

# Overlap Gap Property for Submatrix Recovery

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# Statistical/Computational Tradeoffs

# Intriguing Phenomena

- In many high-dimensional problems, inference might be information theoretically possible, but computationally intractable!
- Believed to be widespread — e.g. Community Detection, Planted Clique, sPCA, tensor recovery, regression etc.
- Accumulating evidence using diverse approaches:
  - Avg. case reductions
  - State Evolution Analysis
  - Sum of Squares lower bounds and the low-degree likelihood method
  - Query Lower bounds
  - ...

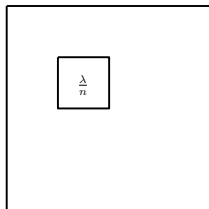
# Planted submatrix recovery

# Setup

- We observe  $M = \mathbb{R}^{n \times n}$ ,  $M = M^T$ .

$$M = \frac{\lambda}{n} \mathbf{v} \mathbf{v}^T + W.$$

- $\mathbf{v} \in \{0, 1\}^n$ ,  $\sum_{i=1}^n v_i = n\rho$ ,  
 $\rho \in (0, 1)$ .
- $W = (W_{ij})$ ,  $\{W_{ij} : i \leq j\}$  ind.  
 $W_{ij} \sim \mathcal{N}(0, \frac{1}{n})$   $i < j$ .

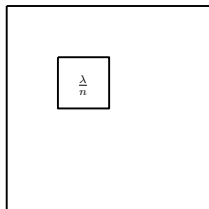


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Today:  $\rho \rightarrow 0$ , following  $n \rightarrow \infty$ .

# Natural Questions

- Can we detect the presence of this planted submatrix?
- Can we recover the submatrix?
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# History of the problem

- Deshpande-Montanari '14 — Bayesian formulation, low rank matrix recovery, AMP optimal for some  $\rho > \rho_c$ .
- Lesieur-Krzakala-Zdeborova '15, '17 — Replica analysis of Mutual Information, State evolution analysis of AMP.

# Likelihood Landscape Analysis

# The MLE

$$\begin{aligned} & \max \langle x, Mx \rangle \\ \text{subject to } & x \in \{0, 1\}^n, \sum_i x_i = n\rho. \end{aligned}$$

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## Definition (Reliable Recovery)

An estimator  $\hat{v} := \hat{v}(M) \in \{0, 1\}^n$  *reliably recovers* the vector  $v$  if there exists  $c > 0$  such that w.h.p.

$$\langle \hat{v}, v \rangle \geq cn\rho.$$

## Theorem (Gamarnik, Jagannath, S. '19)

*If  $\lambda = o\left(\sqrt{\frac{1}{\rho} \log \frac{1}{\rho}}\right)$ , the MLE does not recover the support reliably. On the other hand, if  $\lambda > (2 + \varepsilon)\sqrt{\frac{1}{\rho} \log \frac{1}{\rho}}$  for some  $\varepsilon > 0$ , there exists  $c := c(\varepsilon) > 0$  such that  $\langle \hat{v}_{\text{MLE}}, v \rangle > cn\rho$  w.h.p.*

# A naive Spectral Algorithm

- (i) Let  $\tilde{\mathbf{v}}$  be the top eigenvector of  $M$  with  $\|\tilde{\mathbf{v}}\|_2^2 = n\rho$ . Set  $S = \{i \in [n] : \tilde{v}_i \geq \frac{\delta}{2}\}$ .
- (ii) If  $|S| < n\rho$ , sample  $(n\rho - |S|)$  elements from  $S^c$  and augment to  $S$ . Define the new set  $\tilde{S}$ .
- (iii) If  $|S| > n\rho$ , sample  $n\rho$  elements randomly from  $S$  to construct  $\tilde{S}$ .
- (iv) Set  $\hat{\mathbf{v}} = \mathbf{1}_S$ .

## A naive Spectral Algorithm (cont'd)

### Theorem

*For any  $\varepsilon > 0$  and  $\lambda > (1 + \varepsilon)^{\frac{1}{\rho}}$ , there exists  $\delta := \delta(\varepsilon) > 0$  and  $c := c(\varepsilon, \delta) > 0$  such that w.h.p.  $\langle \hat{v}, v \rangle \geq cn\rho$ .*

What happens for intermediate  $\rho$ ?



# The Restricted Likelihood

$$E_n(q : \rho, \lambda) := \frac{1}{n} \max \langle x, Mx \rangle$$

$$\text{subject to } \langle x, v \rangle \approx nq, x \in \{0, 1\}^n, \sum_i x_i = n\rho.$$

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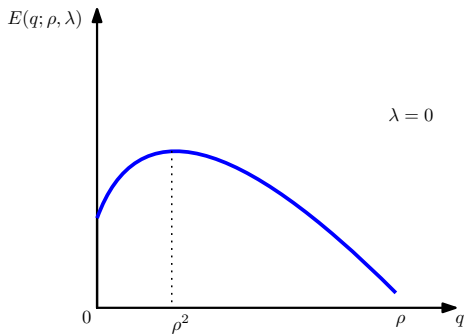
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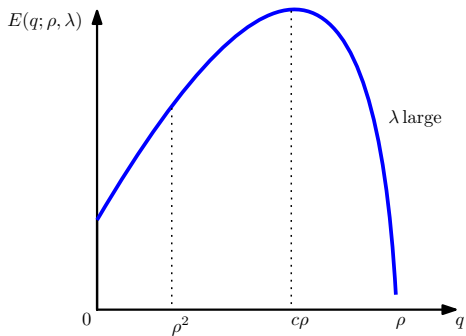
For any  $q \in [0, \rho]$ , there exists  $E(\cdot; \rho, \lambda)$  such that

$$E_n(q; \rho, \lambda) \rightarrow E(q; \rho, \lambda).$$

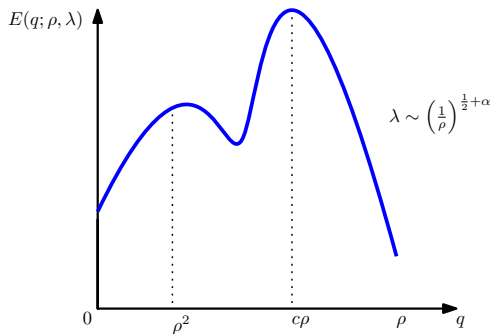
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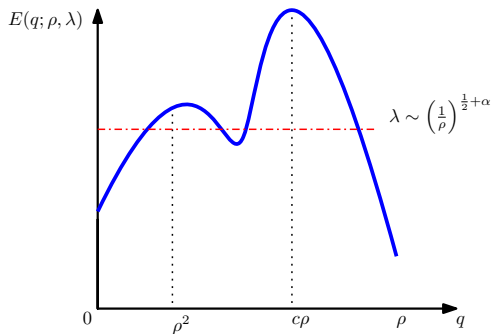
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# Consequences for Markov Chains

- (i) Fix  $\beta > 0$ , and let  $\pi_\beta(x) \propto \exp(\beta \langle x, Mx \rangle)$ .
- (ii) Let  $\{X_t : t \geq 0\}$  be a *local* Markov Chain, reversible with respect to  $\pi_\beta$ .

## Theorem (Gamarnik, Jagannath, S. '19 (informal))

Let  $\lambda \sim \left(\frac{1}{\rho}\right)^{\frac{1}{2}+\alpha}$ ,  $\alpha > 0$  sufficiently small. Define

$$A = \left\{ x \in \{0, 1\}^n : \sum_i x_i = n\rho, an < \sum_i x_i v_i < bn \right\}.$$

For  $\rho > 0$  sufficiently small, there exists  $a < \rho^2 < b = o(\rho)$  such that

$$\int_A Q_x(\tau_{A^c} \leq T) d\pi_\beta(x) \leq T \exp(-\Theta(N)).$$

# Glimpse of Proof

- A Parisi type formula for the restricted likelihood.
- Overlap Gap Property based on variational analysis of ground state formulae for spin glasses.
- Reliable recovery thresholds based on moment based arguments.



# Subsequent Developments

- Barbier, Macris '19, Barbier, Macris, Rush '20 — Mutual information in the regime  $\rho_n \rightarrow 0$ .
- Ben Arous, Wein and Zadik (COLT 2020) — Extended analysis for  $k \times k$  planted matrices with  $k = o(n)$ .

*THANK YOU!*