

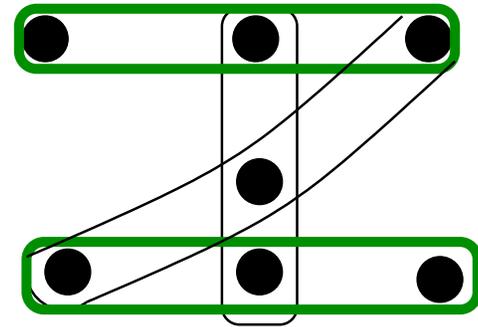
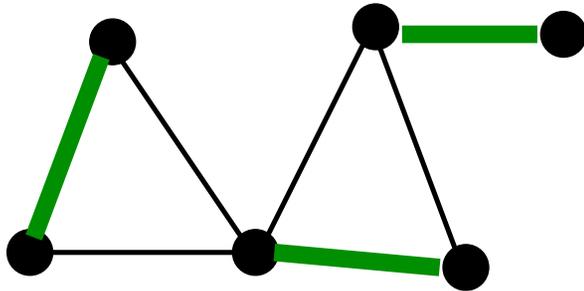
Packing and Covering in Uniform Hypergraphs

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Packing

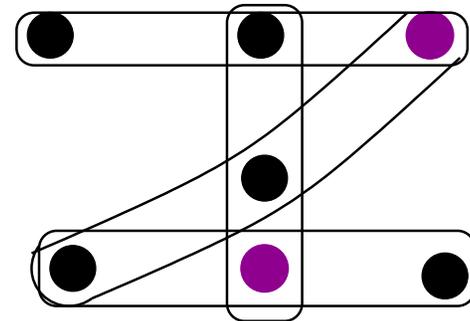
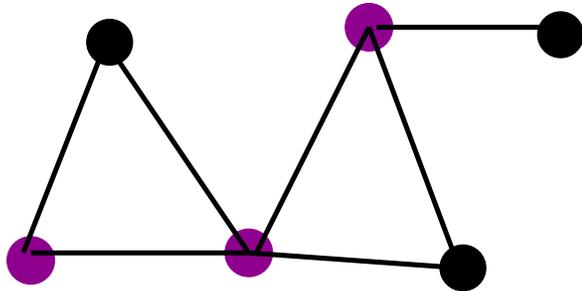
Let \mathcal{H} be a hypergraph. A **packing** or **matching** of \mathcal{H} is a set of pairwise disjoint edges of \mathcal{H} .



The parameter $\nu(\mathcal{H})$ is defined to be the maximum size of a packing in \mathcal{H} .

Covering

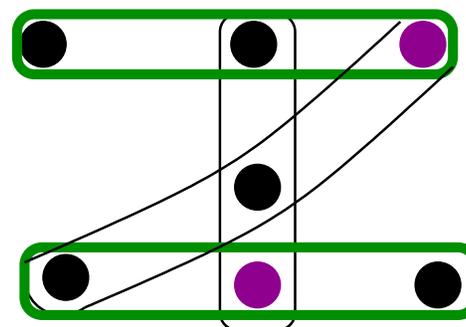
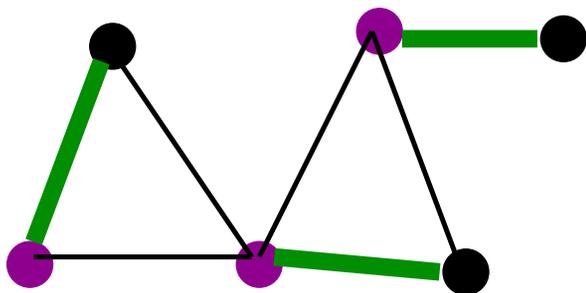
A **cover** of the hypergraph \mathcal{H} is a set of vertices C of \mathcal{H} such that every edge of \mathcal{H} contains a vertex of C .



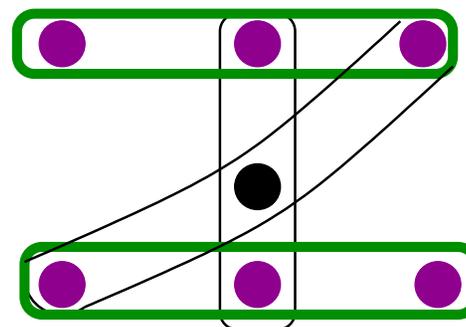
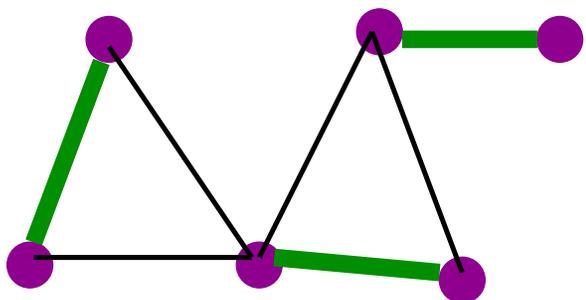
The parameter $\tau(\mathcal{H})$ is defined to be the minimum size of a cover of \mathcal{H} .

Comparing $\nu(\mathcal{H})$ and $\tau(\mathcal{H})$

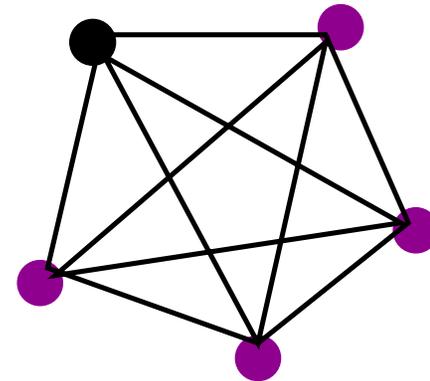
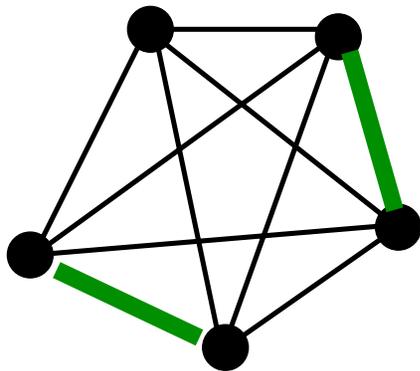
For **every** hypergraph \mathcal{H} we have $\nu(\mathcal{H}) \leq \tau(\mathcal{H})$.



For every **r -uniform** hypergraph \mathcal{H} we have $\tau(\mathcal{H}) \leq r\nu(\mathcal{H})$.



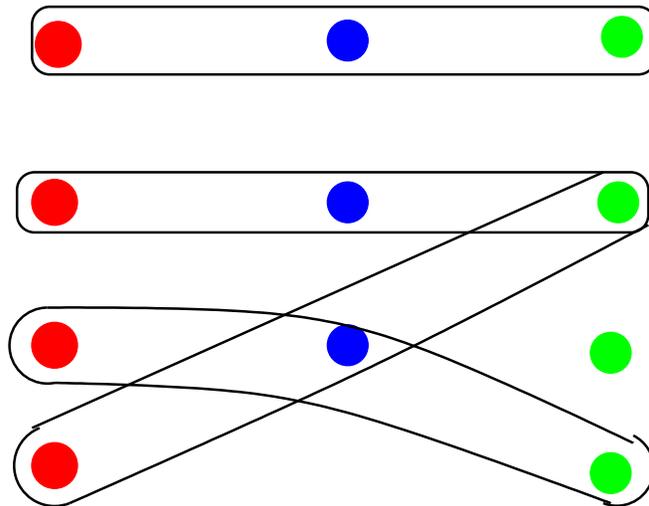
The upper bound $\tau(\mathcal{H}) \leq r\nu(\mathcal{H})$ is attained for certain hypergraphs, for example for the complete r -uniform hypergraph \mathcal{K}_{rt+r-1}^r with $rt+r-1$ vertices, in which $\nu = t$ and $\tau = rt$.



Ryser's Conjecture

Conjecture: Let \mathcal{H} be an r -partite r -uniform hypergraph. Then

$$\tau(\mathcal{H}) \leq (r - 1)\nu(\mathcal{H}).$$



This conjecture dates from the early 1970's.

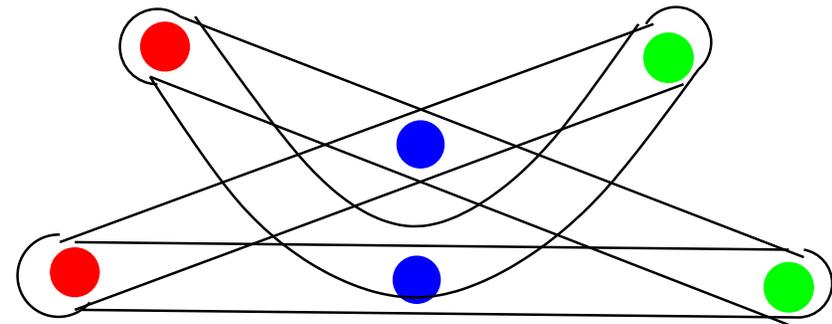
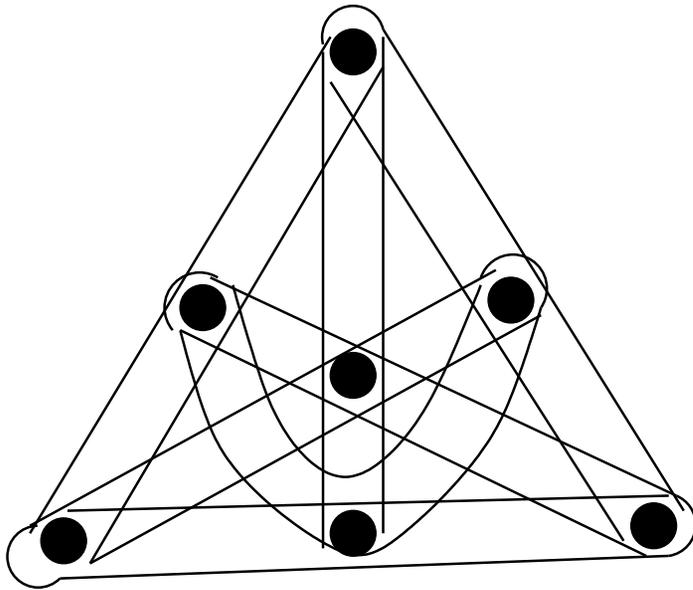
A stronger conjecture

Conjecture (Lovász): Let \mathcal{H} be an r -partite r -uniform hypergraph. Then there exists a set S of at most $r - 1$ vertices such that the hypergraph \mathcal{H}' formed by removing S satisfies

$$\nu(\mathcal{H}') \leq \nu(\mathcal{H}) - 1.$$

Results on Ryser's Conjecture

- $r = 2$: This is König's Theorem for bipartite graphs.
- $r = 3$: Known (proved by Aharoni, 2001)
- $r = 4$ and $r = 5$: Known for small values of $\nu(\mathcal{H})$, namely for $\nu(\mathcal{H}) \leq 2$ when $r = 4$ and for $\nu(\mathcal{H}) = 1$ when $r = 5$. (Tuza)
- whenever $r - 1$ is a prime power: If true, the upper bound is best possible.



Here $\nu(\mathcal{H}) = 1$ and $\tau(\mathcal{H}) = r - 1$.

On Ryser's Conjecture

Theorem (PH, Scott 2012) There exists $\epsilon > 0$ such that for every 4-partite 4-uniform hypergraph \mathcal{H} we have

$$\tau(\mathcal{H}) \leq (4 - \epsilon)\nu(\mathcal{H}).$$

Theorem (PH, Scott 2012) There exists $\epsilon > 0$ such that for every 5-partite 5-uniform hypergraph \mathcal{H} we have

$$\tau(\mathcal{H}) \leq (5 - \epsilon)\nu(\mathcal{H}).$$

On Ryser's Conjecture for $r = 3$

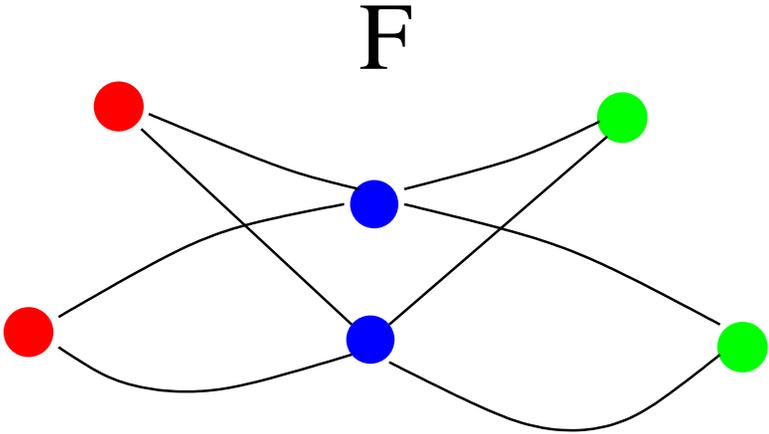
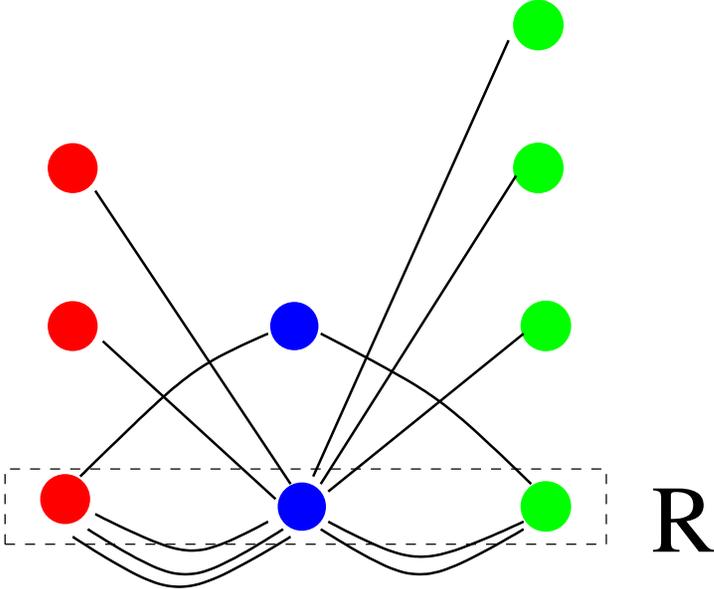
Theorem (Aharoni 2001): Let \mathcal{H} be a 3-partite 3-uniform hypergraph. Then

$$\tau(\mathcal{H}) \leq 2\nu(\mathcal{H}).$$

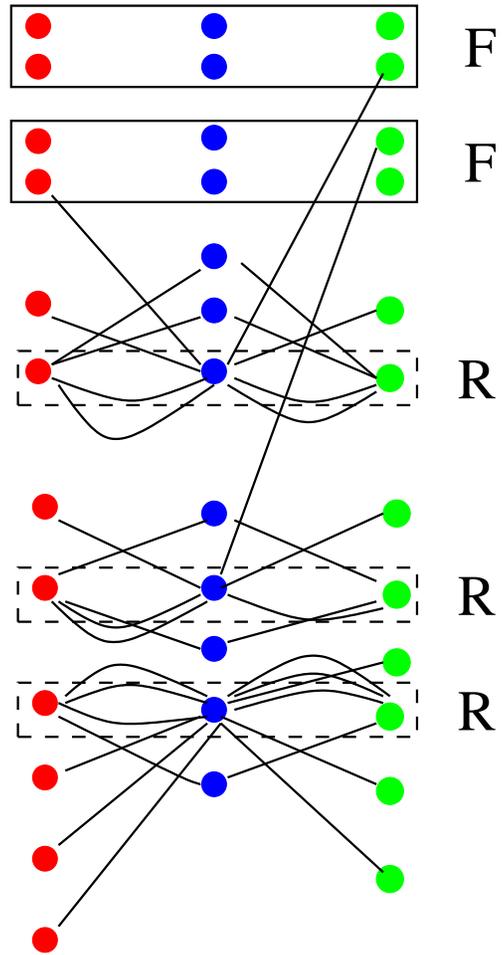
Proof: Uses topological connectivity of matching complexes of bipartite graphs.

Q: What is \mathcal{H} like if it is a 3-partite 3-uniform hypergraph with $\tau(\mathcal{H}) = 2\nu(\mathcal{H})$?

Extremal hypergraphs for Ryser's Conjecture



Home base hypergraphs



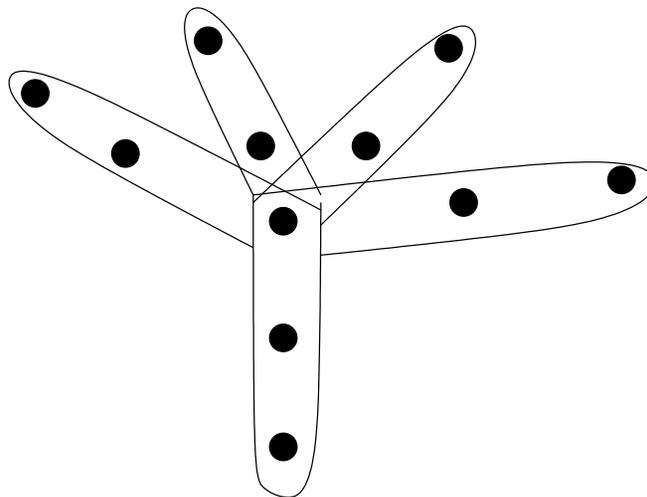
Extremal hypergraphs for Ryser's Conjecture

Theorem (PH, Narins, Szabó): Let \mathcal{H} be a 3-partite 3-uniform hypergraph with $\tau(\mathcal{H}) = 2\nu(\mathcal{H})$. Then \mathcal{H} is a home base hypergraph.

Some proof ingredients

The **nontrivial bounds** for Ryser's conjecture for $r = 4$ and $r = 5$ use Tuza's result that the conjecture is true for **intersecting hypergraphs**, together with several classical results.

A **sunflower** \mathcal{S} with centre C in a hypergraph is a set of edges satisfying $S \cap S' = C$ for all $S \neq S'$ in \mathcal{S} .



The Sunflower Theorem

Theorem (Erdős and Rado): Every hypergraph of rank r with more than $(t - 1)^r r!$ edges contains a sunflower with t petals.

Fact: If an intersecting hypergraph \mathcal{H} of rank r contains a sunflower \mathcal{S} with $r + 1$ petals with centre C , then $\mathcal{H} \setminus \mathcal{S} \cup \{C\}$ is also intersecting.

The Bollobás Theorem

Theorem (Bollobás): Let $\{A_1, \dots, A_t\}$ and $\{B_1, \dots, B_t\}$ be such that

- $A_i \cap B_i = \emptyset$ for each i ,
- $A_i \cap B_j \neq \emptyset$ for each $i \neq j$,
- $|A_i| \leq k$ and $|B_i| \leq l$ for each i .

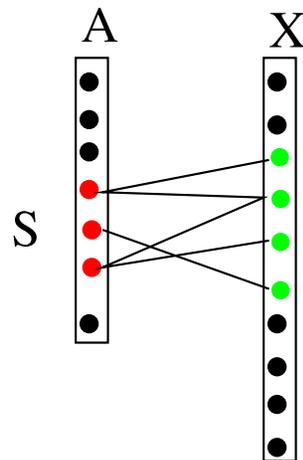
Then $t \leq \binom{k+l}{k}$.

Some proof ingredients

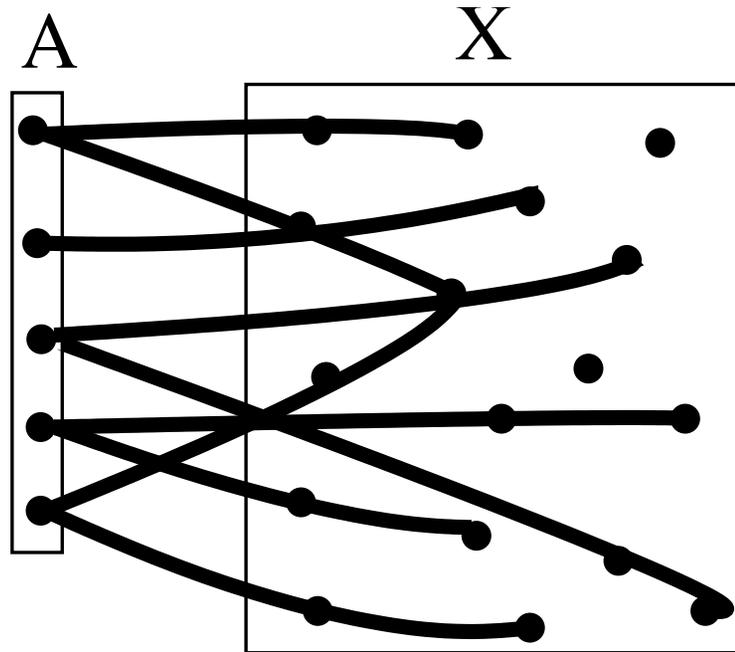
The **extremal result** for Ryser's conjecture for $r = 3$ initially follows Aharoni's proof of the conjecture for $r = 3$, which uses **Hall's Theorem for hypergraphs** together with König's Theorem.

Hall's Theorem: The **bipartite graph** G has a **complete matching** if and only if: For every subset $S \subseteq A$, the **neighbourhood** $\Gamma(S)$ is **big enough**.

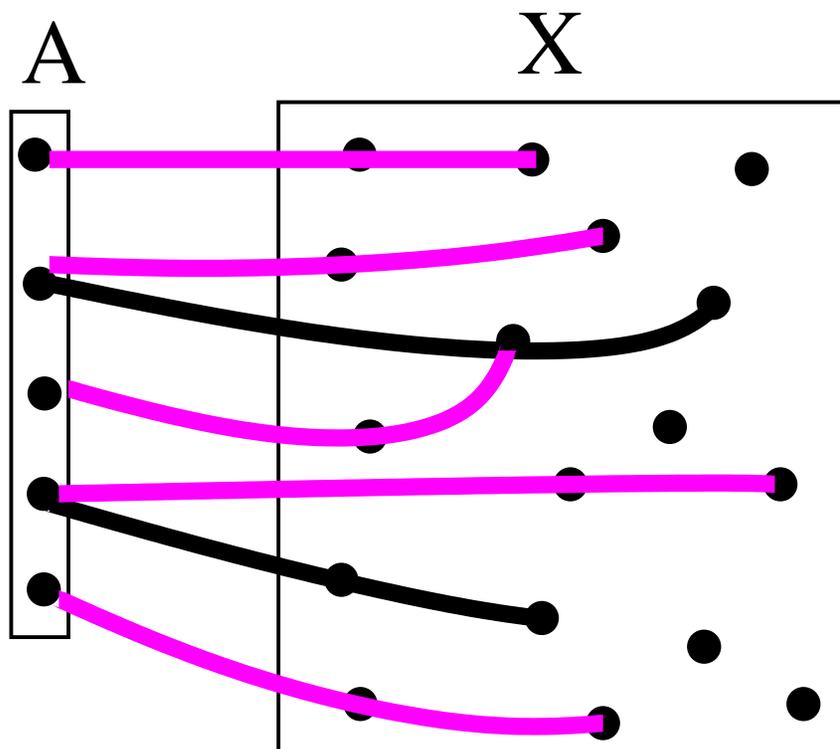
Here **big enough** means $|\Gamma(S)| \geq |S|$.



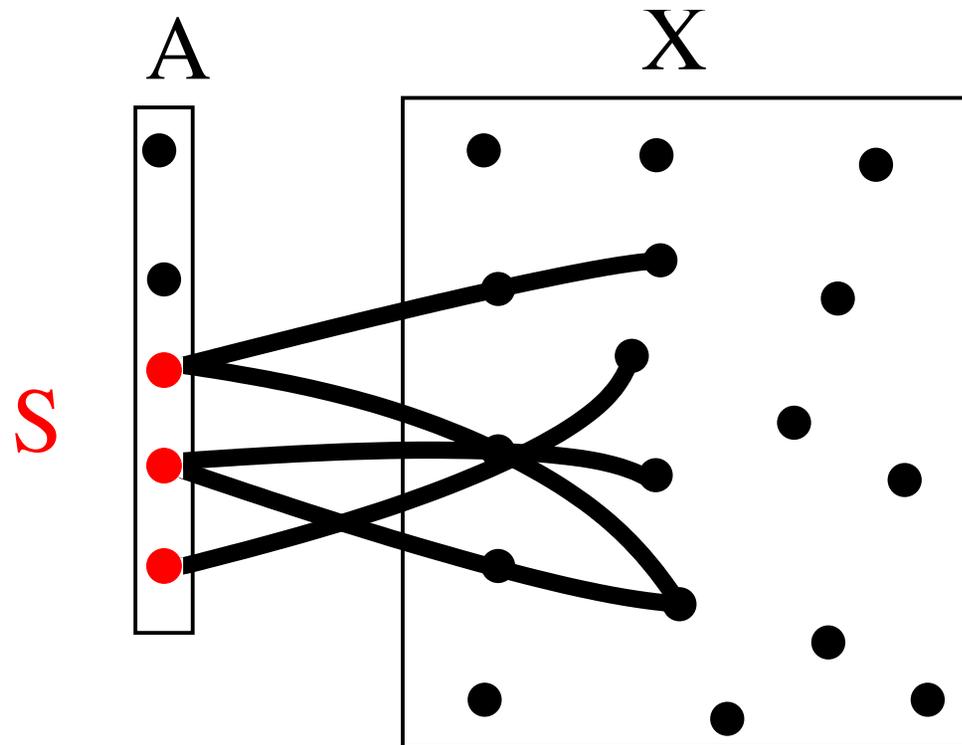
bipartite 3-uniform hypergraphs

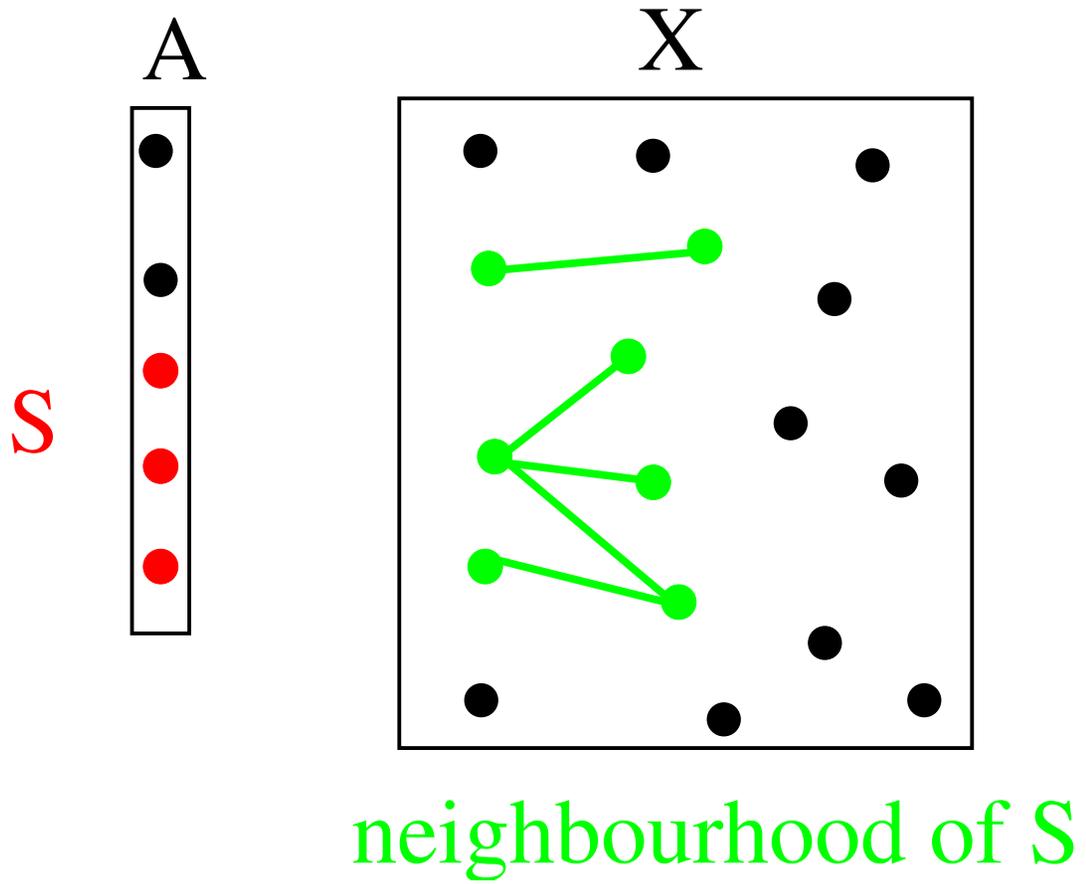


def: A complete packing:

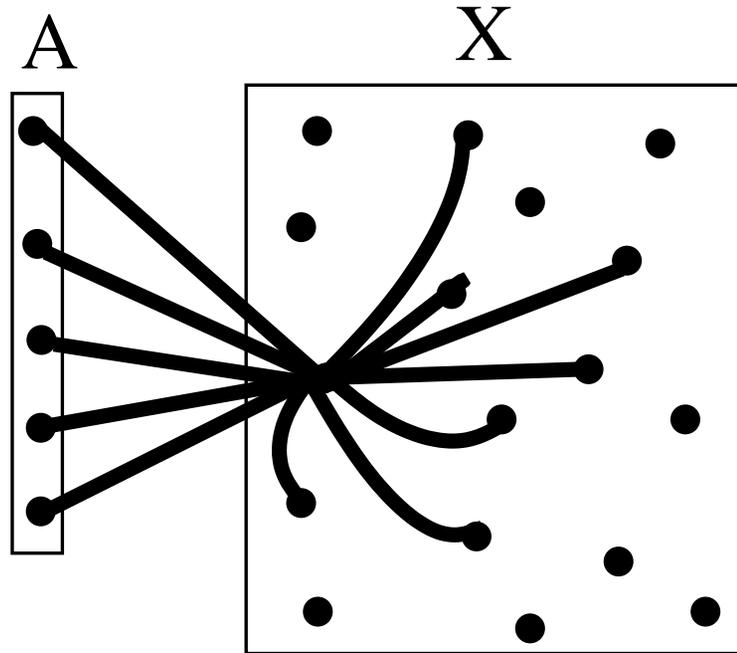


def: The **neighbourhood** of the subset S of A is the **graph** with vertex set X and edge set $\{\{x, y\} : \{z, x, y\} \in H \text{ for some } z \in S\}$.



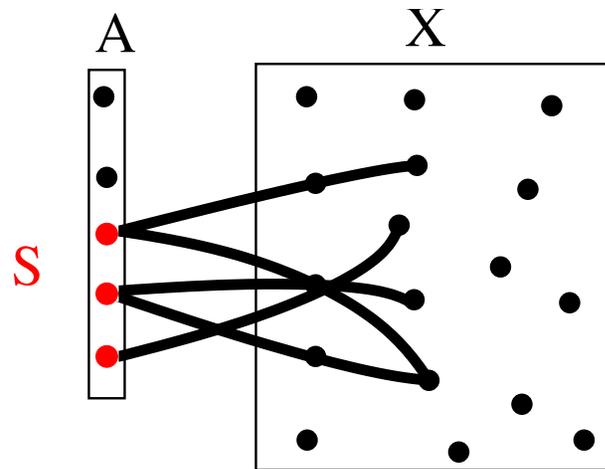


What should **big enough** mean?

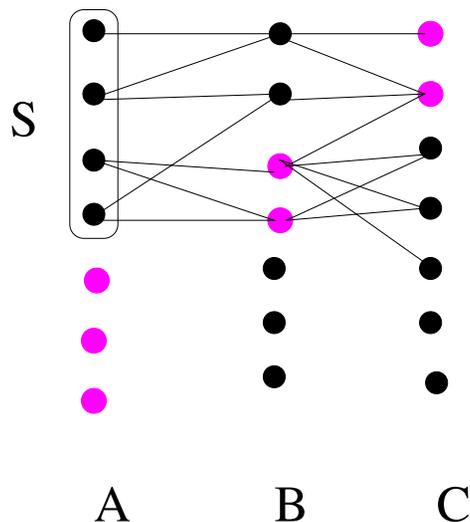


Hall's Theorem for 3-uniform hypergraphs

Theorem (Aharoni, PH, 2000): The bipartite 3-uniform hypergraph H has a **complete packing** if: For every subset $S \subseteq A$, the **neighbourhood** $\Gamma(S)$ has a matching of size at least $2(|S| - 1) + 1$.



Aharoni's proof of Ryser for $r = 3$



Let H be a 3-partite 3-uniform hypergraph. Let $\tau = \tau(H)$. Then by König's Theorem, for every subset S of A , the neighbourhood graph $\Gamma(S)$ has a matching of size at least $|S| - (|A| - \tau)$.

Then by a defect version of Hall's Theorem for hypergraphs, we find that H has a packing of size $\lceil \tau/2 \rceil$.

Proof of Hall's Theorem for hypergraphs

The proof has two main steps.

Step 1: The bipartite 3-uniform hypergraph H has a complete packing if: For every subset $S \subseteq A$, the topological connectivity of the matching complex of the neighbourhood graph $\Gamma(S)$ is at least $|S| - 2$.

Step 2: If the graph G has a matching of size at least $2(|S| - 1) + 1$ then the topological connectivity of the matching complex of G is at least $|S| - 2$.

The matching complex of G is the abstract simplicial complex with vertex set $E(G)$, whose simplices are the matchings in G .

Topological connectivity

One way to describe topological connectivity of an abstract simplicial complex Σ , as it is used here:

We say Σ is **k -connected** if for each $-1 \leq d \leq k$ and each triangulation T of the boundary of a $(d+1)$ -simplex, and each function f that labels each point of T with a point of Σ such that the set of labels on each simplex of T forms a simplex of Σ , the triangulation T can be extended to a triangulation T' of the whole $(d+1)$ -simplex, and f can be extended to a full labelling f' of T' with the same property.

Hall's Theorem for hypergraphs uses this together with Sperner's Lemma.

The topological connectivity of the matching complex of G is **not a monotone parameter**.

Extremal hypergraphs for Ryser's Conjecture

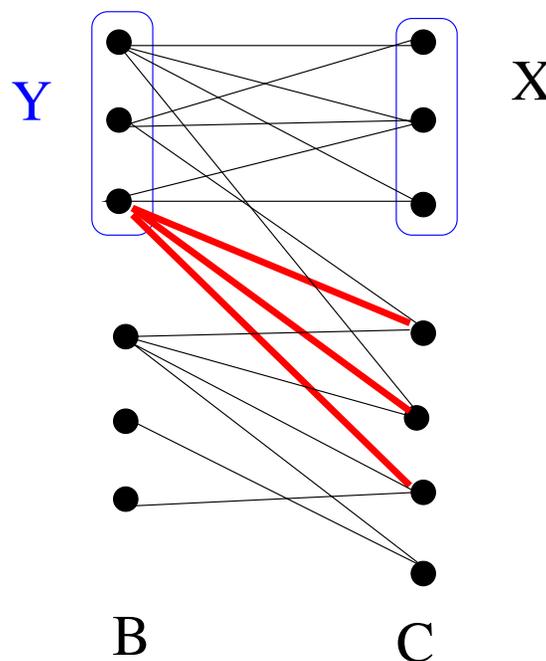
Two main parts are needed in understanding the extremal hypergraphs for Ryser's Conjecture for $r = 3$.

Part A: Show that any bipartite graph G that has a matching of size $2k$ but whose matching complex has the smallest possible topological connectivity (namely $k - 2$) has a very special structure.

Part B: Analyse how the edges of the neighbourhood graph G of A (which has this special structure) extend to A .

Part B (one case)

There exists a subset X of C with $|Y| \leq |X|$, where $Y = \Gamma_G(X)$, such that for each $y \in Y$, if we erase the $(y, C \setminus X)$ edges of G , the topological connectivity of the matching complex goes up.



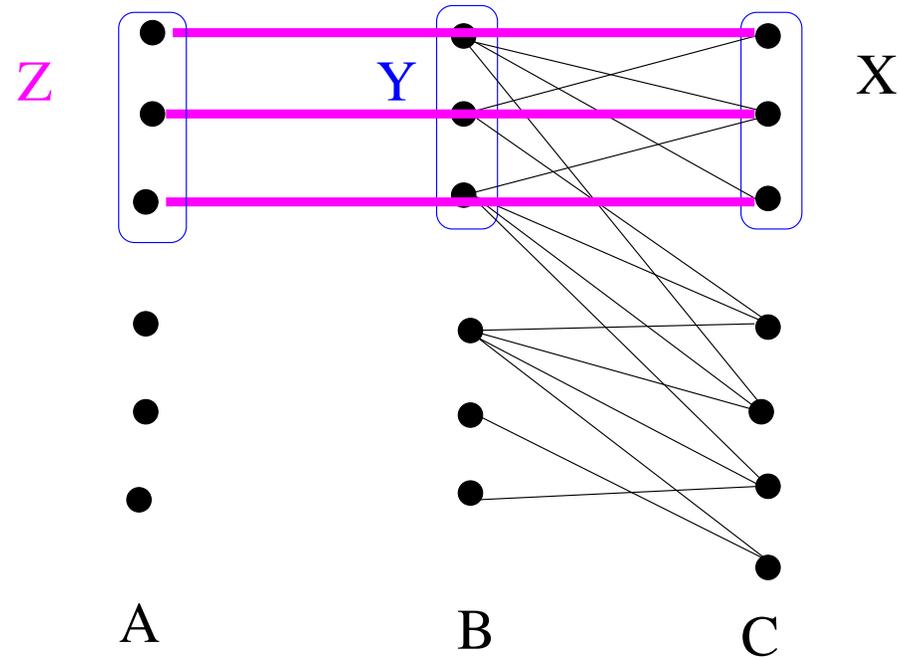
If for each $S \subset A$, the topological connectivity of the matching complex of $\Gamma(S)$ did not go down, then we find H has a packing larger than $\nu(H)$.

So for some S_y , erasing the $(y, C \setminus X)$ edges causes the connectivity to decrease.

Properties of S_y :

- $|S_y| \geq |A| - 1$, which implies $S_y = A \setminus \{a\}$ for some $a \in A$,
- every maximum matching in $\Gamma(S)$ uses an edge of $(y, C \setminus X)$.

What these properties imply



Removing the vertices in Y and Z causes ν to decrease by $|Y|$ and τ to decrease by $2|Y|$. Then we may use induction.

Triangle hypergraphs

Let G be a graph. The triangle hypergraph $\mathcal{H}(G)$ of G is the 3-uniform hypergraph with vertex set $E(G)$. Three edges of G form a hyperedge of $\mathcal{H}(G)$ if and only if they form the edge set of a triangle in G .

Thus a packing in $\mathcal{H}(G)$ is a set of edge-disjoint triangles in G .

A cover in $\mathcal{H}(G)$ is a set S of edges in G such that every triangle contains an edge in S .

Tuza's Conjecture

Conjecture (Tuza 1984): Let \mathcal{H} be a triangle hypergraph. Then

$$\tau(\mathcal{H}) \leq 2\nu(\mathcal{H}).$$

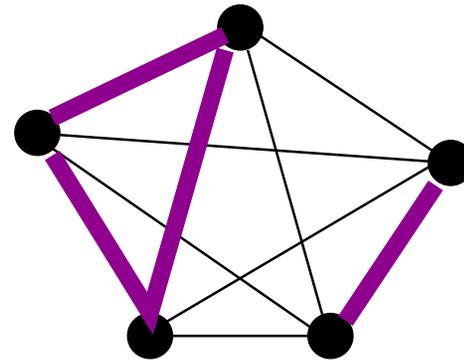
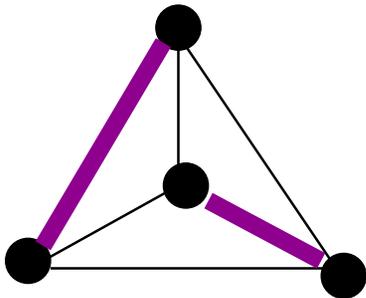
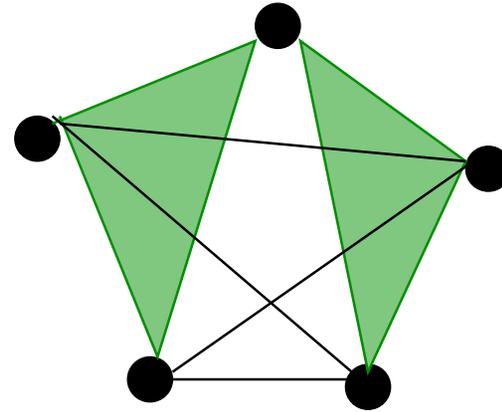
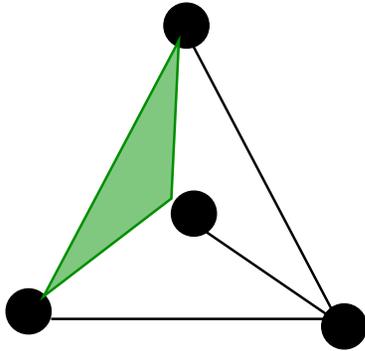
In other words:

Conjecture (Tuza 1984): Suppose the maximum size of a set of pairwise edge-disjoint triangles in a graph G is ν . Then there exists a set of at most 2ν edges in G whose removal makes the graph triangle-free.

Results on Tuza's Conjecture

- known for certain special classes of graphs, including K_5 -free chordal graphs (Tuza 1990), odd-wheel-free and four-colourable graphs (Aparna Lakshmanan, Bujtás and Tuza 2011)
- known for planar graphs (Tuza 1990), and more generally graphs without subdivisions of $K_{3,3}$ (Krivelevich 1995)
- weighted versions of the problem have been studied (Chapuy, DeVos, McDonald, Mohar and Scheide 2011)
- for every graph G the triangle hypergraph \mathcal{H} satisfies $\tau(\mathcal{H}) \leq (3 - \frac{3}{19})\nu(\mathcal{H})$.
- If true, Tuza's Conjecture is best possible.

Tuza's Conjecture



Fractional versions

Let \mathcal{H} be a hypergraph. A **fractional packing** of \mathcal{H} is a function p that assigns to each hyperedge e of \mathcal{H} a non-negative real number, such that for each vertex v of \mathcal{H} we have

$$\sum_{e \ni v} p(e) \leq 1.$$

Thus a **packing** (a set \mathcal{S} of disjoint hyperedges) corresponds to a fractional packing in which each hyperedge in \mathcal{S} gets value 1 and all others get 0.

A **fractional cover** of \mathcal{H} is a function c that assigns to each vertex of \mathcal{H} a non-negative real number, such that for each hyperedge e of \mathcal{H} we have

$$\sum_{v \in e} c(v) \geq 1.$$

Thus a **cover** of \mathcal{H} (a set C of vertices that meets every hyperedge) corresponds to a fractional cover in which each vertex in C gets value 1 and all other vertices get 0.

The fractional parameter $\nu^*(\mathcal{H})$ is defined to be the maximum of $\sum_{e \in \mathcal{H}} p(e)$ over all fractional packings p of \mathcal{H} .

The parameter $\tau^*(\mathcal{H})$ is the minimum of $\sum_{v \in \mathcal{H}} c(v)$ over all fractional covers c of \mathcal{H} .

Then we know that $\nu(\mathcal{H}) \leq \nu^*(\mathcal{H})$ and $\tau(\mathcal{H}) \geq \tau^*(\mathcal{H})$.

The Duality Theorem of linear programming tells us that

$$\tau^*(\mathcal{H}) = \nu^*(\mathcal{H}).$$

Fractional versions

Theorem (Krivelevich 1995): Let \mathcal{H} be a triangle hypergraph. Then

- $\tau^*(\mathcal{H}) \leq 2\nu(\mathcal{H})$.
- $\tau(\mathcal{H}) \leq 2\nu^*(\mathcal{H})$.

A closer look - the role of K_4

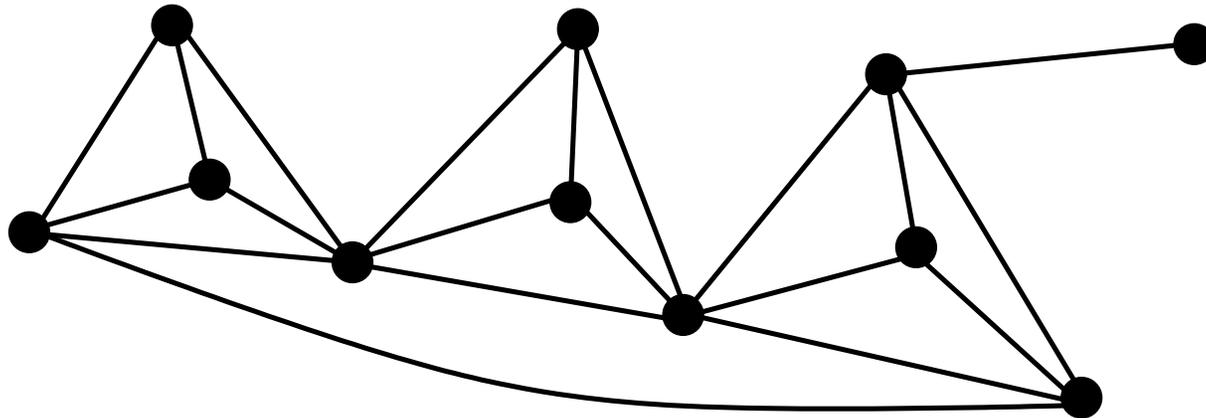
- (A)** Tuza's Conjecture is true for planar graphs, and best possible because of K_4 . What can we say about planar graphs for which $\tau(\mathcal{H})$ is close to $2\nu(\mathcal{H})$? Are they close to being disjoint unions of K_4 's?
- (B)** The fractional result $\tau^*(\mathcal{H}) \leq 2\nu(\mathcal{H})$ of Krivelevich is best possible because of K_4 . What can we say about graphs for which $\tau^*(\mathcal{H})$ is close to $2\nu(\mathcal{H})$? Are they close to being disjoint unions of K_4 's?

On Question (A)

Theorem (Cui, PH, Ma 2009) Let \mathcal{H} be the triangle hypergraph of a planar graph G , and suppose

$$\tau(\mathcal{H}) = 2\nu(\mathcal{H}).$$

Then G is an **edge-disjoint union of K_4 's and edges**, such that every triangle is contained in exactly one of the K_4 's.

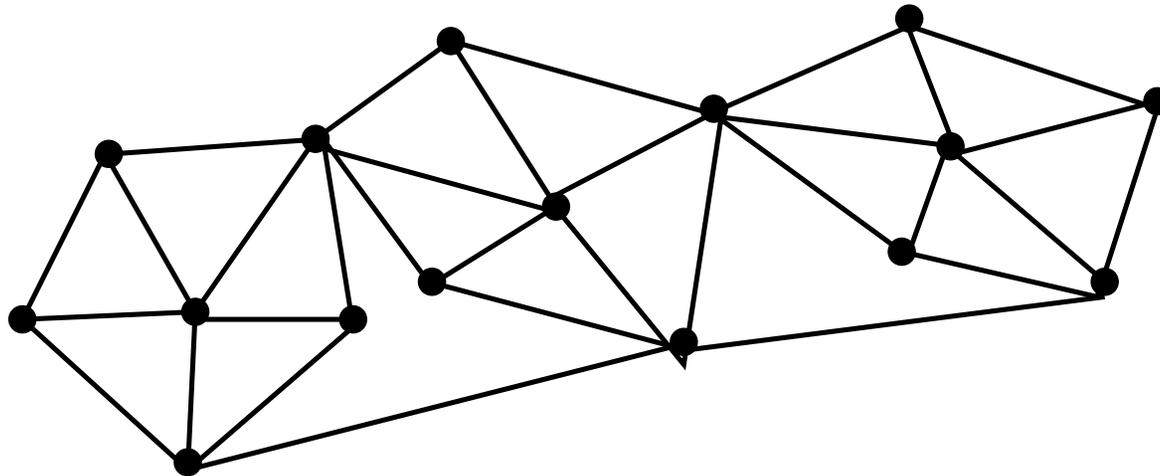


On Question (A)

Theorem (PH, Kostochka, Thomassé 2011) Let \mathcal{H} be the triangle hypergraph of a K_4 -free planar graph G . Then

$$\tau(\mathcal{H}) \leq \frac{3}{2} \nu(\mathcal{H}).$$

Moreover if equality holds then G is an edge-disjoint union of 5-wheels (plus possibly some edges that are not in triangles).



(B): A stability theorem

Theorem (PH, Kostochka, Thomassé 2011) Let G be a graph such that the triangle hypergraph \mathcal{H} satisfies $\tau^*(\mathcal{H}) \geq 2\nu(\mathcal{H}) - x$. Then G contains $\nu(\mathcal{H}) - \lfloor 10x \rfloor$ edge-disjoint K_4 -subgraphs plus an additional $\lfloor 10x \rfloor$ edge-disjoint triangles.

Note that just these K_4 's and triangles witness the fact that

$$\tau^*(\mathcal{H}) \geq 2\nu(\mathcal{H}) - \lfloor 10x \rfloor.$$

The proof also shows that if G is K_4 -free then

$$\tau^*(\mathcal{H}) \leq 1.8\nu(\mathcal{H}).$$

Stability for Tuza's conjecture

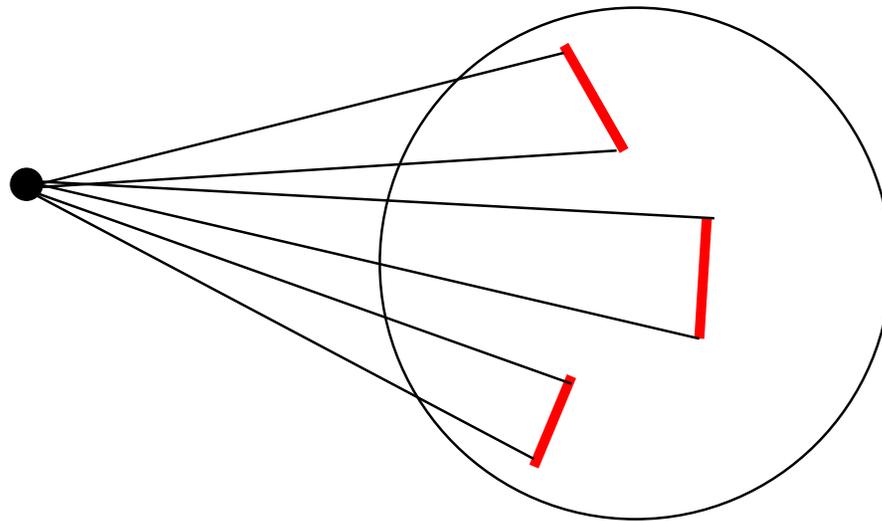
Could there be a similar stability theorem for Tuza's Conjecture?

The only known graphs for which equality holds for Tuza's Conjecture are (disjoint unions of) K_4 and K_5 . Could it be true that every graph for which $\tau(\mathcal{H})$ of the triangle hypergraph is close to $2\nu(\mathcal{H})$ contains many K_4 's?

NO.

For each $\epsilon > 0$, there exists a K_4 -free graph G_ϵ such that the triangle hypergraph \mathcal{H}_ϵ satisfies $\tau(\mathcal{H}_\epsilon) > (2 - \epsilon)\nu(\mathcal{H}_\epsilon)$.

For large n , let J be an n -vertex triangle-free graph with independence number $\alpha(J) < n^{2/3}$. ($R(3, t)$ is of order $t^2 / \log t$.)



Form a graph G by adding a new vertex v_0 and joining it to all vertices in J .

Then a packing in the triangle hypergraph \mathcal{H} corresponds to a matching in J , so

$$\nu(\mathcal{H}) \leq n/2.$$

A cover in \mathcal{H} corresponds to the complement of an independent vertex set in J . Thus

$$\tau(\mathcal{H}) \geq n - n^{2/3}.$$