Density Matrix Renormalization Group

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Spin systems
Exact diagonalization
Entanglement
Area Law
Matrix Product States
DMRG

Julia for fast development C++/ITensor for state of the art efficiency

Overview: QM	of many	electron
systems		
All electrons Weak Correlation {	DFT Model OMFT, etc	Hamiltonians
Quc ED		
		* DMRG *
Background:	ED, Ma	odel Systems,
New Connec	tr; Qua	utum Informal
Todays Topics		Entanglement
ED/Lanczo.	s /c	elic programming
Entanglement		rple exercises
Tomorrow: Matr. DMRG, ITen	son library	5+eto, (C++)

Exact Diagonalization of Small clusters of spin-1/2's, Review: 5= 2, Paul: matrices Spin operator $\vec{S} = \frac{\vec{h}}{2} \vec{\sigma}$ Spin basis: $\Lambda(a) = a/17 + 5/47$ $\Lambda(b) = a/17 + 5/47$ Oz = (1 0) Oz is diagony in this 2-54515, e.g. $\sigma_2 | \uparrow \rangle = | \uparrow \rangle$ But ox 11) = 11> not diagonal Two S=1's: Suppose H=\$1.52. What are the energy lovels? Algebraic Stotal = S, + S2 spins 1 and 2 5, + 3, commute, e.g. [Siz, 527] = 0 52 = Stotal · Stotal = (S, +32) · (S, +32) $=\vec{S}_{1}^{2}+\vec{S}_{2}^{2}+Z\vec{S}_{1}\cdot\vec{S}_{3}$

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Let
$$\vec{S} = \frac{1}{2}\vec{\sigma}$$
, $\vec{S} = S_x \pm iS_y$, etc
 $S_z = 1(\frac{1}{2}\vec{o})$ $S = (\frac{1}{2}\vec{o})$
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Solution H= J (4 0 0 0) Viagnalize: Singlet: \$\frac{1}{\pi_1} = \frac{1}{\pi_2} Julia: H=[1/4 0 0 0; 0-1/4 1/2 0; 0 1/2 -1/4 0: 0 0 0 1/4 eig fact (H)

Many spins: Hilbert Space = { 15, 52... Su} The key difficulty of many pHe QM is the exponentially large size of the Hilbert Space $|7\rangle = 2^{N} - \log vector = \begin{cases} q & \uparrow \uparrow \uparrow \\ 5 & \uparrow \uparrow \downarrow \\ c & \uparrow \downarrow \downarrow \uparrow \\ d & \uparrow \downarrow \downarrow \\ \end{cases}$ H = 2" x 2" matrix (Find lovest E eigenvector) Directly treating N=100 this way is hopeless. But lets see how to do a cluster of 3 or 4 or 10 spires. N=2 come: algebra al reals done note substantial sparseness The elements are < s's' | H | 5,52) w.th 5,= 1 a J, etc.

N=3 H=(3,-32+52·53) J 5,.52 only operates on 142 - It leves spin 3 alone. Thus < 5, 5, 5, 5, 5, 5, 5, 5 oc 65,5, The Z-Z part of 5,5 doent chaye Si or Sz. It give a factur of # 4 for parallel on opp. spirs. The = (5+5+55+) + fl. 71 the spres and puts in +1. These obserations turn into simple rules for the elements of H 1) Diagonal elements are I. (# powallel -)

opposite
heishhars

2) Off diag are zero if more than
two spins differ. It two n.n.
spins are flipped, get + 1
Otherwise, zero. > (11) (17) not 1773

				(23					
	Ex.	azile	N=3	-	• — •		Let	0=	1, 1=	<i>J</i>
		600	001	010	011	100	101	110	(111_	7
	000	2-4	0	0	0	0	0	0	0	
	001	0	0	1/2	0	0	0	D	0	
14	010	0	1/2	-2-4	0	1/2	0	0	0	
H=	011	O	0	O	0	0	1/2	0	0	1.5
	100	0	O	1/2	Ø	0	0	O	0	
	101	0	0	0	1/2	0	-2.4	1/2	0	
	110	0	6	O	0	0	1/2	0	0	
	111	6	0	0	0	0	0	0	2.4	
-	Service Control									

We con call our eigensolver on this to get the energies + eigenvectors

Conservation of State

H does not change the number of up and down spirs.

[H, StoH] = 0

This makes H Block diagnal

For example the 111 block for N-7 111 117 111

This is a big help, but the block sizer are still $O\left(\frac{2^{N}}{N}\right)$ For 5% + old = 0, there are 6 basis 5 totos, 0011, 0101, etc. = 1711 1111 Find the 6 x6 Hblock, diagonalize it with julia, get the ground state energy + vector. Exercise: For a chan of N spins (507 up to N=10) write a julia function to Hard give H and find the ground's state energy,

Thus $2^N = 30K \rightarrow N^{N-15}$ Story Calculator time: Full diagonalizator, MXN $NM^3 = CPU-tre$ $M \sim 10^{4.5-3}$ $M \sim 10^{13.5}$	Memory of Calculation Time
A desktop has $n/0'0$ bytes I double prec ~ 10 bytes (actually 8) Can stope a 10^9 reals; on $10^{4/5} \times 10^{4/5}$ matrix or ~30k × 30k Thus $2^N = 30^{\frac{1}{5}} \rightarrow N^{-1/5}$ storage Calculated time: Full diagonalization, $M \times N$ $N = N^3 = CPU + tree $	For an MXM metrix, storage is M2
Can store a 109 rects; on 1045 x 104 matr. or ~30K x 30K Thus 2N = 30K -> N~15 storage Calculator time: Full diagonalizator, Mx N NM3 = CPU-tre M~104.5-3 ~ 1013.5 Calc 1013.5 time ~ 1010 ~ # of floating pt. op. 10 GFLOPS ~10 GFLOPS ~103.5 ~ 3000 sec ~ 1 hr. OK How Can we do better? 1) Spirseness: Each row of H has only O(N) non zero els, say N~20 Storage -> N 2N ~ D N+4	A Jesktop has ~1000 bytes
Can store a 109 rects; on 1045 x 104 matr. or ~30K x 30K Thus 2N = 30K -> N~15 storage Calculator time: Full diagonalizator, Mx N NM3 = CPU-tre M~104.5-3 ~ 1013.5 Calc 1013.5 time ~ 1010 ~ # of floating pt. op. 10 GFLOPS ~10 GFLOPS ~103.5 ~ 3000 sec ~ 1 hr. OK How Can we do better? 1) Spirseness: Each row of H has only O(N) non zero els, say N~20 Storage -> N 2N ~ D N+4	I double prec ~ 10 bytes (actually 8)
Thus $2^{N} = 30^{\frac{1}{K}} \rightarrow N^{-15}$ Storage Calculator time: Full diagonalization, $M \times N^{-1}$ $N = N^{\frac{3}{2}} = CP4$ the $N^{\frac{3}{2}} = N^{\frac{3}{2}} =$	Can store a 109 reals; or 10 4.5 x 10 mater
Calculator time: Full diagonalizator, MXN NM3 = CPU-time 3 ~ 104.5-3 ~ 1013.5 Calc 1013.5 time ~ 1010 ~ # of floating pt. op. 10 GFLOPS NOW Can we do better? 1) Sparseness: Each row of H has only O(N) non zero els, say N~20 Storage > N2N ~ D2N+4	or ~30K × 30K
Calculator time: Full diagonalizator, MXN NM3 = CPU-time M2/04.5-3 Calc time n 1013.5 1000 per second 10 GFLOPS NOW Can we do better? 1) Sparseness: Each row of H has only O(N) nonzero els, say N~20 Storage > N2N ~ D2N+4	Thus 2N= 30k -> NN 15 Storage
Calc 10 ^{13.5} time ~ 10 ^{13.5} 10'0 # of floating pt. of 10 GFLOPS ~10 ^{3.5} ~ 3000 sec ~ 1 hr. OK How can we do botter? 1) Sparseness: Each row of H has only O(N) non zero els, say N~20 Storage > N2 ^N ~ D N+4	Calculator time: Full diagonalizator, MXN
time ~ 10'0 / 10'0 # of floating pt. of. 10 GFLOPS ~ 103.5 ~ 3000 sec ~ 1 hr. OK How can we do better? 1) Sparseness: Each row of H has only O(N) non zero els, say N~20 Storage ~ N 2 ~ D > N+4	M = cpu-tre M~104.5-5 13.5
How can we do better? 1) Sparseness: Each row of H has only O(N) nonzero els, say N~20 Storage > N2N ~ D>N+4	time ~ 1013.5 100 ~ # of floating pt. of
How can we do better? 1) Sparseness: Each row of H has only O(N) nonzero els, say N~20 Storage > N2N ~ D>N+4	10 GFLOPS
How can we do botter? 1) Sparseness: Each row of H has only O(N) nonzero els, say N~20 Storage > N2N ~ D>N+4	~103.5 ~ 3000 sec ~ 1 hr. OK
1) Sparseness: Each row of H has only O(N) nonzero els, say N~20 Storage > N2N ~ D>N+4	
Storage > N2"~ DON+4	HOW Can we do better?
Storage > N2"~ DON+4	1) Sparseness: Each row of H has
Storage > N2"~ DON+4	only O(P) nonzero els, say N~20
But can't use full diag: millions if his	Storge > N2"~ DON+4
But can't use full drag: millions of b.	say up to N~25 much better 1
Store, mours	But can't use full diag: millions of hours,

Power Methol

Let H = (1 - EH). For E small enough, the largest magnitude eigenvalue is 1- E Egr, with same "grout" state Then 16.5> = 1im H 18> Where 13> is any vector not 1 to 165). The power method only stores (sparse) Hand one or two IV> IV'>= HIV> iterate. But it is slow; lainge k needed.

Lanczos Method

Given 18>=8, the Krylou space is Z's, H's, H's, H's, ... } + all linear (vectors opece) Note (1- EH) & is in the Knylou space. Lanczos is an efficient way of finding the lowest energy vector in a truncate

Krylov space (up to Hkg for some le)

Lanctos varely needs k bigger than k-100-200 for the ground state, much better than the power method Lanczos besis: d chasen so $\hat{X}_2 = \mathcal{N}_2 \left(H \hat{X}_1 - \alpha_2 \hat{X}_1 \right)$ X, L Xz normalizatu fata $X_3 = \mathcal{N}_s (H\vec{X}_2 - \alpha_3 \hat{x}_1 - \beta_3 \hat{x}_1)$ X4 = 74 (H x3 - 4x3 - 8 x2 - 84 x,) this mokes an orthonormal basis <x; |x; > = Si; It turns out that 8, etc' are all zero: only X + B. Also, H = < X; | H | X; > is tridiagonal. Calc tree ~ M. (~30). (~100)

[M~109 OK!]

Sparsers k= # of steps

(

Exercise: Looky the julia Lanczas method. Mod. fy your general dias method for 5=1/2 chairs to 1) use sparse matrices (julia-buitt-in) and the julia Lances method. How big a system can you do.

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	Product States
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	The simplest type of states are product state Examples: 11 1> = 11> 11> = 11> 11
	点(11)~-14>,)(11>2+日14>2)·点
	In general, if you can write 147 &
	In general, if you can write 147 & 14>= 14>= 14> 14>= 14> 2/3>3 it is a product state
	Entanglement It 14) is not a
	product state, the system is entangled. Usually we consider a system divised into two parts, L (left) and R (rsh)
	Usually we consider a system divised
	L+R disjoint (left) and R(rsh
	L+R disjoint spaces? Then 147 = 142>14R>
	() () () () () () () () () ()
	Suppose ô is an op. acting only on L
	(41014) = < 0,1 < 0,1 0 10, >
	= <\p\ \p\ \p\ \p\ \p\ \p\ \p\ \p\ \p\ \p\
Co	nclusta: Product states describe:
	systems. Independent

Entangled states - how do you tell it a state is enterfel? Example: 2 =pms Let 1A7 = = 1/1/> += 161 18> = シーカインナシノインナシノレクンナシノレム Which is entangled? 18>= た(ハンナル>) た(ハンナル>) un entangled. It is easy to see (A) is entangled (exercise). Show there is no d, B, J, 5 so 14>= (x11>+ p(1>) (x/1) + 5/6>)

To know in general if a 5trte is entangled you need to do a singular value decomposition.

The state of	
-	Singular Value Decomposition (SVD)
	Let M be any complex mxn matrix with nzm. (If n <m, mt<="" on="" svd="" th=""></m,>
	Then there exists $U = m \times m$, $D = m \times m$ with D diagonal, $V = m \times n$ with
	with D diagond, V= mxn with
1	$M = U \cap V / () = () () () ()$
	with Dii 20, U= unitary, V= row-
	(VV+=1) (rows of Vorthonormal).
	This is the SVD. The Di: are the Singular values, unique. Di: = 1;
	Singular values, unique. Dii = 1;
	Another form: D=mxn (0)
- 3	then VIS unitary, nxn
	$M = U \tilde{D} \tilde{V}$
	SVP's have many uses. One i's matrix compression, Suppose only a
(CONSISTED OF	matrix compression. Suppose only a
iller armer i	few Di are non negliste
1	[[[() () () () () () () () (
/	O) [mun] = () (mun) Orop rest of matrices
	and the second s

Schmilt Decomposition
Let 14) = 5 4er 12>1r>
Ildin L, Irdin R (Z) R
Ter is a wavefunction, but treat it as a matrix, do SVA
$4 = \int u \tilde{n} \tilde{v} \tilde{v}$
Normalizata & 17/2-1 = tv & 7 7 7 3 Let
1 ot matrix
Let i \geq = \geq U l \geq
li>R = S Vir Ir>
14)= \(\lambda \) \(\lambda
$= \leq \hat{D}_{ii} i\rangle_{i} i\rangle_{k} = \leq \lambda_{i} i\rangle_{k}$
this is the Schmidt = 17/2 Jecomposith
de compos, th

Normalizeth: $1 = tr \{ y^{t}y \} = tr \{ V + \hat{v} + \hat{v} + \hat{v} + \hat{v} + \hat{v} \}$ $= tr \{ \hat{v} + \hat{v} \} \implies \frac{1}{\sum_{i=1}^{n} 1}$ Di is the probab. 1. ty of the Schmilt-state pair 112/12 If 14) = 10> 15>, it is already in Schmidt Jecomp form, w. H d=1, xis1 =0. 50 +45 tells us it 14) is entangled. Von Neumann Entanglement Entropy 5= - 5 2: In 1? VN ent.
entropy This is the stat mech tormula for entropy usry i: = prob of 1i) 5 measures entanglement. 5=0=>
product 5/4e

Example Two 5=2's IS 147 = 1/2 1117> = 1/2 11 117 entyled? $7 = 7 \left(\frac{1}{52} - \frac{1}{52} \right)$ Need to do SVD (by had) = (10) (10) (or thou on)

U

U

V S=-1h1 -0h0 So 1=1 = 0 Unentagled Product for :190= 17> (= 11) - = 16>