



INTERNATIONAL ATOMIC ENERGY AGENCY
UNITED NATIONS EDUCATIONAL, SCIENTIFIC AND CULTURAL ORGANIZATION



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SMR/169 - 17

WORKSHOP ON OPTICAL FIBER COMMUNICATION
(24 February - 21 March 1986)

CIRCUIT DESIGN FOR OPTICAL TRANSMISSION SYSTEMS (II PART)

presented by

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These are preliminary lecture notes, intended only for distribution to participants.

LOCAL AREA NETWORKS LANs

- DEFINITIONS AND SCOPE
- FIELDS OF APPLICATION
- BASIC CONSTITUENTS
- SOLUTIONS ON METALLIC CABLES
- IMPACT OF THE OPTICAL TECHNOLOGY
 - TOPOLOGY
 - ACCESS PROTOCOLS
 - DEVICES
- POWER BUDGET

7

FROM THE ORIGIN ...

A LOCAL AREA NETWORK IS A DATA COMMUNICATIONS SYSTEM WHICH ALLOWS A NUMBER OF INDEPENDENT DATA DEVICES TO COMMUNICATE DIRECTLY WITH EACH OTHER WITHIN A MODERATE SIZED GEOGRAPHIC AREA (SUCH AS A SINGLE OFFICE BUILDING, A WAREHOUSE OR A CAMPUS) OVER A PHYSICAL COMMUNICATION CHANNEL OF MODERATE TO HIGH DATA RATES AND CONSISTENTLY LOW ERROR RATE.



IEEE LOCAL AREA NETWORK STANDARD COMMITTEE
(PROJECT 802)

CSELT 09588/42

... TO THE PRESENT POSITION

A LOCAL AREA NETWORK IS AN INFORMATION TRANSPORTATION SYSTEM THAT SUSTAINS SPEECH, DATA, TEST, GRAPHICS, AND VIDEO, AND PROVIDES FOR HIGH-SPEED CONNECTION BETWEEN USERS WITHIN A SINGLE BUILDING OR A CAMPUS-THROUGH A COMMON WIRING SYSTEM, COMMON COMMUNICATION ADAPTER, COMMON ACCESS PROTOCOL TO ALLOW CONNECTION BETWEEN USERS, AND COMMON SHARED RESOURCES.



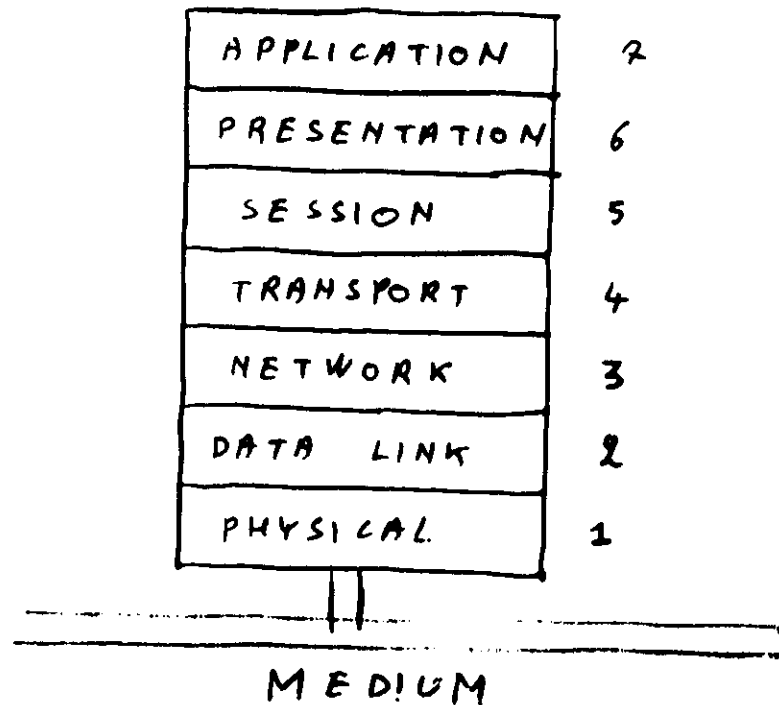
ECONOMIC EUROPEAN COMMUNITY ESPRIT PROGRAM

CSELT 09588/43

Business Service Characterization

SERVICE CLASS	SERVICE DESCRIPTION	SOURCE TRAFFIC TYPE	BIT-RATE kbit/s	BUSY-HOUR UTILIZATION	BUSY-HOUR THROUGHPUT kbit/s
0: TELEPHONY	STANDARD DIGITAL TEL. DSI, COMPRESSION TECHN.	STREAM	64	0.25 0.125	16 8
1: LOW-SPEED DATA	FACSIMILE VIDEOTEX INTERACTIVE DATA TEXT PROCESSING ...	BURSTY	1.2 + 9.6	0.12	0.6
2: MEDIUM-SPEED DATA	FACSIMILE (high-speed) HOST-TO-HOST COMM. LINE PRINTING ...	BURSTY	19.2 + 64	0.2	12.8
3: HIGH-SPEED DATA	MAINFRAME COMM. FILE TRANSFER HIGH-VOLUME PRINTING ...	BURSTY	500 + 1000	0.3	150
4: COMPRESSED DIGITAL VIDEO AND GRAPHICS	DESK-TOP VIDEOTELEPHONY VIDEOCONFERENCE HIG-RES. GRAPHICS ...	STREAM	> 2000	0.15	300
5: HIGH-QUALITY VIDEO	ANALOG or DIGITAL TV	-	6 Mbit/s > 100 Mbit/s	-	-

OSI REFERENCE MODEL
OPEN SYSTEM INTERCONNECTION



LANs APPLICATIONS

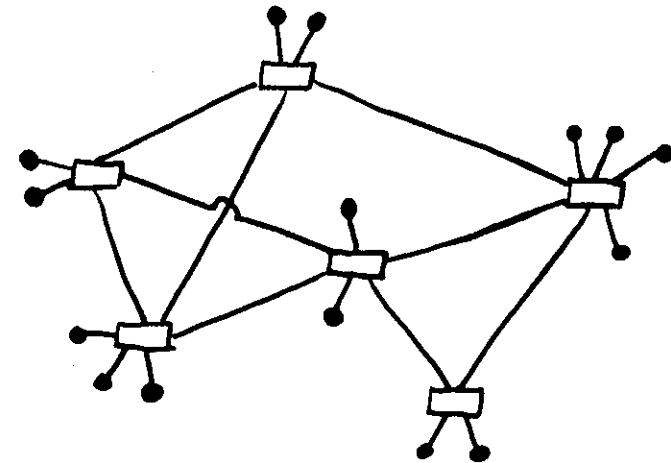
- OFFICE AUTOMATION
- NEWSPAPER EDITING SYSTEMS
- FACTORY AUTOMATION
- PROCESS CONTROL
- HOSPITALS
- UNIVERSITIES AND RESEARCH LABORATORIES

BASIC CONSTITUENTS OF LANs

- TRANSMISSION NETWORK
- MECHANISM FOR THE CONTROL OF THE TRANSMISSION
- INTERFACES TO CONNECT TERMINALS
- PROTOCOLS TO CONTROL THE INFORMATION FLOW
- GATEWAYS TO COMMUNICATE TO THE EXTERNAL PUBLIC TELECOM NETWORK

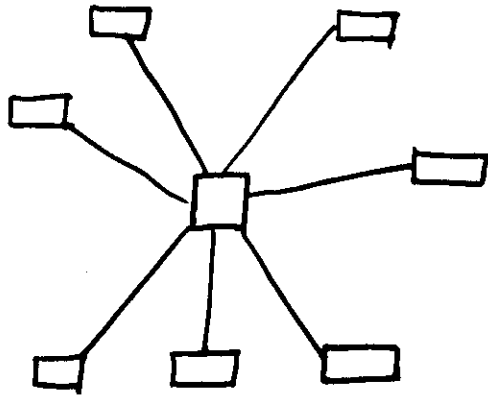
NETWORK TOPOLOGY

PATTERN OF THE INTERCONNECTIONS
BETWEEN THE NODES OF THE
NETWORK



- USER
- NODE
- INTERCONNECTION (LAYER 1)

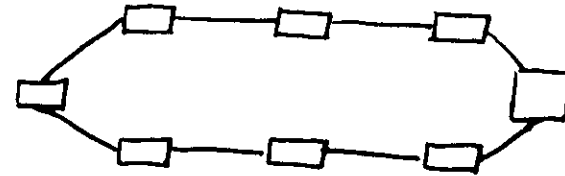
THE STAR NETWORK



- SUITABLE FOR A CENTRALIZED CONTROL
- RELIABILITY FOR THE CENTRAL NODE
- MORE COMPLEX WIRING

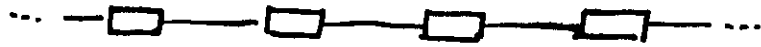
□ NETWORK CONTROL CENTRE

THE RING NETWORK



- NO ROUTING DECISIONS
- NEED FOR RECOGNIZING THE ADDRESSES
- DISTRIBUTED CONTROL POSSIBLE
- RELIABILITY FOR EVERY NODE

THE BUS NETWORK



- REDUCED AMOUNT OF WIRING
- NO ROUTING DECISIONS
- NEED OF RECOGNIZING THE ADDRESS
- DISTRIBUTED CONTROL POSSIBLE
- MANY DIFFERENT VARIATIONS:
 - + UNIDIRECTIONAL/BIDIRECTIONAL
 - + ONE/TWO CABLES

NETWORK CONTROL

PROBLEM: WHICH NODE MAY TRANSMIT
AT ANY GIVEN TIME?

SOLUTION: ACCESS PROTOCOLS

RULES THAT ALLOW THE USE OF
A COMMON COMMUNICATION RESOURCE
BY MANY DISTANT NODES

PROTOCOL CLASSIFICATION

- FIXED / VARIABLE ALLOCATION
- TOPOLOGY
- NUMBER OF SERVED NODES
- DISTRIBUTED / CENTRALIZED CONTROL

CHOICE

- TOPOLOGY
- TRAFFIC REQUIREMENTS
- EFFICIENCY

$$E = \frac{\text{CAPACITY USABLE BY A STATION}}{\text{TOTAL CHANNEL CAPACITY}}$$

- TECHNOLOGY

CONTROL PASSING

THE TRANSMISSION CHANNEL IS ASSIGNED IN CYCLE TO THE STATIONS

A SILENT STATION RELEASES IT IMMEDIATELY

POLLING - CONTROL DISTRIBUTED BY A MASTER STATION (NCC).
SUITABLE FOR STAR NETWORKS

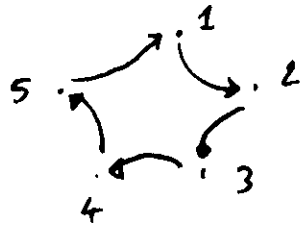
TOKEN - DAISY CHAIN

CONTROL PASSED DIRECTLY FROM A STATION TO THE NEXT ONE

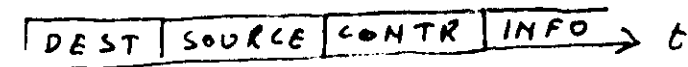
EXPLICIT TOKEN: MESSAGE WITH ADDRESS, INFORMATION, CONTROLS, ETC. USED IN BUS

IMPLICIT TOKEN: PARTICULAR BIT PATTERN, FIXED. USED IN RINGS

CONTROL TOKEN



EXPLICIT TOKEN



IMPLICIT TOKEN



RANDOM ACCESS PROTOCOLS

WITH MANY STATIONS, CONTROL PASSING PROTOCOLS BECOME LESS EFFICIENT

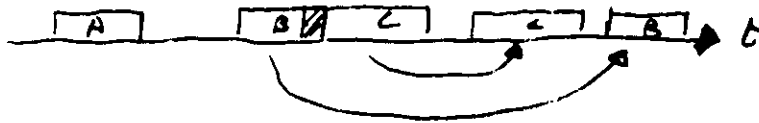
ACCESS TIME CAN BE REDUCED CONSIDERING THAT ONLY A FEW STATIONS WANT TO TRANSMIT AT A GIVEN INSTANT.

THE ACTIVITY OF THE CHANNEL IS MONITORED TO DETECT:

- ABSENCE OF TRANSMISSION
- CORRECT TRANSMISSION
- COLLISION - SIMULTANEOUS TRANSMISSION ATTEMPT BY SEVERAL STATIONS

ALHOA (1970)

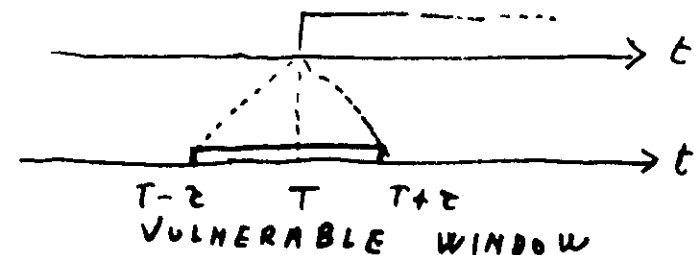
- A PACKET IS TRANSMITTED AS SOON AS AVAILABLE
- AFTER A COLLISION, A RETRANSMISSION IS MADE AFTER A RANDOM TIME



- + SIMPLEST ALGORITHM
- + SUITABLE FOR ALMOST EVERY TO-POLOGY
- VARIABLE DELAY - NOT SUITABLE FOR STREAM TRAFFIC.

CARRIER SENSE MULTIPLE ACCESS CSMA

- A PACKET IS TRANSMITTED ONLY IF THE CHANNEL IS FREE



T = END TO END TRANSMISSION DELAY

NO COLLISION OUTSIDE THE WINDOW

AFTER A COLLISION, NEW ATTEMPT AFTER A RANDOM DELAY

IF THE CHANNEL IS BUSY, STATUS IS MONITORED (WITH PROBABILITY P) OR A NEW ATTEMPT IS MADE (WITH PROBABILITY $1-P$) AFTER A RANDOM DELAY

PROTOCOL SUITABLE FOR BUSES

CSMA-CD (COLLISION DETECTION)

WHEN A COLLISION IS DETECTED
THE TRANSMISSION IS STOPPED

THE TIME IN WHICH COLLISIONS
TAKE PLACE IS NO MORE WASTED.
EFFICIENCY IS IMPROVED.

COLLISION MUST BE DETECTED
IN TIME; A MORE COMPLICATED
MONITORING RECEIVER IS NEEDED.

SPECIFIC ACCESS PROTOCOLS

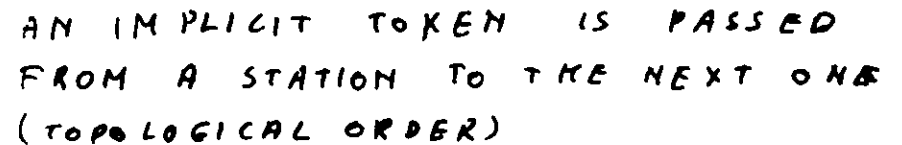
WHEN THE TOPOLOGY IS A RING
OR A BUS AND

- THE STATIONS CAN BE LABELED
IN ORDER
- THE BIT TIMING IS CONTINUOUS-
LY AVAILABLE

* THE EFFICIENCY CAN BE
FURTHER IMPROVED

* IT IS EASY TO ACCOMMODATE
BOTH PACKET AND STREAM
SERVICES.

CONTROL TOKEN



EACH STATION MAY APPEND:

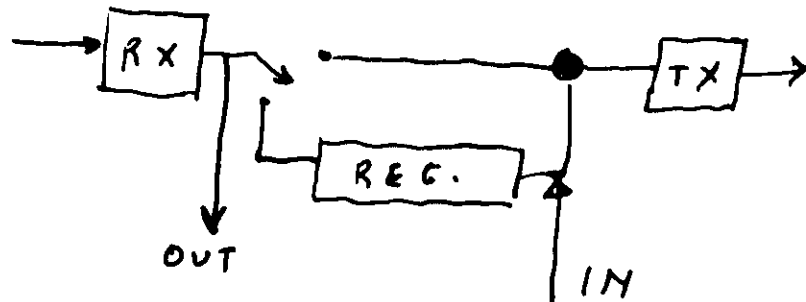
- NO MESSAGE
- A FIXED-LENGTH MESSAGE
- A VARIABLE-LENGTH MESSAGE

PROBLEMS

- IN BUSES! A HEAD STATION MUST ISSUE
A NEW TOKEN AT EVERY CYCLE

IN RINGS : THE TOKEN MUST BE
UNIQUE AND MUST BE PROTECTED
AGAINST THE ERRORS

REGISTER INSERTION



A STATION TRANSMITS WHEN THE CHANNEL IS FREE. INCOMING MESSAGES ARE BUFFERED IN THE REGISTER. TRANSMITTED MESSAGES ARE REMOVED WHEN THEY COME BACK IN THE REGISTER.

- + SIMPLE, FULLY DISTRIBUTED ALGORITHM
- + HIGH EFFICIENCY
- MESSAGES OF FIXED LENGTH
- SUPERVISION AGAINST INDEFINITE CIRCULATION OF CORRUPTED MESSAGES

TRANSMISSION MEDIA

COPPER PAIR



- + LOW COST
- + EASY INSTALLATION
- BANDWIDTH UP TO 5 MHz
- ATTENUATION MEDIUM TO HIGH
- NEEDS SCREENING

TRANSMISSION MEDIA

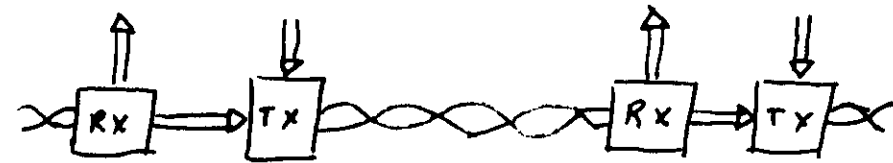
COAXIAL CABLE



- + SCREENED
- + 5-50 MHz BANDWIDTH PER CARRIER
- + CATV TECHNOLOGY AVAILABLE (USA)
- MODERATE COST
- CAREFUL INSTALLATION

EXAMPLE OF IMPLEMENTATION

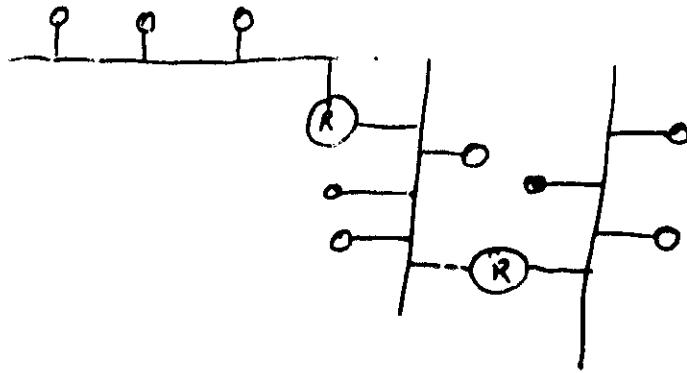
IBM ZURICH RING



- RING WITH ACTIVE STATIONS
- POINT TO POINT UNIDIRECTIONAL LINKS
- STATION DISTANCE $\leq 1 \text{ km}$
- TOKEN PASSING PROTOCOL
- TWISTED PAIR, 4 Mbit/s
- ERROR RATE 10^{-9}

EXAMPLE OF IMPLEMENTATION

XEROX¹, ETHERNET



- MULTISEGMENT (MAX. 3) BUS
- COAXIAL CABLE, 2.5 Km MAX.
- CSMA-CD AT 10 Mbit/s
- MAX. 300 STATIONS

ADVANTAGES OF OPTICAL FIBRES

- LOW ATTENUATION $\leq 3 \text{ dB/Km}$
- WIDE BANDWIDTH $\geq 600 \text{ MHz} \cdot \text{Km}$
- REDUCED SIZE AND WEIGHT
- IMMUNITY TO ELECTRIC INTERFERENCE
- ELECTRICAL INSULATION
- NO PECULIAR CONSTRAINTS ON ACCESS PROTOCOLS

DRAWBACKS OF FIBRES

- REDUCED COUPLING EFFICIENCY (SIDE COUPLING)
- ~~NO~~ HIGH IMPEDANCE TAPPING IS NOT POSSIBLE

WHICH FIBRE FOR LANs ?

- G.I. 50/125 μm MULTIMODE
- S.I. 10/125 μm SINGLE MODE
- S.I. LARGE CORE MULTIMODE ($>100 \mu\text{m}$)

CHOICE ACCORDING TO :

- ATTENUATION
- BANDWIDTH
- SPLICE LOSS
- COST

SINGLE FIBRE INTERCONNECTION CABLE

- G.I. 50/125 μm
- ATTENUATION: 2.5 dB/Km @ 860 nm
- BANDWIDTH: 500 MHz $\cdot$ Km
- EXTERNAL DIAMETER: 2.5 mm
- TENSILE STRENGTH: 20 kg
- BENDING RADIUS: 30 mm MIN
- TEMPERATURE RANGE: $-20^{\circ}\text{C} \div +60^{\circ}\text{C}$
- SECTION LENGTH: 250 $\div$ 1000 m

OPTICAL SOURCES

TYPE	WAVELENGTH	PEAK POWER	BANDWIDTH	SPECTRAL WIDTH
LED	0.8 - 0.9 μm	50 μW	50 MHz	40 nm
	1.3 μm	50 μW	100 MHz	100 nm
LASER	0.8 - 0.9 μm	2 mW	≥ 500 MHz	≥ 4 nm
	1.3 μm	2 mW	≥ 500 MHz	≥ 4 nm
	1.5 μm	2 mW	≥ 500 MHz	≥ 4 nm

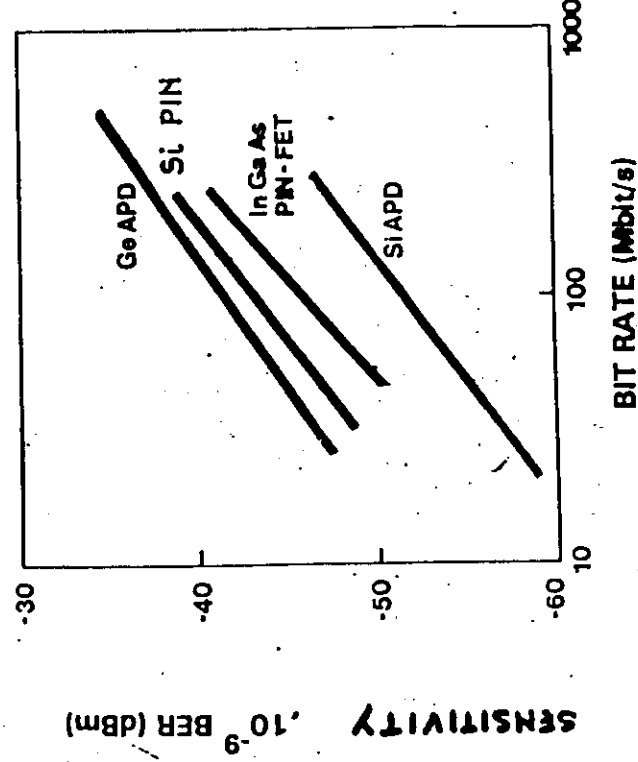
CSELT 09688/92

OPTICAL DETECTORS

TYPE	WAVELENGTH	GAIN	BANDWIDTH
Si-ADP	0,8 - 0,9 μm	~ 100	≥ 500 MHz
Ge-APD	1,1 - 1,6 μm	~ 10	≥ 500 MHz
p-i-n InGaAs InGaAsP	1,3 - 1,6 μm	1	≥ 500 MHz

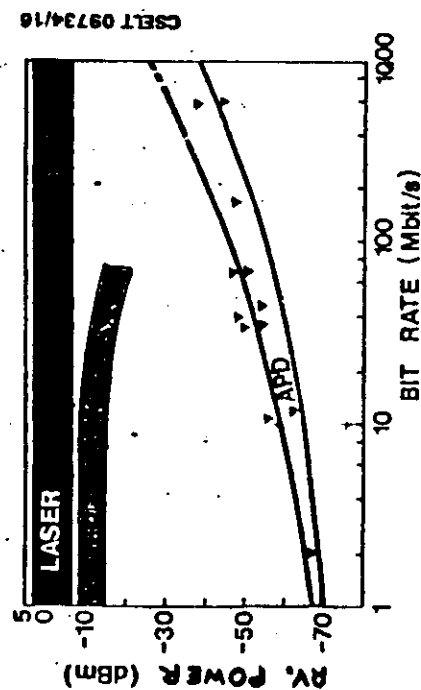
CSRLT 0966000

OPTICAL RECEIVERS SENSITIVITY



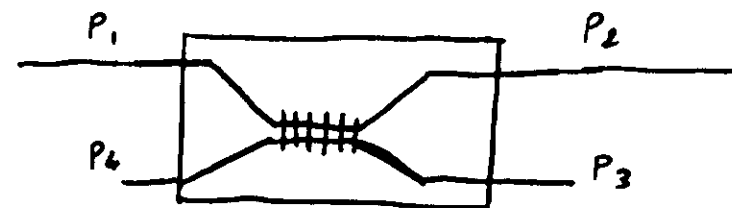
OPTICAL SOURCES AND RECEIVERS

average power vs. bit rate



OPTICAL COUPLERS

MULTI PORT DEVICES FOR LAUNCHING OR EXTRACTING POWER FROM A FIBRE



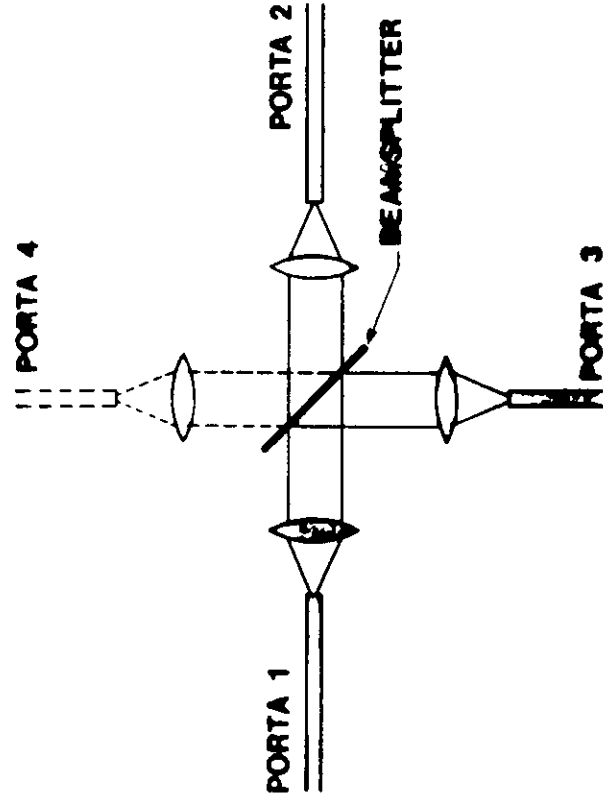
COUPLING $1 \rightarrow 2$ P_2 / P_1

COUPLING $1 \rightarrow 3$ P_3 / P_4

OUTPUT RATIO P_2 / P_3

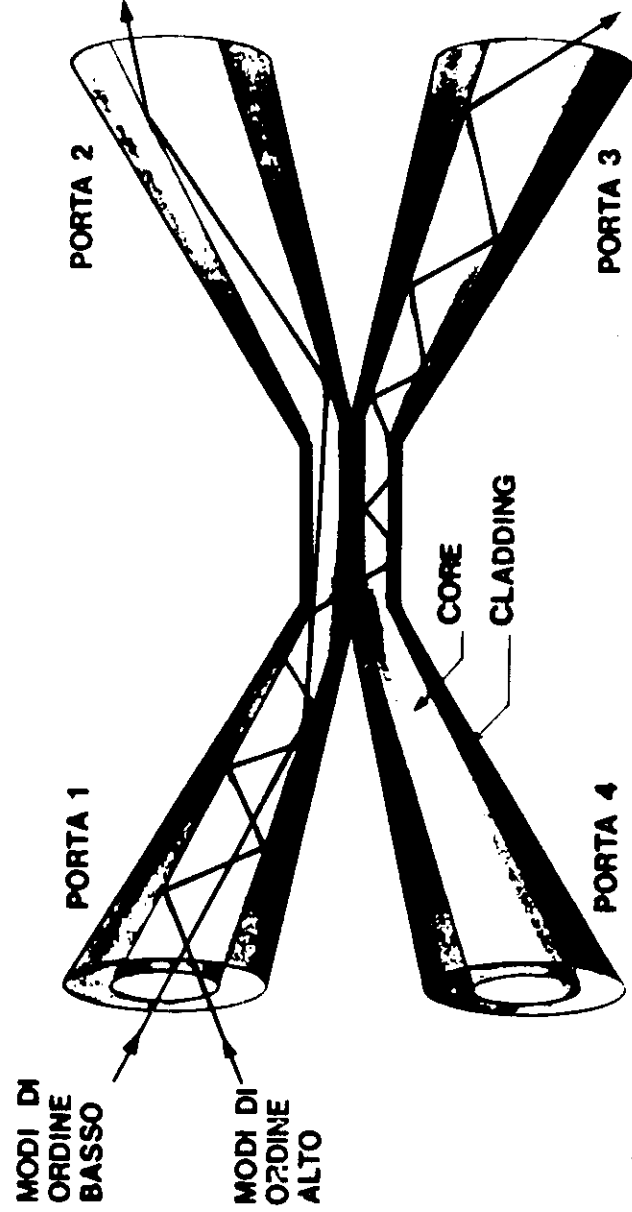
INSERTION LOSS $\frac{P_1 - (P_2 + P_3)}{P_1}$

ESEMPIO DI UN ACCOPIATORE DEL TIPO A RIFLESSIONE



CSELT 10000

ESEMPIO DI ACCOPIATORE DEL TIPO AD ACCOPPIAMENTO DI MODI



CSELT 10000

OPTICAL COUPLERS

POWER BUDGET

- RING
- BIT RATE 34 MBIT/S
- $P_{TX} = -5 \text{ dBm}$ (LASER)
 $S_{RX} = -49 \text{ dBm}$
- INSERTION LOSS WITH CONNECTORS
3.5 dB

OPTIMAL OUTPUT RATIO 10/1

AFTER 5 COUPLERS A REGENERATOR IS REQUIRED

TOPOLOGIES FOR OPTICAL LANs

RING

- WITH OPTICAL COUPLERS

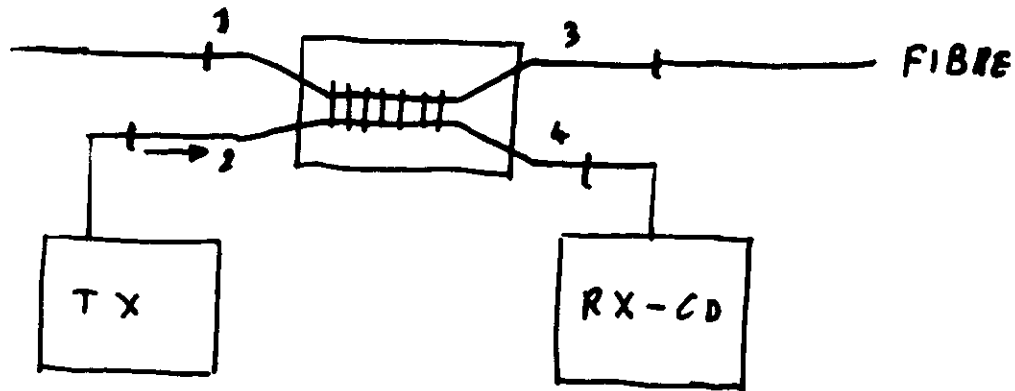
- * REPEATERS EVERY 5 ÷ 10 STATIONS
- * PROBLEMS WITH CIRCUITS /
TIMING, DYNAMIC RANGE

- FULLY ACTIVE

- * SEQUENCE OF POINT-TO-POINT LINKS
- * PROBLEMS OF RELIABILITY AND FAILURE

OPTICAL COUPLERS

EXAMPLE OF INSERTION



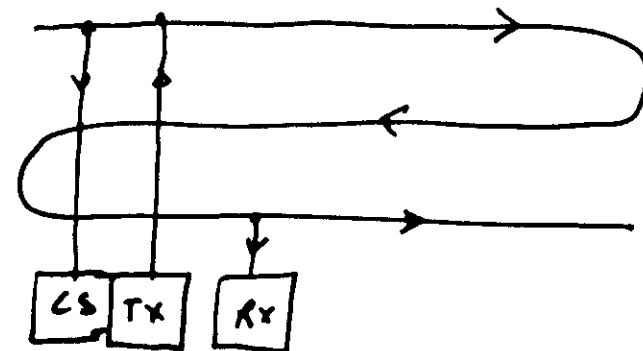
$$\frac{P_4}{P_1} = \frac{P_3}{P_2} \quad \frac{P_3}{P_1} = \frac{P_4}{P_2}$$

COUPLERS ARE RECIPROCAL AND SYMMETRICAL

COUPLING RATIO IS TO BE OPTIMIZED ACCORDING TO THE APPLICATION

TOPOLOGIES FOR OPTICAL LANs BUS

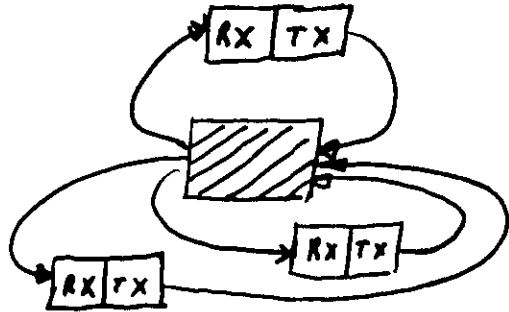
- MONODIRECTIONAL BUS FAVOURED LESS REPEATERS (WITH COUPLERS, PROTOCOLS WITH BETTER PERFORMANCES (ORDERED STATIONS))
- REPEATERS REQUIRED
- SAME RELIABILITY AND CIRCUIT PROBLEMS AS FOR RING



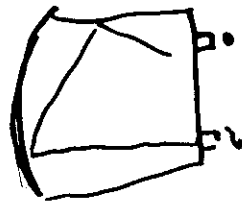
2-BUS

TOPOLOGIES FOR OPTICAL LANs

PASSIVE STAR



- FULLY PASSIVE STRUCTURE
HIGH RELIABILITY
- EQUIVALENT TO A BUS
THE BUS PROTOCOLS CAN BE USED,
EXCLUDING THOSE REQUIRING AN OR-
DERING IN THE STATIONS
- LIMITS : NUMBER OF PORTS



LINE CODES FOR LANs

REQUIREMENTS

- SPECTRUM SHAPING
- LOW REDUNDANCY ADDED
- MANY TRANSITIONS TO EXTRACT
TIMING
- WORDS FOR COMMANDS
- EASY WORD ALIGNMENT
- PERFORMANCE MONITORING
- EASE OF IMPLEMENTATION
- 2 LEVELS ONLY

EXAMPLES OF CODES

• 8 B 1 C

$\bar{D}_0$ D_0 P D_1 D_2 D_3 D_4 D_5 D_6 D_7

- SIMPLE TO BE IMPLEMENTED
- DC COMPONENT IN THE CONTINUOUS SPECTRUM
- COMMAND WORD BY CODE VIOLATION
- ERROR MONITORING AND FRAME ALIGNMENT ON $\bar{D}_0 D_0$

EXAMPLES OF CODES

16 B 18 B

D_0 D_1 D_2 D_{15} I P

I = INVERSION

P = PARITY

- RUNNING DIGITAL SUM IS LIMITED
NO DC COMPONENT IN THE SPECTRUM
- ERROR MONITORING AND ALIGNMENT BASED ON I P
- COMMAND WORDS BY VIOLATIONS

EXAMPLES OF CODES

FIBER DISTRIBUTED DATA INTERFACE

4B-5B

16 DATA WORDS

8 COMMAND WORDS

8 PROHIBITED WORDS

FRAME ALIGNMENT AND ERROR
MONITORING WITH PROHIBITED
WORDS

SUITABLE UP TO 20-30 MB/s

EXAMPLES OF CODES

CMI - CODED MARK INVERSION

0 → 01

1 ↗ 00
↘ 11 ALTERNATE

- SIMPLE CIRCUITS
- EASY CLOCK RECOVERY
- MONITORING OF ERRORS ON ALTER-
NATIONS
- COMMAND WORDS BY VIOLATIONS

DRAWBACK: THE BANDWIDTH
IS DOUBLED

SYNCHRONIZATION IN LANs

THE RECEIVER EXTRACTS TIMING FROM THE INCOMING SIGNAL

- THE ONLY POSSIBLE SOLUTIONS ON BUSES WITH BURSTY ACTIVITY
- PREAMBLE FOR TRAINING NEEDED
- DIFFICULT MATCHING WITH HIGHER LEVEL CLOCK

SYNCHRONIZATION IN LANs

THE RECEIVER IS PERMANENTLY LOCKED TO THE SIGNAL FROM THE PREVIOUS TRANSMITTER AND FEEDS ALL THE NODE.

- SUITABLE IN RINGS
- RELIABLE MASTER CLOCK REQUIRED
- JITTER IF THE NODES ARE MANY
- JITTER REDUCER IN THE RECEIVER OF THE MASTER STATION

SYNCHRONIZATION IN LANs

THE RECEIVER IS LOCKED TO THE PREVIOUS NODE, BUT IT HAS ALSO A LOCAL CLOCK

DATA ARE WRITTEN IN AN ELASTIC BUFFER AND READ BY THE LOCAL CLOCK. THE DIFFERENCE IN FREQUENCY IS COMPENSATED IN IDLE PERIODS

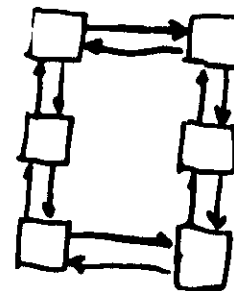
- SUITABLE FOR RINGS
- NO MASTER CLOCK
- UNTITTERED CLOCK IN THE NODES
- SMALL OVERHEAD
- MORE COMPLEX CIRCUITS

RELIABILITY OF A LAN

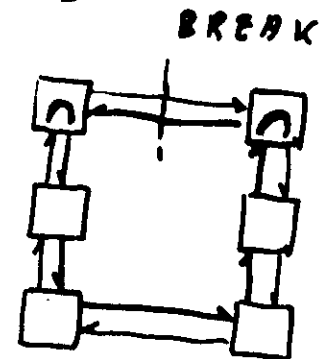
MALFUNCTION	COUNTERMEASURE
NODE FAILURE	BYPASS
CABLE BREAK	CABLE DUPLICATION

IN GENERAL

- NETWORK DUPLICATION IS EXPENSIVE
- NETWORK RECONFIGURATION DEPENDS ON THE PROTOCOLS



NORMAL
RING



EMERGENCY
LOOP
RECONFIGURED RING

