

Introduction to

Real-time Operating System for Xilinx ZYNQ SoC

freeRTOS & lwIP & FatFS

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Systeme der Elektronik (ZEA-2)

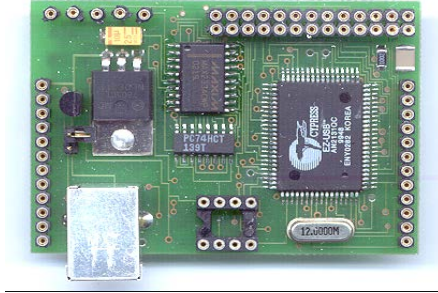
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Motivation

μ Controller

vs.

PC



Embedded System (μ Controller, single Chip Solutions)

- Internal Memory: Flash/RAM (program / data memory)
- Internal I/O
- **„Bare-Metal“ Programmierung**
Without the support of a Operating system
- **Communication**
Serial RS232 (no Ethernet)
- **Data storage**
EEPROMs (limited size, no Filesystem)

PC Processor Systems: (Multiple Chip Solutions)

- External Memories
- External I/O's
- **Operating systems**
Windows, embedded Linux, ...
- **Communication**
Full featured Networking / Ethernet
- **Data storage**
Hard discs, File-System, ...



The gap ...

μController



PC



- Simple „single Chip“ Hardware
- As „Building-Block“ for own HW developments
- **today's μC are powerful** (ARM processor@1GHz)
- **OS Support**
 - Multi-Tasking
 - Real-time
- **Ethernet**
- **Mass storage, Filesystem**

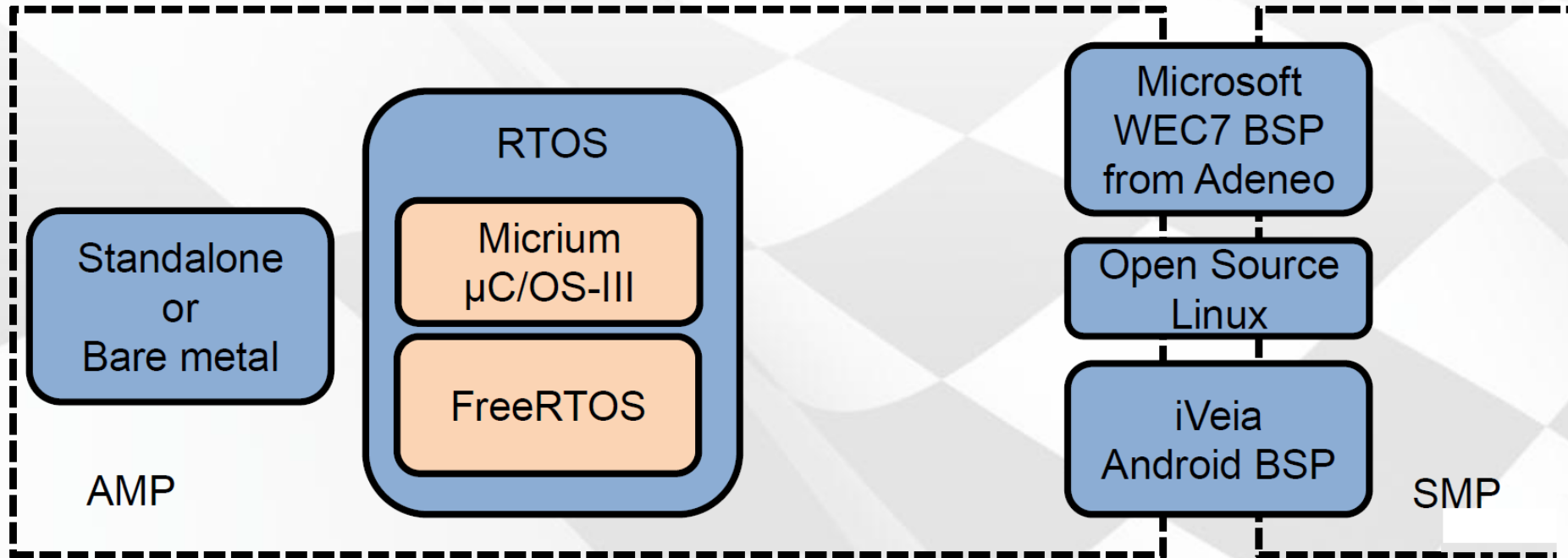
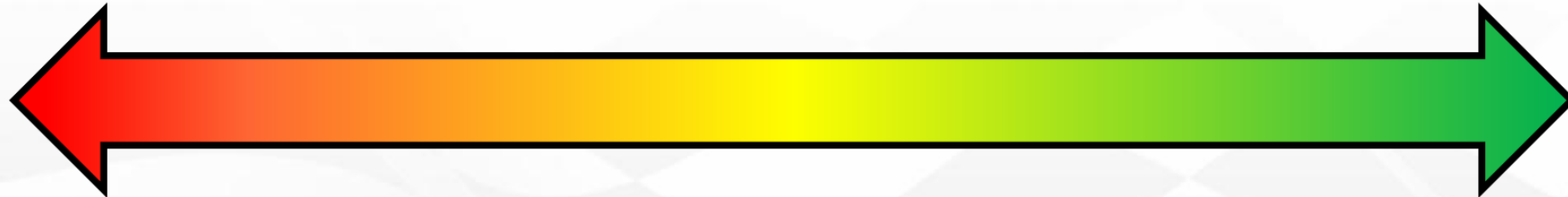
Multi-Tasking / Real-time / Graphical user Interface

What you need ?

- **What are your software application requirements?**

Real-Time Performance

High System Performance



Operating System (OS) Considerations

Bare-Metal / Standalone System

- Software system without an operating system
- Best deterministic behavior (no overhead, fastest interrupt response, ...)
- No support of advanced features (no driver layer, no networking, USB, ...)

→ Minimal complexity

Real-Time Operating Systems

RTOS

- deterministic time behaviour
- predictable response time
- For timing sensitive applications
- Multitasking Support
 - Static Task links,
 - all Task code in image
- Tcp/IP Stacks available

→ Medium complexity

GUI based Operating Systems

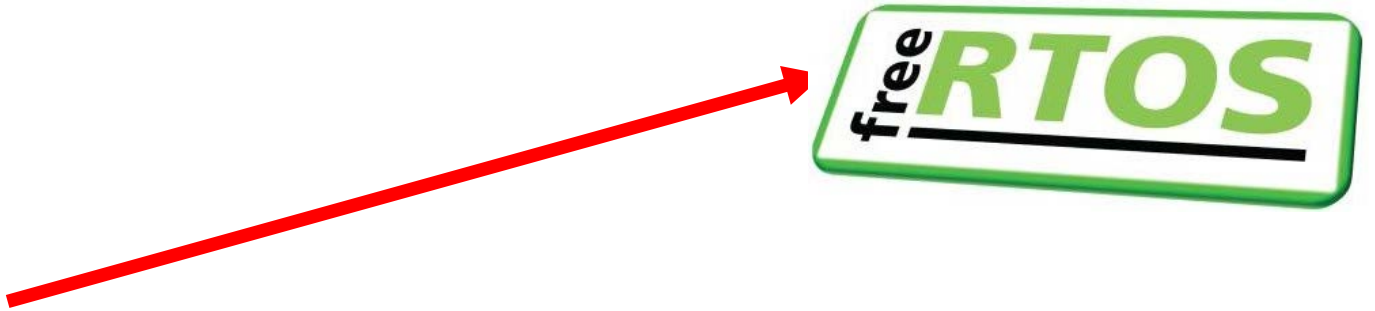
Linux (Windows, ..)

- open-source operating system
- used in many embedded designs
- Full-featured operating system
 - Memory Management Unit (MMU)
 - Full support of all standard interfaces
 - Network, USB, ...
 - and File-System
- not “real” behaviour

→ High complexity

As processing speed has continued to increase for embedded processing, the overhead of an operating system has become mostly negligible in many system designs.

ZYNQ OS support



<https://xilinx-wiki.atlassian.net/wiki/spaces/A/pages/18842118/Zynq-7000+AP+SoC+Operating+Systems>

<https://developer.arm.com/solutions/os>

Solution: freeRTOS, lwIP, FatFS

Operating System

+

Ethernet TCP/IP

+

Filesystem



de.wikipedia.org/wiki/FreeRTOS

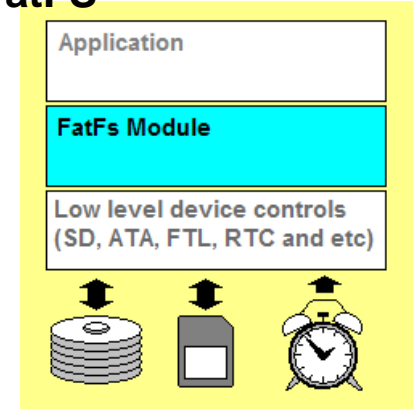
- Open-Source-Echtzeitbetriebssystem
- for embedded MicroControllers
- Multitasking fähig
- präemptive und kooperativer Scheduler



lwip.wikia.com/wiki/LwIP_Wiki

- IP (Internet Protocol)
- ICMP (Internet Control Message Protocol)
- IGMP (Internet Group Management Protocol)
- UDP (User Datagram Protocol)
- TCP (Transmission Control Protocol)
- BSD Berkeley-like socket API
- DNS (Domain names resolver)
- SNMP (Simple Network Management Protocol)
- DHCP (Dynamic Host Configuration Protocol)
- PPP (Point-to-Point Protocol)
- ARP (Address Resolution Protocol) for Ethernet

FatFS



http://elm-chan.org/fsw/ff/00index_e.html



This on

- μ Controller
-

• and **ZYNQ** (ARM μ C)

freeRTOS

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Characteristics of freeRTOS (Operating System)

Born in 2003 and initially conceived for microcontrollers

- Really light
- Really simple: the core of the O.S. are just 3 C files
- Minimal processing overhead
- Ported to a large number of architectures
- Open Source MIT license



- FreeRTOS is a “Embedded Operating System” for
- Embedded MicroController
- Software that provides multitasking facilities.
- FreeRTOS allows to run multiple tasks
- and has a simple scheduler to switch between tasks.

A screenshot of the FreeRTOS website homepage. The browser address bar shows "www.freertos.org". The page features a navigation menu with links like "Quick Start", "Supported MCUs", "PDF Books", "Trace Tools", "Ecosystem", "TCP & FAT", and "Training". The main content area includes the FreeRTOS logo, a navigation menu, and a central banner with the text "FreeRTOS™" and "The Market Leading, De-facto Standard and Cross Platform Real Time Operating System (RTOS). Don't Let Your RTOS Lock You In." Below this, there are several paragraphs of text describing the system's features and benefits, along with a list of links for more information.

FreeRTOS features:

- Priority-based multitasking capability
- Queues to communicate between multiple tasks
- Semaphores to manage resource sharing between tasks
- Utilities to view CPU utilization, stack utilization etc.

More information at: www.FreeRTOS.org

Supported CPUs (Ports): http://www.freertos.org/RTOS_ports.html

Programming styles

Standalone/Baremetal: Based on one superloop

```
void main ()
{
    Init_all();
    while (1)
    {
        do_A();
        do_B();
        do_C();
    }
}
```

freeRTOS:

```
void main()
{
    xTaskCreate (Task_A, ...);
    xTaskCreate (Task_B, ...);
    xTaskCreate (Task_C, ...);
    xTaskStartScheduler ();
}
```

The main program only initializes the needed tasks and starts the scheduler. After this the tasks (in this example 3 Tasks) are now working in parallel. Each Task can have his own initializing part. Finally each tasks operates in a own while loop, given the feeling of having several main programs in parallel.

```
void Task_A (void *p)
{
    Init_A();
    while (1)
    {
        do_A();
    }
}
```

```
void Task_B (void *p)
{
    Init_B();
    while (1)
    {
        do_B();
    }
}
```

```
void Task_C (void *p)
{
    Init_C();
    while (1)
    {
        do_C();
    }
}
```

FreeRTOS Configuration

through a header file: FreeRTOSConfig.h

(<http://www.freertos.org/a00110.html>)

```
#define configUSE_PREEMPTION 1
#define configUSE_MUTEXES 1
#define INCLUDE_xSemaphoreGetMutexHolder 1
#define configUSE_RECURSIVE_MUTEXES 1
#define configUSE_COUNTING_SEMAPHORES 1
#define configUSE_TIMERS 1
#define configUSE_IDLE_HOOK 0
#define configUSE_TICK_HOOK 0
#define configUSE_DAEMON_TASK_STARTUP_HOOK 0
#define configUSE_TICKLESS_IDLE 0
#define configTASK_RETURN_ADDRESS NULL
#define INCLUDE_vTaskPrioritySet 1
#define INCLUDE_uxTaskPriorityGet 1
#define INCLUDE_vTaskDelete 1
#define INCLUDE_vTaskCleanUpResources 1
#define INCLUDE_vTaskSuspend 1
#define INCLUDE_vTaskDelayUntil 1
#define INCLUDE_vTaskDelay 1
#define INCLUDE_eTaskGetState 1
#define INCLUDE_xTimerPendFunctionCall 1
#define INCLUDE_pcTaskGetTaskName 1
#define configMAX_API_CALL_INTERRUPT_PRIORITY (18)
```

Tasks can be interrupted by others with higher priority

This will include a timer service task

Hooks are used to trigger the execution of functions upon the happening of certain events

Some functionality can be optionally included/excluded from the core of the O.S.

by the mss file in the FreeRTOS BSP generated in the SDK

Board Support Package Settings

Control various settings of your Board Support Package.

Configuration for OS: freertos10_xilinx

Name	Value	Default	Type	Description
SYSINTC_SPEC	*			
SYSTMTR_DEV	*			
SYSTMTR_SPEC	true			
stdin	ps7_uart_1	none	peripheral	stdin peripheral
stdout	ps7_uart_1	none	peripheral	stdout peripheral
enable_stm_event_trace	false	false	boolean	Enable event tracing through System Trace M
hook_functions	true	true	boolean	Include or exclude application defined hook (c
use_daemon_task_sta	false	false	boolean	Set true for kernel to call vApplicationDaemoi
use_idle_hook	false	false	boolean	Set to true for the kernel to call vApplicationic
use_malloc_failed_hoc	true	true	boolean	Only used if a FreeRTOS memory manager (h
use_tick_hook	false	false	boolean	Set to true for the kernel to call vApplicationT
kernel_behavior	true	true	boolean	Parameters relating to the kernel behavior
kernel_features	true	true	boolean	Include or exclude kernel features
software_timers	true	true	boolean	Options relating to the software timers functi
tick_setup	true	true	boolean	Configuration for enabling tick timer

Cancel OK

A Task

- Task's are parallel operating MAIN routines
- is a simple C function
- a pointer to parameters (void *) as input
- Creates a forever loop (while (1))
- The tasks are controlled by the Scheduler (freeRTOS internal function)

A task can be preempted (swapped out)

- because of a more priority task
- because it has been delayed (call to vTaskDelay())
- because it waits for a event (semaphore, ...)

When a task can run

- state is set "Ready"

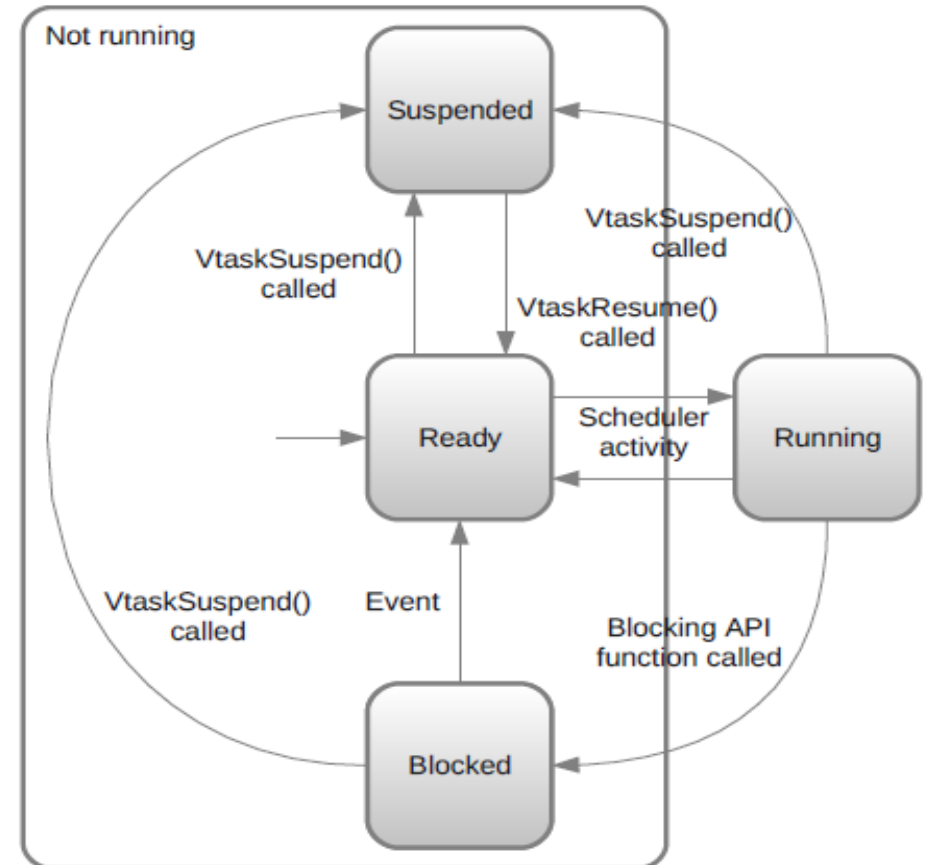
Task will start (swapped in) when

- No more priority task running at this time
- No Delay or Blocking condition

Finally

- a call to vTaskSuspend() stops the Task at all
- vTaskResume() brings him back to the scheduler)

```
void my_task(void* p)
{
  while (1)
  {
    .....
  }
}
```



Creating a Task

The Task function itself:

```
void task_MySampleTask( void *p )
{
int xx[32];
    // do initialization
    while (1)
    {
        // Task execution code
    }
}
```

Remark for Stacksize:

- Each task has his own Stack
- Local variables and function calls are using stack
- 'printf', consumes around 1024 bytes of stack
- Running out of Stack will crash freeRTOS (freeze)
- **At minimum allocate min. 1024 bytes for Stack**

Install the Task (in main.c):

```
portBASE_TYPE xTaskCreate (
    pdTASK_CODE    pvTaskCode,           // pointer to the Task
    char*          pcName,               // String: name of Task
    unsigned short usStackSize,         // Stacksize
    void *         pvParameters,        // pointer to Parameters (can be NULL)
    unsigned short uxPriority,           // Priority (default 2)
    xTaskHandle*   pxCreatedTask );     // Pointer to receive Task handle (can be NULL)
    ???
```

Simple Main and Task

- Simple task example that prints a message once a second.
- **vTaskStartScheduler()** never returns and FreeRTOS will begin servicing the tasks at this point.
- Every task must have an **infinite loop and NEVER EXIT**.

```
void task_HelloWorld (void* p)
{
    while(1)
    {
        printf("Hello World!");
        vTaskDelay(1000);
    }
}

void main(void)
{
    xTaskCreate (task_HelloWorld, „Hello", 1024, NULL, 1, NULL);

    vTaskStartScheduler();
    // never comes here
}
```

Controlling Tasks

- Task with highest priority will run first, and never give up the CPU until it sleeps
- If two or more tasks with the same priority do not give up the CPU (they don't sleep), then FreeRTOS will share the CPU between them (time slice).
- Here are some of the ways you can give up the CPU:
 - **vTaskDelay (..)** This simply puts the task to "sleep,,,. You decide how much you want to sleep.
 - **vTaskDelayUntil (..)** Delay a task until a specified absolute time.
 - **xQueueSend (..)** If the Queue you are sending to is full, this task will sleep (block).
 - **xQueueReceive (..)** If the Queue you are reading from is empty, this task will sleep (block).
 - **xSemaphoreTake (..)** You will sleep if the semaphore is taken by somebody else.

A typical freeRTOS application

```
int main ( void )
{
    // do needed Platform initialization

// Now we deal with RTOS. Create the Tasks and start the scheduler

// 1) Start LED 1 toggle
    xTaskCreate (Task_LEDS, (signed char*) "LEDs", 64, NULL, 1, NULL);

// 2) Start SWITCH
    xTaskCreate (Task_SWITCH, (signed char*) "Sws", 64, NULL, 1, NULL);

// 3) Start LCD-Anzeige
    xTaskCreate (Task_LCD, (signed char*) "LCD", 1024, NULL, 1, NULL);

// Finally: Start FreeRTOS
    vTaskStartScheduler();

// Will only reach here if there was insufficient memory to create the idle task
    return 0;
}
```


Stop a Task / Scheduler

Task:

- A Task can stop himself or also other tasks

```
// Example: Stopping a task
...
vTaskDelete (NULL);
```

Scheduler:

- one method to create a critical section:
 - prevent a task from preempting it
 - but let interrupts to do their job

→ Stopping the scheduler (= stopping all other tasks)

→ Do code in 'critical section'

→ Restart the scheduler

Notice: No FreeRTOS API functions can be called when the scheduler is stopped !

```
// Example: Create a critical section
...
vTaskSuspendAll();
{
    printf( "%s", pcString );
    fflush( stdout );
}
xTaskResumeAll();
...
```

Software Timers

```
TimerHandle_t xTimerCreate (
    const char * const pcTimerName,           // name of the timer
    const TickType_t xTimerPeriod,           // Period in ticks (ms)
    const UBaseType_t uxAutoReload,          // FALSE=OneShoot / TRUE=Repeat
    void * const pvTimerID,                  // ID number
    TimerCallbackFunction_t pxCallbackFunction ); // Timer Callback function
```

```
#include <timers.h>

TimerHandle_t xTimer;

void vTimerCallback ( TimerHandle_t pxTimer ) // Timer callback function
{
    static cnt = 0;
    cnt++;
    if (cnt > 10) xTimerStop( pxTimer, 0 );
}

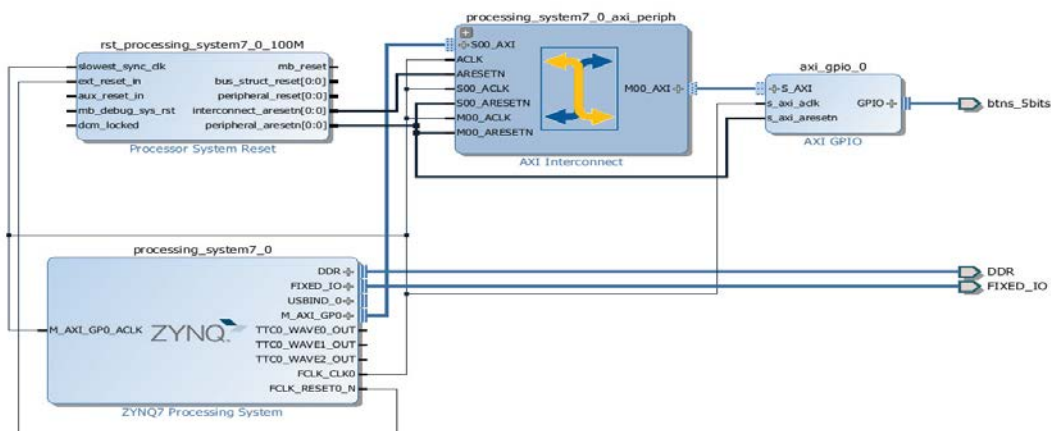
void main( void )
{
    xTimer = xTimerCreate ("Tim1", 200, TRUE, (void*)1, vTimerCallback ); // Install the timer
    xTimerStart (xTimer, 0 ); // Start timer

    ... // do more (Create other tasks ....)
    vTaskStartScheduler();
}
```

Project 1: First (simple) FreeRTOS application

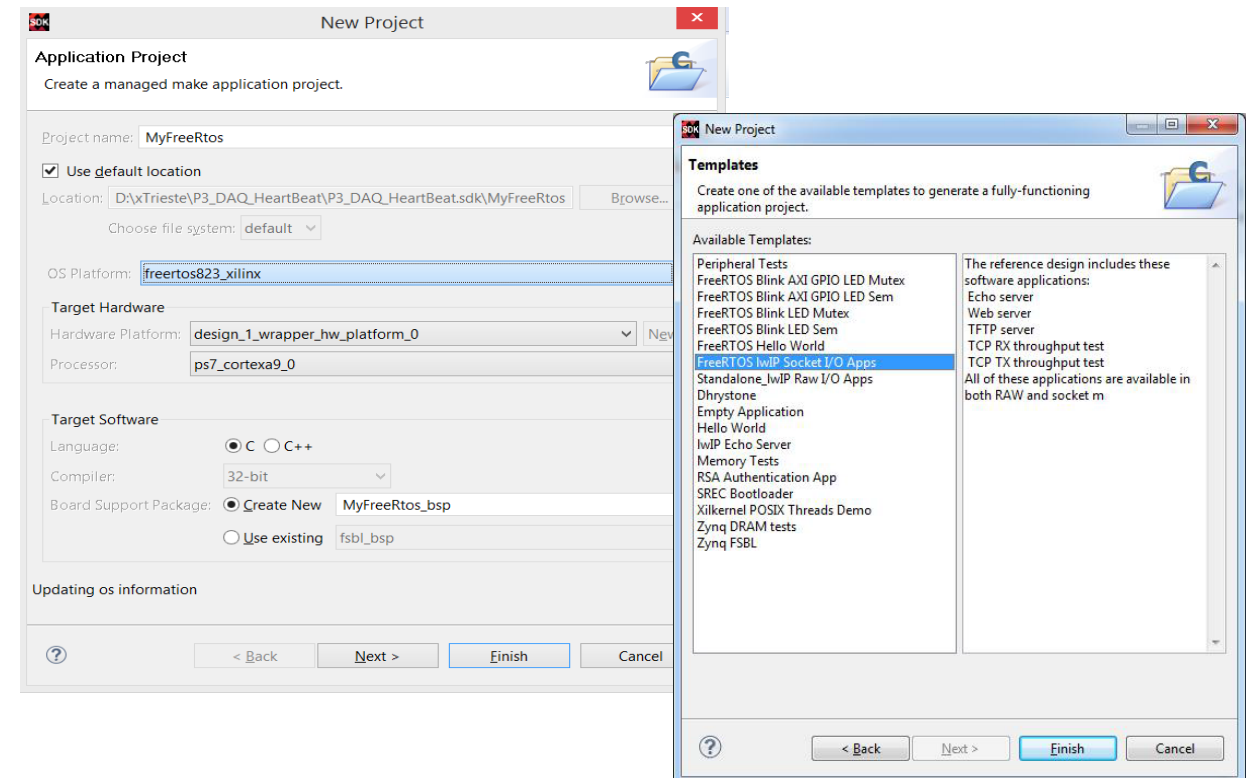
1) VIVADO Block Design

- Configure a Vivado project with the IPI (IP Block Integrator)
- Add the “ZYNQ7 Processing System”
 - (use the standard MIO Signals (PS Multiplexed IO))
- Add a GPIO to the LEDs
- Make a HDL Wrapper,
- Implement the Design,
- Generate Bitstream,
- Export to SDK and launch SDK



2) SDK Application

- File -> New -> Application Project
- Provide Project Name
- Target Hardware: Choose the Hardware Platform
- **OS Platform: freertosXXX_Xilinx**
- Click Next
- For first example: “FreeRTOS Hello World”
- Click Finish



Project 1: simple freeRTOS example

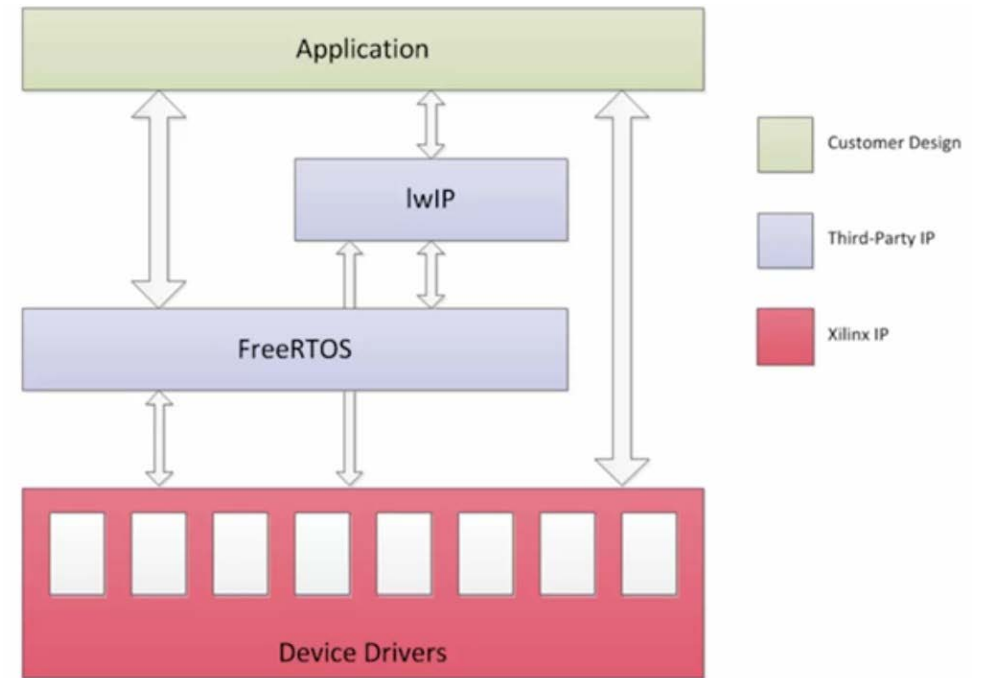
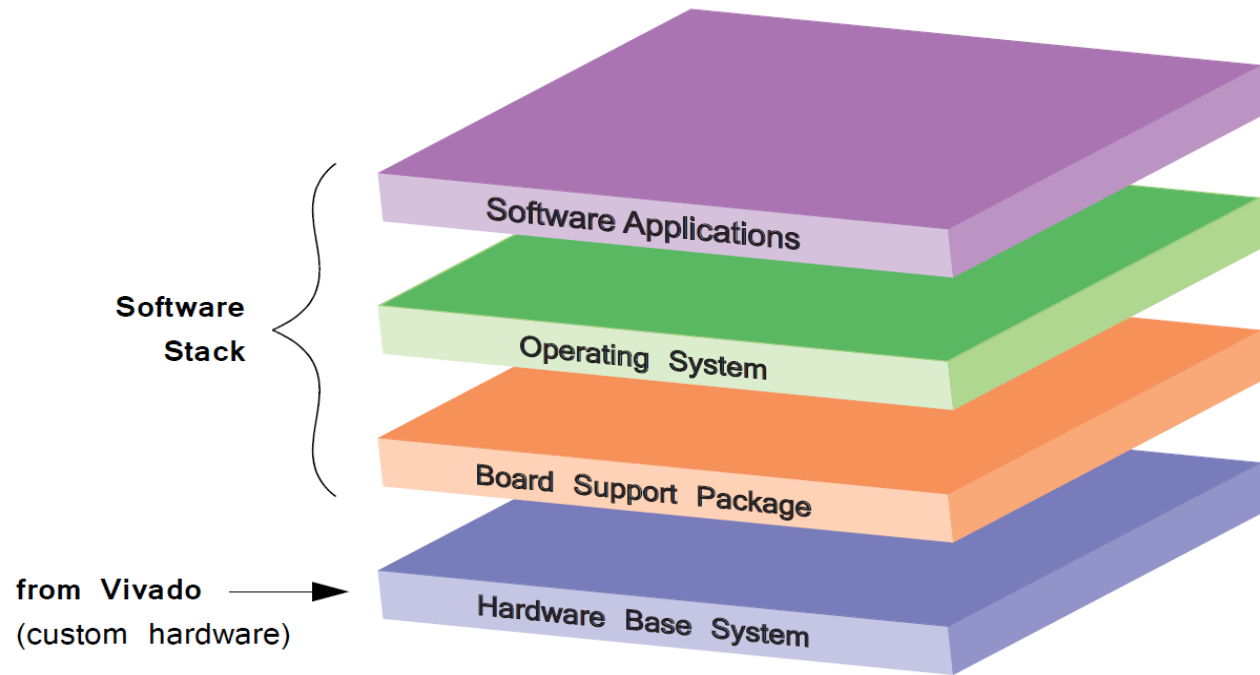
```
#include ...  
int tick=0;
```

```
void Task_LED (void* p)  
{  
    while (1)  
    {  
        Xil_Out32 (aGPIO, tick);  
        vTaskDelay (100);  
    }  
}
```

```
void Task_Print (void* p)  
{  
    while (1)  
    {  
        printf („Tick is %d \n“, tick);  
        vTaskDelay (500);  
        tick++;  
    }  
}
```

```
int main ( void )  
{  
    // 1) Start LED 1 toogle  
    xTaskCreate (Task_LED, (signed char*) "LEDs", 1024, NULL, 1, NULL);  
  
    // 2) printf  
    xTaskCreate (Task_Print, (signed char*) "Print", 1024, NULL, 1, NULL);  
  
    // Finally: Start FreeRTOS  
    vTaskStartScheduler();  
  
    // Will only reach here if there was insufficient memory to create the idle task  
    return 0;  
}
```

Appendix: Software Stack



lwIP

TCP/IP Ethernet-stack for freeRTOS

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The TCP/IP Stack

lwIP stands for 'Lightweight IP'

- full scale TCP protocol stack
- small memory footprint (for embedded systems, μ C)
- Open Source (C Code)

Supports a large number of protocols and APIs

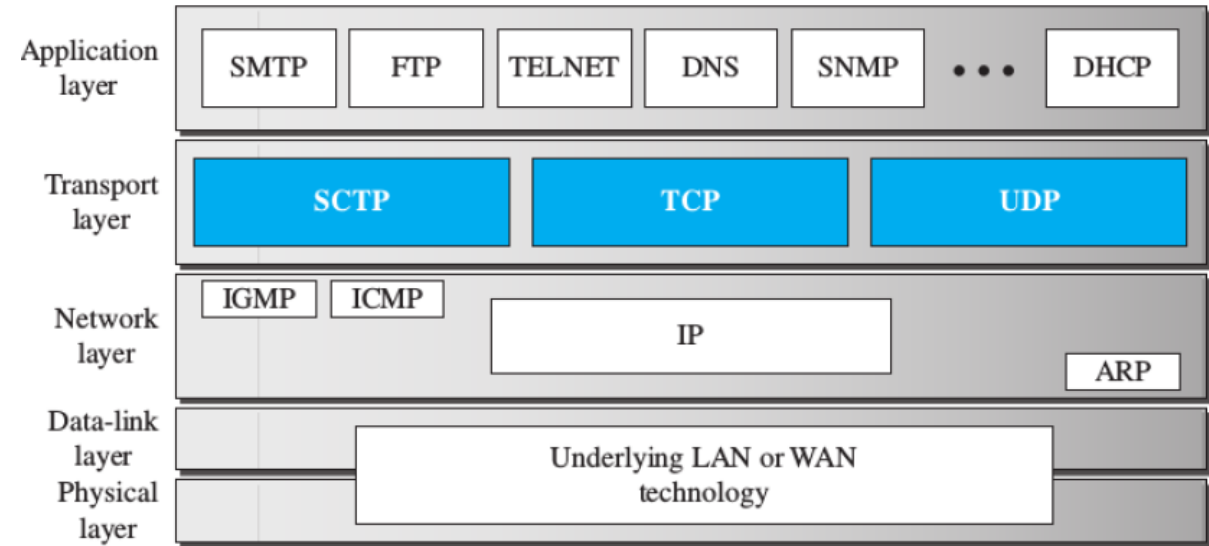
- TCP Transport Control Protocol
- UDP User Datagram Protocol
- IP Internet Protocol
- ICMP Internet Control Message Protocol
- ARP Address Resolution Protocol
- DHCP Dynamic Host Configuration Protocol
- Raw API and Berkeley sockets (requires an multitasking O.S.)

Included in Xilinx SDK

- includes driver for Xilinx Ethernet driver
- XAPP1026 is the reference application note

Application level

- HTTP(S) server, SMTP client, SMTP(S) client, ping, TFTP, ...



The network design is organized as a layer stack.

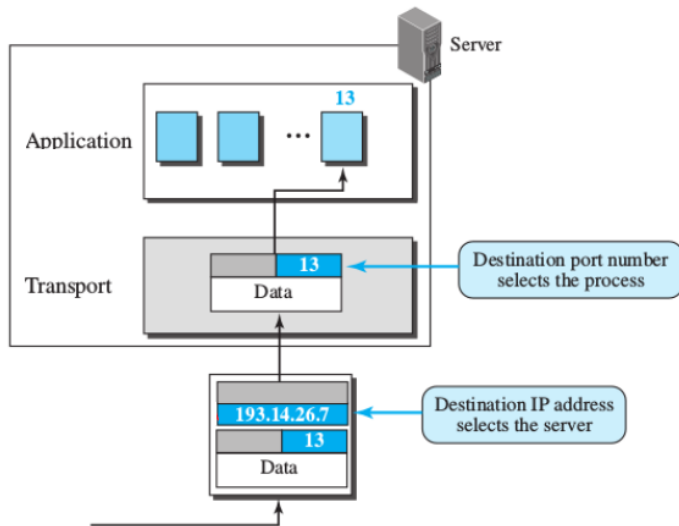
- Each layer provides a set of services to the upper layer and requires services from the lower layer.

BSD Sockets

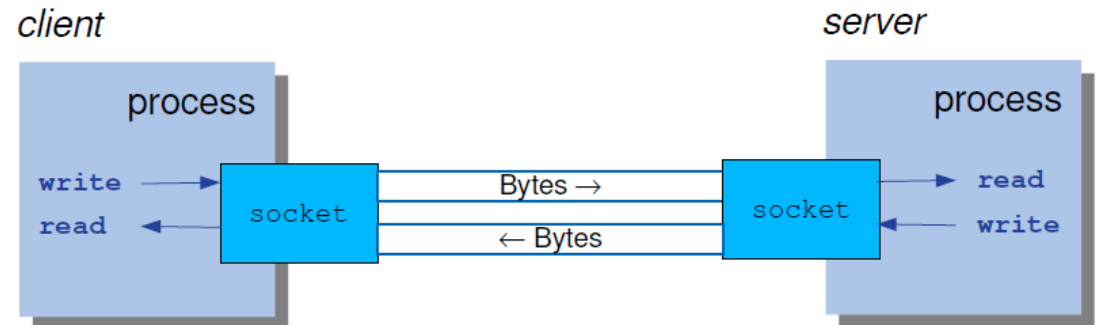
BSD Sockets (Berkeley sockets | POSIX sockets)

- de facto standard API
- Basic abstraction for network programming
- Combination of **IP address + port**
- Inter-process communication

→ use „LwIP Socket API“



```
lwip_socket(AF_INET, SOCK_STREAM, 0)
```



SETUP & STARTUP OF THE NETWORK

Basic template for freeRTOS & lwIP:

➤ Start a “*task_StartEthernet*”

- Initializes lwip
- Configures a network interface
- Start the interface and a reception task
- Installs any other network tasks
 - Example: Web-Server, Echo-server
- Finally the start up task deletes himself

After initialization several threads are active:

- Reception task
- Web-server
- Echo-server

```
int main()
{
    printf ("Start Ethernet \n\r");
    xTaskCreate (task_StartEthernet, (char*)"Start_Eth", 2048, NULL, DEFAULT_THREAD_Prio, NULL);

    vTaskStartScheduler();
    while(1);
    return 0;
}
```

```
void task_StartEthernet(void *p)
{
    struct netif *netif;
    struct ip_addr ipaddr, netmask, gw;

    /* the mac address of the board, this should be unique per board */
    unsigned char mac_ethernet_address[] = { 0x00, 0x0a, 0x35, 0x00, 0x01, 0x02 };

    lwip_init();
    printf ("Start Ethernet2 \n\r");

    netif = &server_netif;

    /* initialize IP addresses to be used */
    IP4_ADDR(&ipaddr, 192, 168, 2, 2);
    IP4_ADDR(&netmask, 255, 255, 255, 0);
    IP4_ADDR(&gw, 192, 168, 2, 1);

    /* print out IP settings of the board */
    print("\r\n\r\n");
    print("-----lwIP Socket Mode Demo Application -----r\n");

    print_ip_settings(&ipaddr, &netmask, &gw);
    /* print all application headers */

    /* Add network interface to the netif_list, and set it as default */
    if (!xemac_add(netif, &ipaddr, &netmask, &gw, mac_ethernet_address, PLATFORM_EMAC_BASEADDR))
    {
        xil_printf("Error adding N/W interface\r\n");
        return;
    }
    netif_set_default(netif);
    netif_set_up(netif); // specify that the network if is up

    /* start packet receive thread - required for lwIP operation
    sys_thread_new ("xemacif_input_thread", (void*)(void*)xemacif_input_thread, netif, 2048, DEFAULT_THREAD_Prio);

    //*****

    sys_thread_new("httpd", web_application_thread, 0, 2048, DEFAULT_THREAD_Prio);

    sys_thread_new("echod", echo_application_thread, 0, 2048, DEFAULT_THREAD_Prio);

    vTaskDelete(NULL);
    return;
}
```

UDP

Unreliable protocol

- No error control
- corrupted packets are ignored
- No flow control (Speed)

But:

- Extremely simple (minimum overhead)
- the fastest way (lowest latency)

• UDP socket Programming flow

```
void echo_application_thread()
{
    int sock, new_sd;
    struct sockaddr_in address, remote;
    int size;

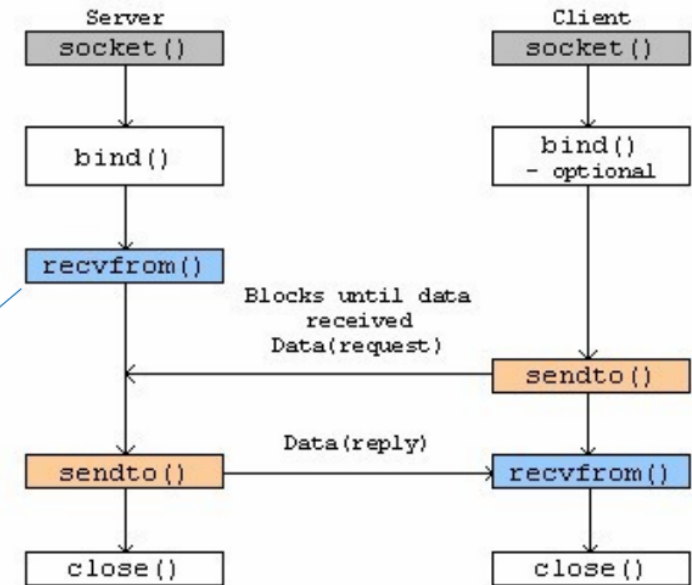
    int RECV_BUF_SIZE = 2048;
    char recv_buf[RECV_BUF_SIZE];
    int n, nwrote;

    if ((sock = lwip_socket(AF_INET, SOCK_DGRAM, 0)) < 0)
        return;

    address.sin_family = AF_INET;
    address.sin_port = htons(echo_port);
    address.sin_addr.s_addr = INADDR_ANY;

    if (lwip_bind(sock, (struct sockaddr *)&address, sizeof (address)) < 0)
        return;

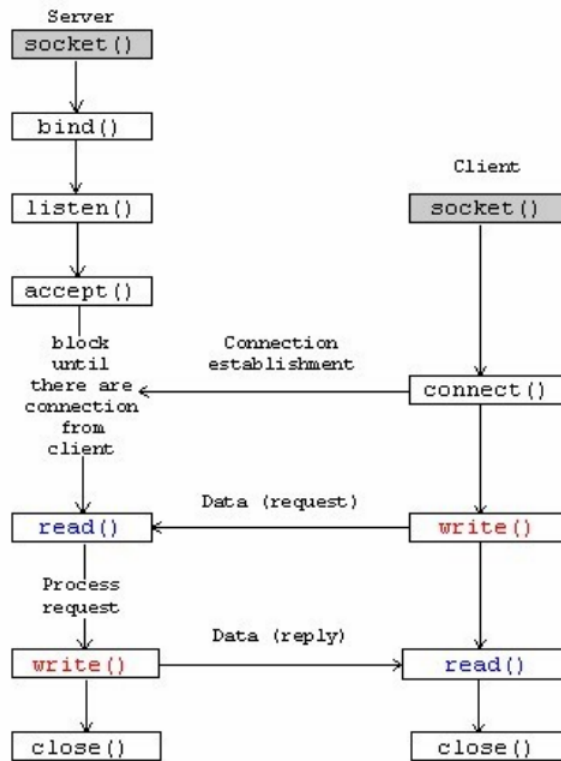
    if ((n = read(sock, recv_buf, RECV_BUF_SIZE)) < 0) {
        xil_printf("%s: error reading from socket %d, closing socket\r\n",
            __FUNCTION__, sock);
    }
}
```



Used in applications where losing of some part of the information can be tolerated, example Video Streaming/conference

TCP

- Connection-oriented protocol
- Reliable, Error free (correction)
 - Retransmission of lost or corrupted packets
- Complex protocol with multiple phases
 - higher latency, lower throughput
 - Connection control



Used when losing information can't be tolerated.
Example: HTTP, E-mail, binary Data, ...

```
void echo_application_thread()
{
    int sock, new_sd;
    struct sockaddr_in address, remote;
    int size;

    if ((sock = lwip_socket(AF_INET, SOCK_STREAM, 0)) < 0)
        return;

    address.sin_family = AF_INET;
    address.sin_port = htons(echo_port);
    address.sin_addr.s_addr = INADDR_ANY;

    if (lwip_bind(sock, (struct sockaddr *)&address, sizeof (address)) < 0)
        return;

    lwip_listen(sock, 0);

    size = sizeof(remote);

    while (1) {
        if ((new_sd = lwip_accept(sock, (struct sockaddr *)&remote, (socklen_t *)&size)) > 0) {
            sys_thread_new("echos", process_echo_request,
                (void *)new_sd,
                THREAD_STACKSIZE,
                DEFAULT_THREAD_PRIORITY);
        }
    }
}

/* thread spawned for each connection */
void process_echo_request(void *p)
{
    int sd = (int)p;
    int RECV_BUF_SIZE = 2048;
    char recv_buf[RECV_BUF_SIZE];
    int n, nwrote;

    while (1) {
        /* read a max of RECV_BUF_SIZE bytes from socket */
        if ((n = read(sd, recv_buf, RECV_BUF_SIZE)) < 0) {
            xil_printf("%s: error reading from socket %d, closing socket\r\n", __FUNCTION__, sd);
            break;
        }

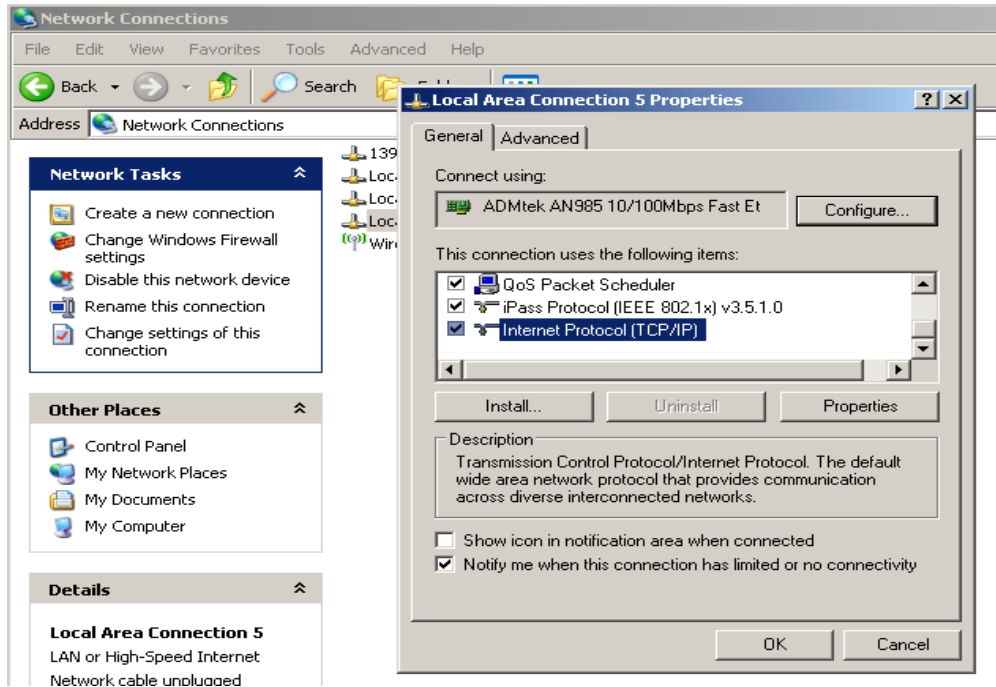
        /* break if client closed connection */
        if (n <= 0)
            break;

        /* handle request */
        if ((nwrote = write(sd, recv_buf, n)) < 0) {
            xil_printf("%s: ERROR responding to client echo request. received = %d, written = %d\r\n",
                __FUNCTION__, n, nwrote);
            xil_printf("Closing socket %d\r\n", sd);
            break;
        }
    }

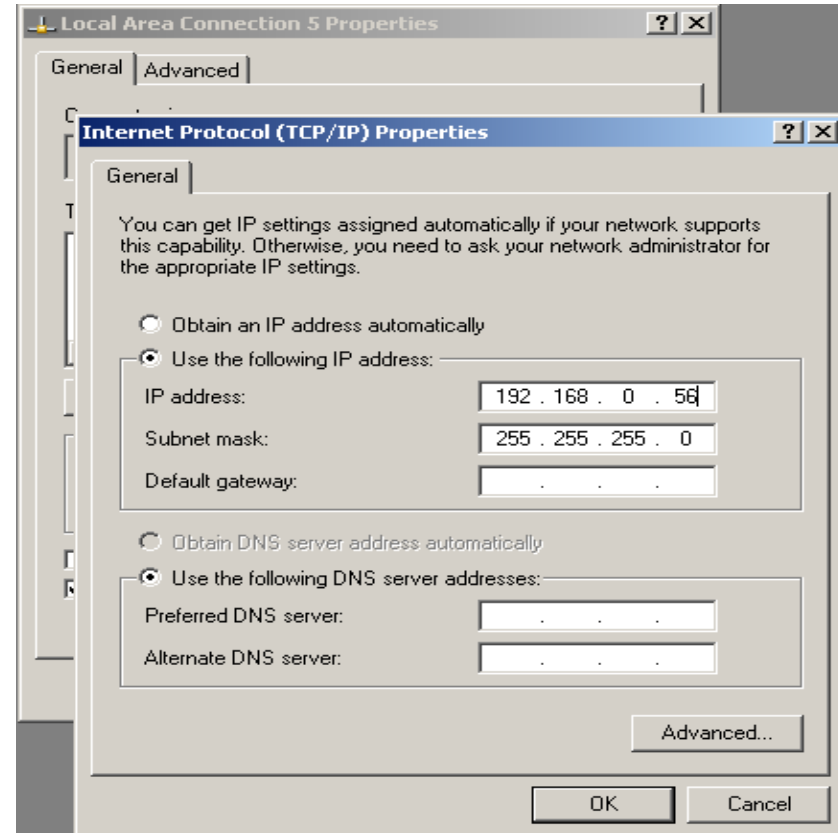
    /* close connection */
    close(sd);
    vTaskDelete(NULL);
}
```

Setup PC for communication

Control Panel → Network connections → Properties



Set PC TCP/IP to address within same subnet:



TROUBLESHOOTING

Wireshark

- widely-used network protocol analyzer.
- see what's happens on your network
- Filters to select specific packets
- Live capture / Offline analysis
- Multi-Protocol
- Decryption of packets headers

A screenshot of the Wireshark network protocol analyzer interface. The main pane displays a list of captured packets with columns for No., Time, Source, Destination, Protocol, and Info. The bottom pane shows the detailed view of a selected packet, including Ethernet II, Internet Protocol, User Datagram Protocol, and NetBIOS Name Service layers. The status bar at the bottom indicates 42 packets displayed and 0 dropped.

No.	Time	Source	Destination	Protocol	Info
1	0.000000	169.254.94.55	169.254.255.255	NBNS	Registration NB WORKGROUP<1e>
2	0.764428	169.254.94.55	169.254.255.255	NBNS	Registration NB WORKGROUP<1e>
3	1.027526	161.67.27.162	255.255.255.255	UDP	Source port: 17500 Destination port: 17500
4	1.030688	161.67.27.162	255.255.255.255	UDP	Source port: 17500 Destination port: 17500
5	1.030834	161.67.27.162	255.255.255.255	UDP	Source port: 17500 Destination port: 17500
6	1.030985	161.67.27.162	161.67.27.255	UDP	Source port: 17500 Destination port: 17500
7	1.128790	161.67.27.141	161.67.27.255	CUPS	ipp://161.67.27.141:631/printers/Impresora_SALA_2 (idle)
8	1.235153	161.67.27.136	161.67.27.255	UDP	Source port: 17500 Destination port: 17500
9	1.377818	161.67.27.183	161.67.27.255	NBNS	Name query NB GRUPO_TRABAJO<1b>
10	1.590410	Cisco_49:33:17	Spanning-tree-(for-br	STP	Conf. Root = 32768/0/00:02:b9:1e:2e:40 Cost = 27 Port = 0x8017
11	1.622479	fe80::d52d:6cea:2ee9:	ff02::1:2	DHCPv6	Solicit
12	2.127825	161.67.27.183	161.67.27.255	NBNS	Name query NB GRUPO_TRABAJO<1b>
13	2.778209	209.85.227.125	161.67.27.210	Jabber/X	Response: \027\003\001\001 R\271\365z\340\260\266\324\341\v\365\366
14	2.877821	161.67.27.183	161.67.27.255	NBNS	Name query NB GRUPO_TRABAJO<1b>
15	2.878097	161.67.27.141	161.67.27.255	BROWSER	Local Master Announcement PIKE, Workstation, Server, Print Queue Serv
16	2.878128	161.67.27.141	161.67.27.255	BROWSER	Domain/Workgroup Announcement WORKGROUP, NT Workstation, Domain Enum
17	2.986303	fe80::55c3:ebb0:6c0c:	ff02::1:2	DHCPv6	Solicit
18	3.041701	Dell_b5:59:a6	Broadcast	ARP	who has 161.67.27.86? Tell 161.67.27.150
19	3.532048	161.67.27.1	224.0.0.5	OSPF	Hello Packet
20	3.604098	Cisco_49:33:17	Spanning-tree-(for-br	STP	Conf. Root = 32768/0/00:02:b9:1e:2e:40 Cost = 27 Port = 0x8017
21	3.815277	161.67.27.137	161.67.27.255	NBNS	Name query NB IMAC-MARD<00>
22	3.817904	161.67.27.137	161.67.27.255	NBNS	Name query NB IMP-ESI3<00>
23	3.818279	Vmware_b3:69:88	Broadcast	ARP	who has 161.67.27.137? Tell 161.67.27.129
24	3.820521	161.67.27.137	161.67.27.255	NBNS	Name query NB IMP-ESI2<00>

Packet details for Frame 1 (110 bytes on wire, 110 bytes captured):

- Ethernet II, Src: Dell_b5:59:a6 (00:21:70:b6:07:ee), Dst: Broadcast (ff:ff:ff:ff:ff:ff)
- Internet Protocol, Src: 169.254.94.55 (169.254.94.55), Dst: 169.254.255.255 (169.254.255.255)
- User Datagram Protocol, Src Port: netbios-ns (137), Dst Port: netbios-ns (137)
- NetBIOS Name Service

Packet contents (hex dump):

```
0000 ff ff ff ff ff 00 21 70 b6 07 ee 08 00 45 00 .....! p.....E.
0010 00 60 00 30 00 00 80 11 88 29 a9 fe 5e 37 a9 fe ..0....).^7..
0020 ff ff 00 89 00 89 00 4c e1 80 82 c2 29 10 00 01 .....L.....)
0030 00 00 00 00 00 01 20 46 48 45 50 46 43 45 4c 45 .....F HEPPCBL
0040 48 46 43 45 50 46 46 46 41 43 41 43 41 43 41 43 HFCEPFFF ACACACAC
0050 41 43 41 43 41 42 4f 00 00 20 00 01 c0 0c 00 20 ACACABD. ....
0060 00 01 00 04 93 e0 00 06 e0 00 a9 fe 5e 37 .....^7
```

Captured packets

Packet encapsulation

Packet contents

FatFS

http://elm-chan.org/fsw/ff/00index_e.html

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FAT File System

FatFs - Generic FAT Filesystem Module

FatFs is a generic FAT/exFAT filesystem module for small embedded systems. The FatFs module is written in compliance with ANSI C (C89) and completely separated from the disk I/O layer. Therefore it is independent of the platform. It can be incorporated into small microcontrollers with limited resource, such as 8051, PIC, AVR, ARM, Z80, RX and etc. Also Petit FatFs module for tiny microcontrollers is available [here](#).

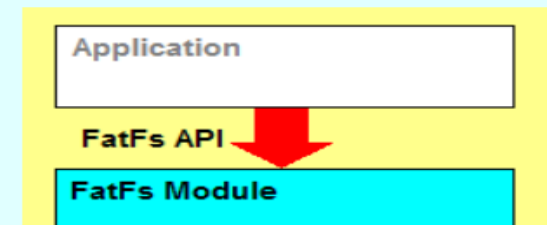
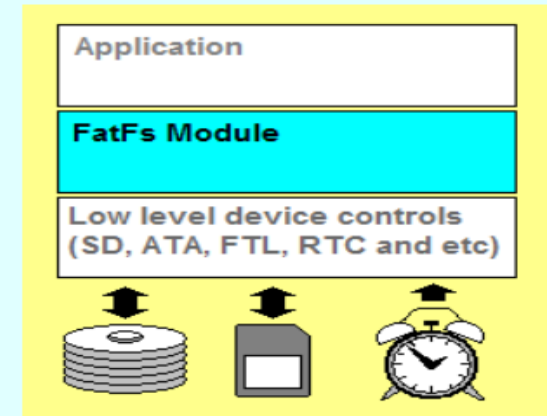
Features

- DOS/Windows compatible FAT/exFAT filesystem.
- Platform independent. Easy to port.
- Very small footprint for program code and work area.
- Various [configuration options](#) to support for:
 - Multiple volumes (physical drives and partitions).
 - Multiple code pages including DBCS.
 - Long file name in OEM code or Unicode.
 - exFAT filesystem.
 - Thread safe.
 - Fixed or variable sector size.
 - Read-only, optional API, I/O buffer and etc...

Application Interface

FatFs provides various filesystem functions for the applications as shown below.

- File Access
 - [f_open](#) - Open/Create a file
 - [f_close](#) - Close an open file
 - [f_read](#) - Read data from the file
 - [f_write](#) - Write data to the file



FatFS programming example

```
//-- Fat-File System
#include "ff.h"

FATFS      fatfs;
TCHAR      *Path = "0:/";

int SD_Mount (void)
{
    FRESULT Res;
    Res = f_mount (&fatfs, Path,1);
    if (Res != FR_OK)
    {
        printf ("SD: Mount failed\n\r");
        return 0;
    }
    printf ("SD mounted\n\r");
    return 1;
}

int SD_Unmount(void)
{
    f_mount(NULL, Path, 1);
    SDMounted = 0;
    return 1;
}
```

```
void Write_SDcard (void)
{
    int Res, wr;
    FIL* hFile;
    Res = f_open (&hFile, "data.bin", FA_CREATE_ALWAYS | FA_WRITE );
    if (Res)
    {
        printf("SD: Open failed \n\r");
        return 0;
    }
    Res = f_write (&hFile, Data, 1024*2, &wr);
    Res = f_close (&hFile);
    return 1;
}
```

```
void Read_SDcard (void)
{
    int Res, rd;
    FIL* hFile;
    Res = f_open (&hFile, "data.bin", FA_READ );
    if (Res)
    {
        printf("SD: Open failed \n\r");
        return 0;
    }
    Res = f_read (&hFile, Data, 1024*2, &rd);
    Res = f_close (&hFile);
    return 1;
}
```


Links / Appendix

Links

freeRTOS Documentation:

<http://www.freertos.org>

FreeRTOS API documentation

<http://www.freertos.org/a00106.html>

Books:

- FreeRTOS Reference Manual.pdf
- FreeRTOS_Melot.pdf

LwIP Documentation:

<http://www.nongnu.org/lwip/>

Two Application Program's Interfaces (APIs)

- Netconn API: http://lwip.wikia.com/wiki/Netconn_API
- Socket API: <http://pubs.opengroup.org/onlinepubs/007908799/xnsix.html>
(compatible to posix- / BSD-sockets)

Annex: Accessing memory mapped Hardware-Devices

Using BSP driver functions:

```
#include "xparameters.h"
#include "xgpio.h"

#define LED_CHANNEL    1
#define SW_CHANNEL    2

XGpio Gpio;          /* The Instance of the GPIO Driver */

    // GPIO Initialisation:
XGpio_Initialize      (&Gpio, XPAR_AXI_GPIO_0_DEVICE_ID);    // =0
XGpio_SetDataDirection (&Gpio, LED_CHANNEL,    0xFFFFFFFF0); // 0 = Outputs
XGpio_SetDataDirection (&Gpio, SW_CHANNEL,    0xFFFFFFFFF); // 1 = Inputs

    //GPIO Data:
Data = XGpio_DiscreteRead      (&Gpio, SW_CHANNEL);
XGpio_DiscreteWrite            (&Gpio, LED_CHANNEL, Data);
```

Using direct I/O functions:

```
#include "xparameters.h"
#include "Xil_io.h"

    Xil_Out32 (Addr, Value)
    Xil_In32  (Addr)
```

Pointer usage:

```
#include "xparameters.h"
    *(u32*)Addr = Value;
```

AXI GPIO Registers

Registers

There are four internal registers in the AXI GPIO design as shown in [Table 4](#). The memory map of the AXI GPIO design is determined by setting the C_BASEADDR parameter. The internal registers of the AXI GPIO are at a fixed offset from the base address and are byte accessible.

Table 4: Registers

Base Address + Offset (hex)	Register Name	Access Type	Default Value (hex)	Description
C_BASEADDR + 0x00	GPIO_DATA	Read/Write	0x0	Channel 1 AXI GPIO Data Register
C_BASEADDR + 0x04	GPIO_TRI	Read/Write	0x0	Channel 1 AXI GPIO 3-state Register
C_BASEADDR + 0x08	GPIO2_DATA	Read/Write	0x0	Channel 2 AXI GPIO Data Register
C_BASEADDR + 0x0C	GPIO2_TRI	Read/Write	0x0	Channel 2 AXI GPIO 3-state Register

3-State Register: 1 = Input 0=Output

Default Base-Address: look in **xParameters.h** (0x4120_0000)

Change Echo Server to “Data-Server”

```
u16_t echo_port = 1111;

short int Data[1024];
int tick = 0;

void GetMyDAQ_Data (void)
{
    int i;
        for (i=0; i<1024; i++)
            Data[i] = i + tick*33;
        tick++;
}
```

```
GetMyDAQ_Data ();

nwrote = write (sd, Data, 1024*sizeof(short));

//nwrote = write (sd, recv_buf, n); // handle request by sending data back
```

Labview Client for Visualization

