

Outline

- ❑ Digital CMOS Design
- ❑ Arithmetic Operators
- ❑ Floating Point Arithmetic Operators
 - └ Square Root
 - └ Division



Division – Direct Method

In IEEE floating point standard a real number is represented as :

$$(-1)^S \times M \times 2^E$$

In 32-bit representation :

$$S \in \{0, 1\}$$

$$M \in [1, 2[$$

$$E \in \{-126, \dots, 127\}$$

$$\text{Or } M \in [0, 2[$$

$$\text{if } E = -127$$

normal

sub-normal



Division – Direct Method

Let U and V ($V \neq 0$) be two real numbers

We seek to calculate the real number $\frac{U}{V}$

$$U = (-1)^{S_U} \times M_U \times 2^{E_U}$$

$$V = (-1)^{S_V} \times M_V \times 2^{E_V}$$

then

$$\frac{U}{V} = (-1)^{S_U \oplus S_V} \times \frac{M_U}{M_V} \times 2^{E_U - E_V}$$



Division – Direct Method

If U or V are sub-normal numbers, the calculation of $\frac{U}{V}$ may lead to a loss of precision

Therefore if $M_U = 0 \quad \frac{U}{V} = 0$

and if M_U or $M_V \in]0, 1[$

E_U or E_V are decreased and M_U or M_V are $\times 2$ until they can fit within $[1, 2[$



Division – Direct Method

Since $M_U \in [1, 2[$ and $M_V \in [1, 2[$ $\frac{M_U}{M_V} \in]0.5, 2[$

Therefore if $M_U < M_V$

E_U is decreased and M_U is $\times 2$ so $\frac{M_U}{M_V} \in [1, 2[$



Division – Direct Method

Then, the problem can be stated as :

Given two positive real numbers $A \in [1, 4[$ and $B \in [1, 2[$

such as $A \geq B$ we seek to calculate $X_{Th} = \frac{A}{B}$ $X_{Th} \in [1, 2[$

Let X_n be an approximation of X_{Th} coded on $n+1$ bits

$$X_n = \sum_{j=0}^n x_{-j} \times 2^{-j} \quad \text{such as} \quad X_n \leq \frac{A}{B} < (X_n + 2^{-n})$$

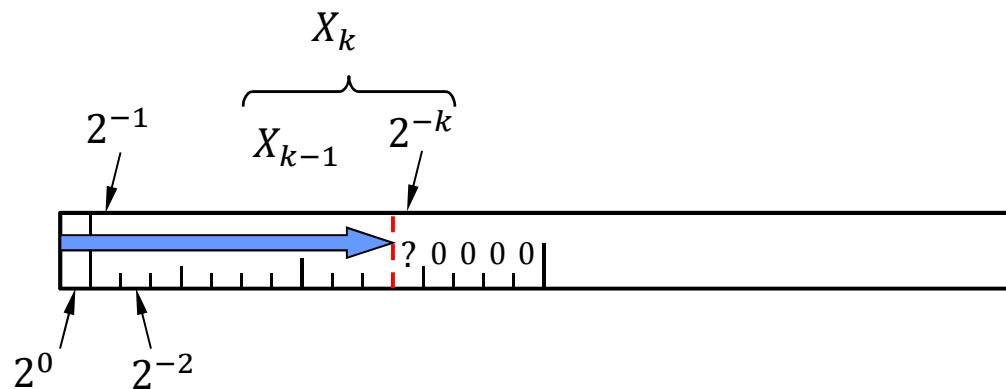


We propose to calculate X_n digit-by-digit

Division – Direct Method

Let $X_{k-1} = \sum_{j=0}^{k-1} x_{-j} \times 2^{-j}$ such as

$$X_{k-1} \leq \frac{A}{B} < (X_{k-1} + 2^{-(k-1)})$$



Division – Direct Method

Let $X_k = \sum_{j=0}^n x_{-j} \times 2^{-j}$ such as $X_k \leq \frac{A}{B} < (X_k + 2^{-k})$

At each iteration X_k is obtained from X_{k-1}

$$X_k = X_{k-1} + x_{-k} 2^{-k} \quad x_{-k} \in \{0, 1\}$$

2^{-k} is denoted W_k



Division – Direct Method

$$X_k \leq \frac{A}{B} < (X_k + 2^{-k})$$

$$BX_k \leq A < B(X_k + 2^{-k})$$

$$0 \leq A - BX_k < 2^{-k}B \quad \text{Let } \Delta_k = A - BX_k$$

$$0 \leq \Delta_k < 2^{-k}B \quad \text{yet } B < 2$$

then $0 \leq \Delta_k < 2^{-k} \times 2$

- $0 \leq \Delta_0 < 2$
- At each iteration the upper bound of Δ_k is divided by 2



Division – Direct Method

$$\Delta_k = A - BX_k$$

$$\Delta_k = A - B(X_{k-1} + x_{-k} 2^{-k})$$

$$\Delta_k = A - BX_{k-1} - x_{-k} 2^{-k} B$$

$$\Delta_k = \Delta_{k-1} - x_{-k} 2^{-k} B$$

$$2^k \Delta_k = 2^k \Delta_{k-1} - x_{-k} B$$

Let $D_k = 2^k \Delta_k$

$$D_k = 2D_{k-1} - x_{-k} B \quad x_{-k} = \begin{cases} 0 & \text{such as } 0 \leq D_k \\ 1 & \end{cases}$$



Division – Direct Method

Iteration scheme :

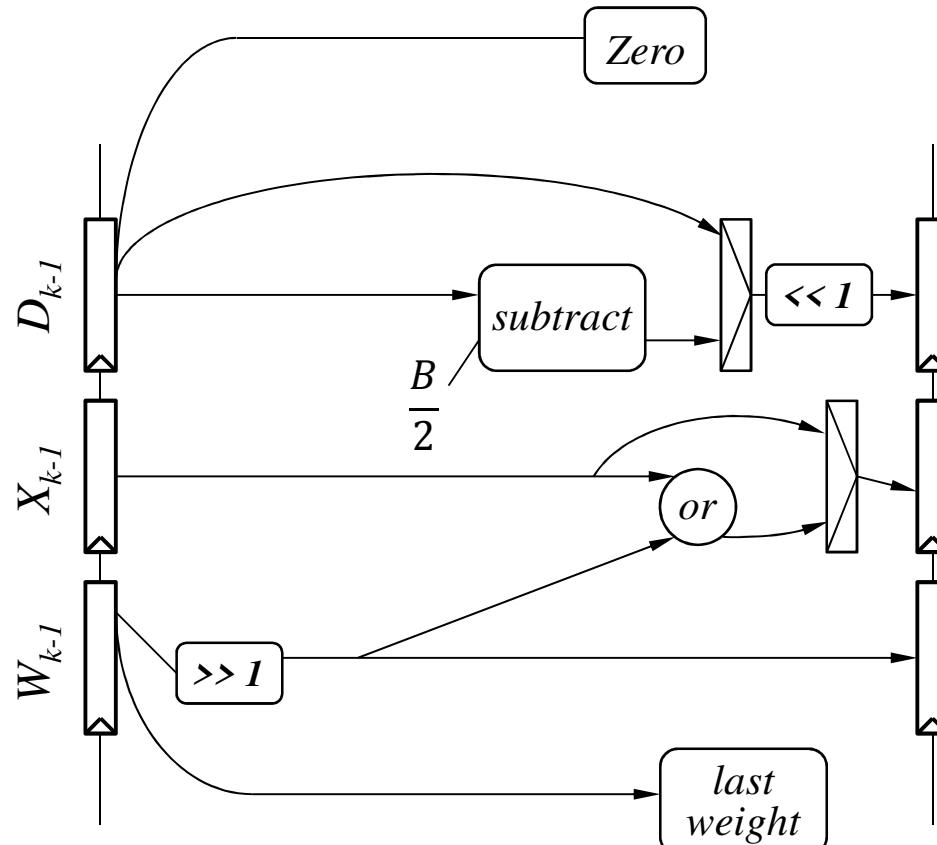
$$\begin{cases} D_k = 2 \left(D_{k-1} - x_{-k} \frac{B}{2} \right) \\ X_k = X_{k-1} + x_{-k} 2^{-k} \end{cases} \quad x_{-k} = \begin{cases} 0 & \text{such as } 0 \leq D_k \\ 1 & \text{otherwise} \end{cases}$$



Division – Direct Method

Implementation

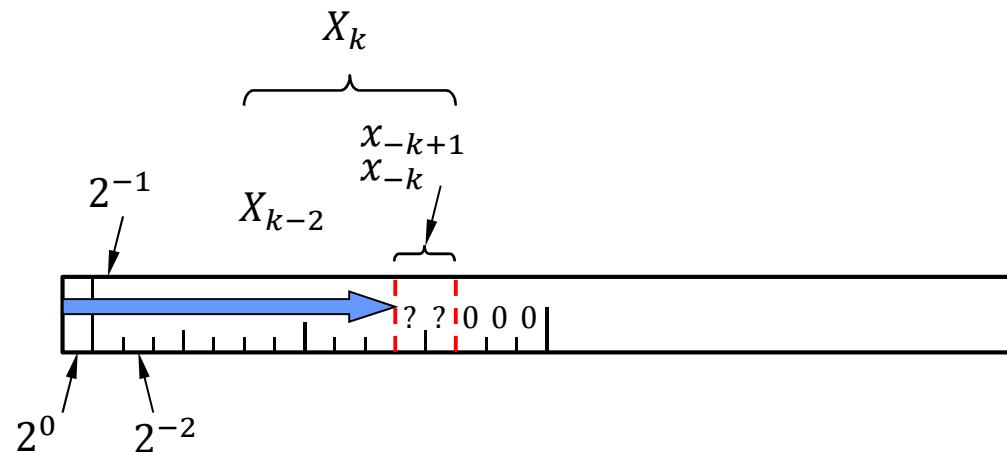
$$\begin{aligned}D_k &= 2 \left(D_{k-1} - x_{-k} \frac{B}{2} \right) \\X_k &= X_{k-1} + x_{-k} 2^{-1} W_{k-1} \\W_k &= 2^{-k} = 2^{-1} W_{k-1}\end{aligned}$$



Division – Direct Method

Implementation improvement – radix 4

$$X_k = X_{k-2} + (2x_{-k+1} + x_{-k})2^{-k}$$



Division – Direct Method

Implementation improvement – radix 4

$$\Delta_k = A - BX_k$$

$$\Delta_k = A - B(X_{k-2} + (x_{-k+1}2^{-k+1} + x_{-k}2^{-k}))$$

$$\Delta_k = \Delta_{k-2} - B(x_{-k+1}2^{-k+1} + x_{-k}2^{-k})$$

$$2^k \Delta_k = 2^k \Delta_{k-2} - B(2x_{-k+1} + x_{-k}) \quad \text{Let } D_k = 2^k \Delta_k$$

$$D_k = 4 \left(D_{k-2} - \frac{B}{4} (2x_{-k+1} + x_{-k}) \right)$$



Division – Direct Method

Implementation improvement – radix 4

Iteration scheme :

$$\begin{cases} D_k = 4 \left(D_{k-2} - \frac{B}{4} (2x_{-k+1} + x_{-k}) \right) \\ X_k = X_{k-2} + (2x_{-k+1} + x_{-k}) 2^{-k} \end{cases}$$

$$x_{-k+1}, x_{-k} = \begin{cases} 00 \\ 01 \\ 10 \\ 11 \end{cases} \text{ such as } 0 \leq D_k$$



Division – Direct Method

Implementation improvement – radix 4

$$x_{-k+1} = 0, \quad x_{-k} = 0 \quad \begin{cases} D_k = 4 D_{k-2} \\ X_k = X_{k-2} \end{cases}$$

$$x_{-k+1} = 0, \quad x_{-k} = 1 \quad \begin{cases} D_k = 4 \left(D_{k-2} - \frac{B}{4} \right) \\ X_k = X_{k-2} + 1 \times 2^{-k} \end{cases}$$

$$x_{-k+1} = 1, \quad x_{-k} = 0 \quad \begin{cases} D_k = 4 \left(D_{k-2} - 2 \frac{B}{4} \right) \\ X_k = X_{k-2} + 2 \times 2^{-k} \end{cases}$$

$$x_{-k+1} = 1, \quad x_{-k} = 1 \quad \begin{cases} D_k = 4 \left(D_{k-2} - 3 \frac{B}{4} \right) \\ X_k = X_{k-2} + 3 \times 2^{-k} \end{cases}$$



Division – Direct Method

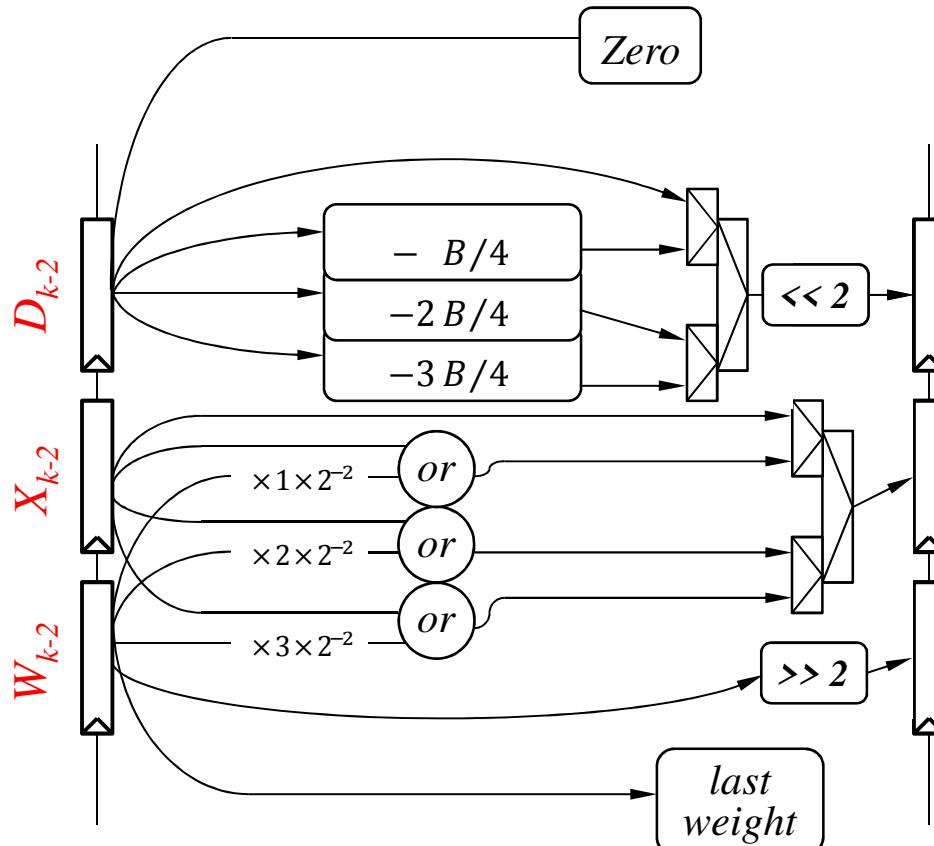
Implementation improvement – radix 4

$$\begin{aligned} x_{-k+1} &= 0, \\ x_{-k} &= 0 \end{aligned} \quad \left\{ \begin{array}{l} D_k = 4 D_{k-2} \\ X_k = X_{k-2} \end{array} \right.$$

$$\begin{aligned} x_{-k+1} &= 0, \\ x_{-k} &= 1 \end{aligned} \quad \left\{ \begin{array}{l} D_k = 4 \left(D_{k-2} - \frac{B}{4} \right) \\ X_k = X_{k-2} + 1 \times 2^{-k} \end{array} \right.$$

$$\begin{aligned} x_{-k+1} &= 1, \\ x_{-k} &= 0 \end{aligned} \quad \left\{ \begin{array}{l} D_k = 4 \left(D_{k-2} - 2 \frac{B}{4} \right) \\ X_k = X_{k-2} + 2 \times 2^{-k} \end{array} \right.$$

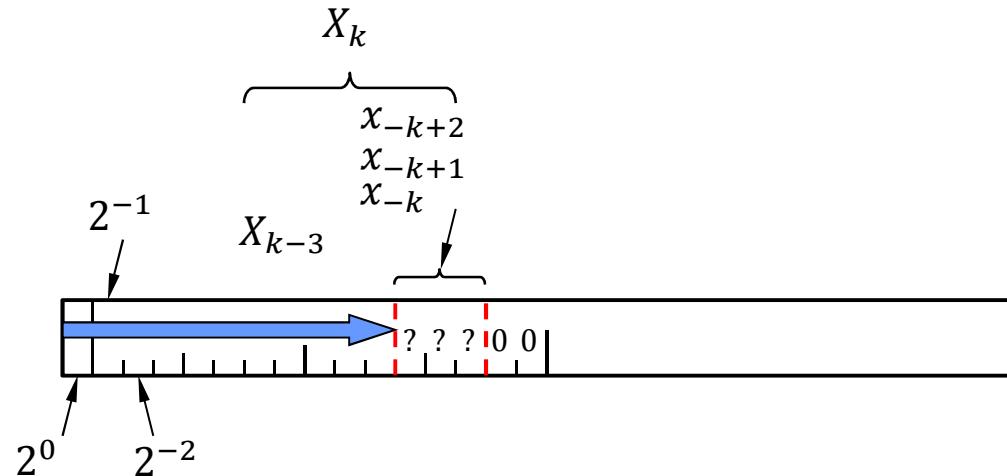
$$\begin{aligned} x_{-k+1} &= 1, \\ x_{-k} &= 1 \end{aligned} \quad \left\{ \begin{array}{l} D_k = 4 \left(D_{k-2} - 3 \frac{B}{4} \right) \\ X_k = X_{k-2} + 3 \times 2^{-k} \end{array} \right.$$



Division – Direct Method

Implementation improvement – radix 8

$$X_k = X_{k-3} + (4x_{-k+2} + 2x_{-k+1} + x_{-k})2^{-k}$$



Division – Direct Method

Implementation improvement – radix 8

$$\Delta_k = A - BX_k$$

$$\Delta_k = A - B(X_{k-3} + (x_{-k+2}2^{-k+2} + x_{-k+1}2^{-k+1} + x_{-k}2^{-k}))$$

$$\Delta_k = \Delta_{k-3} - B(x_{-k+2}2^{-k+2} + x_{-k+1}2^{-k+1} + x_{-k}2^{-k})$$

$$2^k \Delta_k = 2^k \Delta_{k-3} - B(4x_{-k+2} + 2x_{-k+1} + x_{-k}) \quad \text{Let } D_k = 2^k \Delta_k$$

$$D_k = 8 \left(D_{k-3} - \frac{B}{8} (4x_{-k+2} + 2x_{-k+1} + x_{-k}) \right)$$



Division – Direct Method

Implementation improvement – radix 8

Iteration scheme :

$$\begin{cases} D_k = 8 \left(D_{k-3} - \frac{B}{9} (4x_{-k+2} + 2x_{-k+1} + x_{-k}) \right) \\ X_k = X_{k-3} + (4x_{-k+2} + 2x_{-k+1} + x_{-k}) 2^{-k} \end{cases}$$

$$x_{-k+2}, x_{-k+1}, x_{-k} = \begin{cases} 000 \\ 001 \\ 010 \\ 011 \\ 000 \\ 001 \\ 010 \\ 011 \end{cases} \text{ such as } 0 \leq D_k$$



Division – Direct Method

Implementation improvement – radix 8

$$D_k = 8 \left(D_{k-3} - \frac{B}{8} (4x_{-k+2} + 2x_{-k+1} + x_{-k}) \right)$$
$$X_k = 8 (X_{k-3} \times (4x_{-k+2} + 2x_{-k+1} + x_{-k}))$$

