" (omit x_k)

Let: $M = (\ )^{-1/2}(X) - (\ )^{1/2}(\Lambda)$

$$\begin{aligned} \operatorname{Tr} M^*M &= \langle M, M\rangle \\ &= \langle [(\ )^{-1/2}(X) - (\ )^{1/2}(\Lambda)], [(\ )^{-1/2}(X) - (\ )^{1/2}(\Lambda)]\rangle \\ &= \langle X, (\ )^{-1}(X)\rangle - \langle X, \Lambda\rangle - \langle \Lambda, X\rangle + \langle \Lambda, (\ )(\Lambda)\rangle \end{aligned}$$

Choose $M = (L_A + tR_B)^{-1/2}(X) - (L_A + tR_B)^{1/2}(\Lambda)$

$\operatorname{Tr} M^*M =$

$$\operatorname{Tr} X^*(L_A + tR_B)^{-1}(X) - \operatorname{Tr} X^*\Lambda - \operatorname{Tr} \Lambda^*X + \operatorname{Tr} \Lambda^*(L_A + tR_B)(\Lambda)$$

Let $M_j = (L_{A_j} + tR_{B_j})^{-1/2}(X_j) - (L_{A_j} + tR_{B_j})^{1/2}(\Lambda)$. Then

$$0 \le \sum_j \mathrm{Tr}\, M_j^* M_j = \sum_j \mathrm{Tr}\, X_j^*(L_{A_j} + tR_{B_j})^{-1}(X_j)$$
$$-\mathrm{Tr}\,(\textstyle\sum_j X_j^*)\Lambda - \mathrm{Tr}\,\Lambda^*(\sum_j X_j) + \mathrm{Tr}\,\Lambda^* \sum_j (L_{A_j} + tR_{B_j})\Lambda$$
$$= \sum_j \mathrm{Tr}\, X_j^* \frac{1}{(L_{A_j} + tR_{B_j})}(X_j) - \mathrm{Tr}\, X^*\Lambda - \mathrm{Tr}\,\Lambda^* X + \mathrm{Tr}\,\Lambda^*(L_A + tR_B)\Lambda$$

Choose $\Lambda = \dfrac{1}{L_A + tR_B}(X) \qquad X = \sum_j X_j, \ \sum_j L_{A_j} = L_{\sum_j A_j} = L_A$

$$\mathrm{Tr}\,\Lambda^* \textstyle\sum_j (L_{A_j} + tR_{B_j})\Lambda = \mathrm{Tr}\, X^* \frac{1}{L_A + tR_B} X = \mathrm{Tr}\, X\Lambda = \mathrm{Tr}\,\Lambda^* X$$

$$0 \le \sum_j \mathrm{Tr}\, X_j^* \frac{1}{L_{A_j} + tR_{B_j}}(X_j) - \mathrm{Tr}\, X^* \frac{1}{L_A + tR_B}(X)$$

compare elementary C-S ineq:

$$|\sum_k \overline{v}_k w_k|^2 \le \sum_k |v_k|^2 \sum_k |w_k|^2$$

For $a_k > 0$ let $v_k = a_k^{1/2}$, $w_k = a_k^{-1/2} x_k$

$$|\sum_k x_k|^2 \le \sum_k a_k \sum_k \overline{x}_k \frac{1}{a_k} x_k$$

Rewrite $\quad (\sum_k \overline{x}_k) \dfrac{1}{\sum_k a_k} (\sum_k x_k) \le \sum_k \overline{x}_k \dfrac{1}{a_k} x_k$

Lieb and Ruskai (1973) proved operator version

$$(\sum_k X_k^*) \frac{1}{\sum_k A_k} (\sum_k X_k) \le \sum_k X_k^* \frac{1}{A_k} X_k$$

Not suff. for SSA — need Araki rel mod op hidden in L_A and R_B.
Compare proof: $\left|\sum_k v_k + tw_k\right|^2 \ge 0 \ \forall\, t$ choose t to minimize

Remarks on $q \neq 1 - p$

$p, q > 0, \ p + q < 1 \qquad \mathrm{Tr}\, K^* A^p K B^{1-p}$ concave

Write $\mathrm{Tr}\, K^* A^p K B^q = \mathrm{Tr}\, K^* A^p K (B^s)^{1-p} \qquad 0 < s = \frac{1}{1-p} < 1$

B^s is op monotone and op concave for $s \in (0, 1)$

$$(\lambda B_1 + (1-\lambda) B_2)^s > \lambda B_1^s + (1-\lambda) B_2^s$$

$$\begin{aligned}
\mathrm{Tr}\, K^* A^p K B^q &= \mathrm{Tr}\, K^* A^p K \left[(\lambda B_1 + (1-\lambda)B_2)^s\right]^{1-p} \\
&> \mathrm{Tr}\, K^* A^p K (\lambda B_1^s + (1-\lambda) B_2^s)^{1-p} \\
&\ge \lambda \mathrm{Tr}\, K^* A_1^p K (B_1^s)^{1-p} + (1-\lambda) \mathrm{Tr}\, K^* A_2^p K (B_2^s)^{1-p} \\
&= \lambda \mathrm{Tr}\, K^* A_1^p K B_1^q + (1-\lambda) \mathrm{Tr}\, K^* A_2^p K B_2^q
\end{aligned}$$

Note: $f(x)$ strictly concave and op concave $\Rightarrow$ strict op ineq

get equal only for trivial cases, $B_1 = B_2$ or $\lambda = 0, 1$.

Monotonicity under partial traces

Define generalized Pauli (Weyl-heisenberg) operators,

$$Z|e_n\rangle = e^{2\pi i n/d}|e_n\rangle \qquad X|e_n\rangle = |e_{n+1}\rangle$$

$$\sum_j Z^j A Z^{-j} = d\, A_{\mathrm{diag}} \qquad \sum_j X^j A_{\mathrm{diag}} X^{-j} = (\mathrm{Tr}\, A) I$$

$$\tfrac{1}{d} \sum_j \sum_k X^j Z^k A (X^j Z^k)^* = (\mathrm{Tr}\, A)\, I = \tfrac{1}{d} \sum_n W_n A W_n^*$$

$W_n = X^j Z^k$ in some ordering $n = 1, 2, \ldots d^2$, e.g., $n = j + d(k-1)$

$$\tfrac{1}{d_2} \sum_n (I_1 \otimes W_n) A_{12} (I_1 \otimes W_n)^* = A_1 \otimes I_2$$

Discrete version of Uhlmann's observation that partial trace can be obtained by integrating over $SU(n)$ using Haar measure.

$$
\begin{aligned}
J_p(K_2, A_2, B_2) &= J_p(I_1 \otimes K_2, \tfrac{1}{d_1} I_1 \otimes A_2, \tfrac{1}{d_1} I_1 \otimes B_2) \\
&= \tfrac{1}{d_1^2} J_p\Big(K_{12}, \textstyle\sum_n (W_n \otimes I_2) A_{12} (W_n \otimes I_2)^*, \sum_n (W_n \otimes I_2) B_{12} (W_n \otimes I_2)^*\Big) \\
&\leq \tfrac{1}{d_1^2} \sum_n J_p(I_1 \otimes K_2, (W_n \otimes I_2) A_{12} (W_n \otimes I_2)^*, (W_n \otimes I_2) B_{12} (W_n \otimes I_2)^*) \\
&= \tfrac{1}{d_1^2} \sum_n J_p(I_1 \otimes K_2, A_{12}, B_{12}) = J_p(I_1 \otimes K_2, A_{12}, B_{12})
\end{aligned}
$$

used $\quad J_p(I_1 \otimes K_2, A_{12}, B_{12}) \qquad$ wrote $K_{12} = I_1 \otimes K_2$

$\quad = J_p(I_1 \otimes K_2, (W_n \otimes I_2) A_{12} (W_n \otimes I_2)^*, (W_n \otimes I_2) B_{12} (W_n \otimes I_2)^*)$

$J_1(I, A_2, B_2) \leq J_1(I, A_{12}, B_{12})$ gives $H(A_2, B_2) \leq H(A_{12}, B_{12})$

Cor: SSA $\quad H(A_{23}, A_2) \leq H(A_{123}, A_{13})$

"no transparent proof of SSA is known"

p. 645 of *Quantum Computation and Quantum Information*

Michael A. Nielsen and Isaac L. Chuang (Cambridge Press, 2000)

based on B. Simon's version adapted from Uhlmann (1977) of

"elementary" proof of $(A, B) \mapsto \mathrm{Tr}\, K^* A^p K B^{1-p}$ concave

similar argument in Wehrl *Rev. Mod. Phys* (1978). BUT

- MBR, "Lieb's simple proof of concavity ..." quant-ph/0404126

Int. J. Quant Info. **3**, 579–590 (2005) Schwarz + max mod

- Ando's argument described in Carlen's Tucson notes
- Petz – uses Δ_{AB} in book; elem version in quant-ph/0408130
- Proof here based on Schwarz ineq. using L_A, R_B really elem.

based on Lesniewski and Ruskai, JMP; and MBR quant-ph/0604206

Aside: Equality conditions in $J_p(K.A.X)$ convex

$$
\begin{aligned}
\int_0^\infty \mathrm{Tr}\, K^* A \frac{1}{L_A + t R_B} (AK) \nu(t) dt &\leq \int_0^\infty \sum_j \mathrm{Tr}\, (A_j K)^* \frac{1}{L_{A_j} + t R_{B_j}} (A_j K) \nu(t) \\
&= \sum_j \int_0^\infty \mathrm{Tr}\, (A_j K)^* \frac{1}{L_{A_j} + t R_{B_j}} (A_j K) \nu(t) d
\end{aligned}
$$

Equal $\Leftrightarrow$ equal for each term in integ , i.e., $M_j = 0 \quad \forall j, \forall t$

$$(L_{A_j} + t R_{B_j})^{-1}(X_j) = (L_A + t R_B)^{-1}(X) \qquad \forall j, \forall t$$

equality conditions independent of $p \in (0, 2)$

$X = AK \qquad (I + t \Delta_{A_j B_j}^{-1})^{-1}(K) = (I + t \Delta_{AB}^{-1})^{-1}(K) \qquad \forall j, \forall t$

$X = BK \qquad (\Delta_{A_j B_j} + tI)^{-1}(K) = (\Delta_{AB} + tI)^{-1}(K) \qquad \forall j, \forall t$

Recall $\Delta_{AB} = L_A R_B^{-1} > 0$ prod of commuting pos def ops

$$(\Delta_{A_j B_j} + tI)^{-1}(K) = (\Delta_{AB} + tI)^{-1}(K) \qquad \forall j, \forall t$$

$\Delta_{AB} > 0 \quad \Rightarrow \quad (\Delta_{AB} + tI)^{-1}$ anal cont to $\mathbf{C} \backslash (-\infty, 0]$

can apply Cauchy integral Thm. to get

$\Rightarrow \; G(\Delta_{A_j B_j})(K) = G(\Delta_{AB})(K) \quad \forall j \quad G$ anal on $\mathbf{C} \backslash (-\infty, 0]$

allows several useful formulations

$\Rightarrow \; (\Delta_{AB} + tI)$ and $(I + \Delta_{AB}^{-1} t)$ forms equiv.

Equivalent equality conditions

Thm: For fixed K, and $A = \sum_j A_j, B = \sum_j B_j$ TFAE

a) $J_p(K, A, B) = \sum_j J_p(K, A_j, B_j)$ for all $p \in (0, 2)$.

b) $J_p(K, A, B) = \sum_j J_p(K, A_j, B_j)$ for some $p \in (0, 2)$.

c) $(\Delta_{A_j B_j} + tI)^{-1}(K) = (\Delta_{AB} + tI)^{-1}(K) \quad \forall\, j$ and $\forall\, t > 0$.

d) $A_j^{it} K B_j^{-it} = A^{it} K B^{-it} \quad \forall\, j$ and $\forall\, t > 0$.

e) $(\log A - \log A_j)K = K(\log B - \log B_j) \quad \forall\, j$.

In addition when $K = I$, equiv to

f) There are $D_j > 0$ such that $[A_j, D_j] = [B_j, D_j] = 0$, and

$A_j = AD^{-1}D_j, \quad B_j = BD^{-1}D_j$ with $D = \sum_j D_j$

neccessity of (f) uses sufficient subalgebra – developed by Petz

formulation here from Jenčová and Petz, CMP, **263**, 259–276 (2006).

Equality conditions for SSA

use form $\log A_{123} - \log A_{12} - \log A_{23} + \log A_2 = 0$

Easy to see $A_{123} = A_1 \otimes A_{23}$ or $A_{12} \otimes A_3$ will suffice

If $\mathcal{H}_2 = \mathcal{H}_{2_L} \otimes \mathcal{H}_{2_R}$ then $A_{123} = A_{12_L} \otimes A_{2_R 3}$ will suffice

Thm: Equality holds in SSA if and only if

$$\mathcal{H}_2 = \bigoplus_n \mathcal{H}_n^L \otimes \mathcal{H}_n^R \quad \text{and} \quad A_{123} = \bigoplus_n A_n^L \otimes A_n^R$$

with $A_n^L \in \mathcal{B}(\mathcal{H}_1 \otimes \mathcal{H}_n^L), \quad A_n^R \in \mathcal{B}(\mathcal{H}_n^R \otimes \mathcal{H}_3)$

Cor: Equality in $\operatorname{Tr} A_{23}^p A_2^{1-p} \leq \operatorname{Tr} A_{123}^p A_{12}^{1-p}$ iff same cond

$p \in (0, 1) \ \leq \qquad p \in (1, 2) \ \geq$

Other approaches to JC of $\operatorname{Tr} K^* A^p K B^{1-p}$

Lieb; based on maximum modulus principle –

basic result rel easy in finite dim.

Epstein; based on op-monotone or Herglotz functions

deep and not easy – paper must be read backwards !

Uhlmann - Simon elementary, but long and not insightful

gave bad rep – see Nielsen-Chuang quote

Ando; used iso between $\mathcal{B}(\mathcal{H}_B, \mathcal{H}_A) \simeq \mathcal{H}_A \otimes \mathcal{H}_B$ to observe

$\operatorname{Tr} K^* A^p K B^{1-p} = \langle K, A^p \otimes B^{1-p} K \rangle$

where K interp as vec in $\mathcal{H}_A \otimes \mathcal{H}_B$

to give linear algebra proof — first with no complex anal.

Differentiatie $\log A$ for matrix

$$\begin{aligned}
\log(R + xQ) - \log R &= \int_0^\infty \left(\frac{1}{R + uI} - \frac{1}{R + xQ + uI} \right) du \\
&= \int_0^\infty \frac{1}{R+uI} [(R + xQ + uI) - (R + uI)] \frac{1}{R + xQ + uI} du \\
&= \int_0^\infty \frac{1}{R + uI} xQ \frac{1}{R + xQ + uI} du
\end{aligned}$$

Can take $\lim_{x \to 0} \frac{1}{x}(\log(R + xQ) - \log R$ and find

$$\log(R + xQ) = \log R + x \int_0^\infty \frac{1}{R + uI} Q \frac{1}{R + uI} du + O(x^2)$$